

# Tutorial

## *The Perturbation-Duality Scheme in Optimization*

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# Introduction

Duality is widely used in optimization

- ▶ linear programming  
(Lagrangian duality, including optimal transport, etc.)
- ▶ convex programming  
(Lagrangian duality in mathematical programming, minimal cost flow on a graph etc.)
- ▶ conic programming, semidefinite programming, etc.

## The **perturbation-duality scheme** (PDS)

- ▶ Introduced in Rockafellar [1974]
- ▶ Goal: systematically produce dual optimization problems from a given optimization problem by perturbation followed by conjugate duality

# Outline

PDS for linear programming (LP)

PDS for convex programming

PDS for pure integer linear programming (PILP)

Conclusion

# Outline of the presentation

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Fundamental LP duality theorem

LP weak duality through PDS

LP strong duality through PDS

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# Primal and dual problems in standard LP

$$\begin{array}{ccc} \text{dual problem} & & \text{primal problem} \\ \hline \sup & \langle b, \lambda \rangle & \inf \quad \langle x, k \rangle \\ \lambda \in \mathbb{R}^m & & x \in \mathbb{R}^n \\ \lambda^T A \leq k & & Ax = b \\ & & x \geq 0 \end{array}$$

$\leq$   
weak duality

# Strong duality for LP

Adapted from Conforti, Cornuéjols, and Zambelli [2014, Theorem 3.7, Proposition 3.9]

## Theorem

Given a matrix  $A \in \mathbb{R}^{m \times n}$  and vectors  $k \in \mathbb{R}^n$ ,  $b \in \mathbb{R}^m$ ,  
if  $\{x \in \mathbb{R}^n \mid Ax = b, x \geq 0\} \neq \emptyset$  or  $\{\lambda \in \mathbb{R}^m \mid \lambda^T A \leq k\} \neq \emptyset$ ,  
(that is, if the primal or the dual problem is feasible)

then we have

$$\begin{array}{ccc} \sup_{\substack{\lambda \in \mathbb{R}^m \\ \lambda^T A \leq k}} \langle b, \lambda \rangle & \stackrel{=}{=} & \inf_{\substack{x \in \mathbb{R}^n \\ Ax = b \\ x \geq 0}} \langle x, k \rangle \end{array}$$

**strong duality**



## Sketch of the proof

1. Introduce the **Lagrangian**  $\mathcal{L}: \mathbb{R}_+^n \times \mathbb{R}^m \rightarrow \mathbb{R}$  by

$$\mathcal{L}(x, \lambda) = \langle x, k \rangle + \langle b - Ax, \lambda \rangle, \quad \forall x \in \mathbb{R}_+^n, \lambda \in \mathbb{R}^m$$

2. Use **sup-inf inversion** inequality to get weak duality

$$\underbrace{\sup_{\lambda \in \mathbb{R}^m} \inf_{x \in \mathbb{R}_+^n} \mathcal{L}(x, \lambda)}_{\text{dual problem}} \leq \underbrace{\inf_{x \in \mathbb{R}_+^n} \sup_{\lambda \in \mathbb{R}^m} \mathcal{L}(x, \lambda)}_{\text{primal problem}}$$

3. Find a **saddle-point**  $(\bar{x}, \bar{\lambda}) \in \mathbb{R}_+^n \times \mathbb{R}^m$

$$\mathcal{L}(\bar{x}, \lambda) \leq \mathcal{L}(\bar{x}, \bar{\lambda}) \leq \mathcal{L}(x, \bar{\lambda}), \quad \forall x \in \mathbb{R}_+^n, \lambda \in \mathbb{R}^m$$

to prove strong duality, i.e.

$$\underbrace{\sup_{\lambda \in \mathbb{R}^m} \inf_{x \in \mathbb{R}_+^n} \mathcal{L}(x, \lambda)}_{\text{dual problem}} = \underbrace{\inf_{x \in \mathbb{R}_+^n} \sup_{\lambda \in \mathbb{R}^m} \mathcal{L}(x, \lambda)}_{\text{primal problem}}$$

Now, let us deduce the previous duality results  
from a perturbation duality scheme (PDS)

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# Steps of the perturbation-duality scheme

Rockafellar [1974]

1. **Perturb** a minimization problem with a perturbation (primal) variable belonging to a vector space, and a Rockafellian function
2. **Pair** the (primal) perturbation space with a dual space by means of a bilinear form  $\langle , \rangle$
3. **Biconjugate** the perturbation function, and get
  - ▶ a dual problem
  - ▶ weak duality
4. Deduce conditions for **strong duality** by means of either global or local properties of the perturbation function

# Illustration of the scheme in Linear Programming (LP)

- ▶ Constraint matrix  $A \in \mathbb{R}^{m \times n}$
- ▶ Cost vector  $k \in \mathbb{R}^n$
- ▶ Anchor  $\bar{b} \in \mathbb{R}^m$

Initial/original minimization problem

$$\begin{aligned} \inf \quad & \langle x, k \rangle \\ x \in & \mathbb{R}^n \\ Ax = & \bar{b} \\ x \geq & 0 \end{aligned}$$

## Step 1. Perturbation of the initial minimization problem

- ▶ Introduce a perturbation space,  $\mathbb{R}^m$ , and embed the original problem into a family of minimization problems (more on the **Rockafellian** later)
- ▶ Introduce the **perturbation function**

$$\varphi: \mathbb{R}^m \rightarrow \overline{\mathbb{R}} = \mathbb{R} \cup \underbrace{\{-\infty\}}_{\text{unbounded}} \cup \underbrace{\{+\infty\}}_{\text{unfeasible}}$$

$$\forall b \in \mathbb{R}^m, \varphi(b) = \inf_{\substack{x \in \mathbb{R}^n \\ Ax = b \\ x \geq 0}} \langle x, k \rangle$$

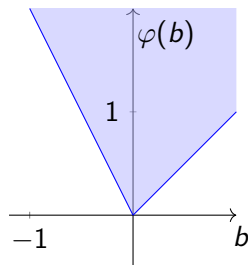
- ▶ The value of the original problem is then  $\varphi(\bar{b})$

## Example of perturbation function's epigraph for LP

Let  $\varphi: \mathbb{R} \rightarrow \mathbb{R}$  be defined as

$$\forall b \in \mathbb{R}, \varphi(b) = \inf_{\substack{x \in \mathbb{R}^2 \\ x_1 - x_2 = b \\ x \geq 0}} x_1 + 2x_2$$

Then  $\varphi(b) = \max\{-2b, b\}$ ,  $\forall b \in \mathbb{R}$



## Step 2. Dual space, coupling and conjugate function

- ▶ Perturbation space:  $\mathbb{R}^m$  — dual space:  $\mathbb{R}^m$  (linear functions)
- ▶ Introduce the bilinear coupling

$$\langle \cdot, \cdot \rangle : \overbrace{\mathbb{R}^m}^{\text{perturbation space}} \times \overbrace{\mathbb{R}^m}^{\text{dual space}} \rightarrow \mathbb{R}$$

- ▶ Deduce the **conjugate function**  $\varphi^* : \mathbb{R}^m \rightarrow \overline{\mathbb{R}}$  of the perturbation function

$$\forall \lambda \in \mathbb{R}^m, \varphi^*(\lambda) = \sup_{b \in \mathbb{R}^m} \{ \langle b, \lambda \rangle - \varphi(b) \}$$



# Conjugate function and Lagrangian

$$\begin{aligned}\varphi^*(\lambda) &= \sup_{b \in \mathbb{R}^m} \{ \langle b, \lambda \rangle - \varphi(b) \} \\ &= \sup_{b \in \mathbb{R}^m} \{ \langle b, \lambda \rangle - \inf_{\substack{Ax=b \\ x \geq 0}} \langle x, k \rangle \} \\ &= \sup_{b \in \mathbb{R}^m} \{ \langle b, \lambda \rangle + \sup_{\substack{Ax=b \\ x \geq 0}} \langle -x, k \rangle \} \\ &= \sup_{x \geq 0} \left\{ \sup_{\substack{Ax=b \\ b \in \mathbb{R}^m}} \langle b, \lambda \rangle - \langle x, k \rangle \right\} \\ &= \sup_{x \geq 0} \{ \langle Ax, \lambda \rangle - \langle x, k \rangle \} \\ &= \langle \bar{b}, \lambda \rangle - \inf_{x \geq 0} \underbrace{\{ \langle x, k \rangle + \langle \bar{b} - Ax, \lambda \rangle \}}_{\text{Lagrangian } \mathcal{L}(x, \lambda)}\end{aligned}$$

### Step 3. Biconjugate and weak duality

- ▶ Biconjugate function  $\varphi^{**'} : \mathbb{R}^m \rightarrow \bar{\mathbb{R}}$

$$\forall b \in \mathbb{R}^m, \varphi^{**'}(b) = \sup_{\lambda \in \mathbb{R}^m} \{ \langle b, \lambda \rangle - \varphi^*(\lambda) \}$$

- ▶ We obtain **weak duality** for all  $b \in \mathbb{R}^m$

$$\underbrace{\sup_{\substack{\lambda \in \mathbb{R}^m \\ \lambda^T A \leq k}} \langle b, \lambda \rangle}_{\text{dual problem}} = \varphi^{**'}(b) \leq \varphi(b) = \inf_{\substack{x \in \mathbb{R}^n \\ Ax = b \\ x \geq 0}} \langle x, k \rangle$$

- ▶ At the anchor  $\bar{b}$

$$\varphi^{**'}(\bar{b}) = \sup_{\lambda \in \mathbb{R}^m} \underbrace{\{ \langle \bar{b}, \lambda \rangle - \varphi^*(\lambda) \}}_{\inf_{x \geq 0} \mathcal{L}(x, \lambda)}$$

We have obtained the LP weak duality result  
What about the LP strong duality result?

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## Step 4. Conditions for strong duality

### Proposition

Let  $A \in \mathbb{R}^{m \times n}$  and  $k \in \mathbb{R}^n$

If  $-\infty < \varphi(0)$ , that is, **the corresponding LP is bounded**

then, **for all  $b \in \mathbb{R}^m$ ,**

$$\left( \sup_{\substack{\lambda \in \mathbb{R}^m \\ \lambda^T A \leq k}} \langle b, \lambda \rangle \right) = \underbrace{\varphi^{**'}(b) = \varphi(b)}_{\text{strong duality}} = \left( \inf_{\substack{x \in \mathbb{R}^n \\ Ax = b \\ x \geq 0}} \langle x, k \rangle \right)$$

### Remark

*This result is true even if  $b \in \mathbb{R}^m$  is such that  $\varphi(b) = +\infty$ , meaning for any unfeasible LPs*

# Proof of strong duality for LP. Sketch of the proof

adapted from Rockafellar [1974, p.24]

- (a) We show that if the LP  $\varphi(0)$  is bounded, then every feasible LP is bounded

$$-\infty < \varphi(0) \implies \varphi \text{ is proper}$$

- (b) We show that  $\text{epi } \varphi$  is a closed convex set (by showing that  $\text{epi } \varphi$  is a polyhedron)

$$\varphi \text{ is a closed convex function}$$

- (c) We apply Fenchel-Moreau Theorem to get strong duality

$$\varphi^{**'} = \varphi$$

# Strong duality for LP. Step (a) Proper functions

## Definition

Let  $f : \mathbb{R}^m \rightarrow \bar{\mathbb{R}} = \mathbb{R} \cup \{-\infty, +\infty\}$

- ▶  $\text{dom } f = \{b \in \mathbb{R}^m : f(b) < +\infty\}$
- ▶ The function  $f$  is said to be **proper** if  $\text{dom } f \neq \emptyset$  and  $-\infty < f(b)$ ,  $\forall b \in \mathbb{R}^m$

## Lemma

If  $-\infty < \varphi(0)$  (the corresponding LP is bounded)

then the value function  $\varphi$  is **proper** (all feasible LPs are bounded)

## Idea of the proof

The recession cone of  $\{x \in \mathbb{R}^n : Ax = b\}$

is given by  $\{r \in \mathbb{R}^n : Ax = 0, r \geq 0\}$

Conforti, Cornuéjols, and Zambelli [2014, Proposition 3.15]

# Strong duality for LP. Step (b) Closed convex epigraph

## Definition

Let  $f : \mathbb{R}^m \rightarrow \overline{\mathbb{R}} = \mathbb{R} \cup \{-\infty, +\infty\}$

The epigraph of the function  $f$  is defined by

$$\text{epi } f = \{(b, t) \in \mathbb{R}^m \times \mathbb{R} : f(b) \leq t\}$$

## Proposition

Let  $A \in \mathbb{R}^{m \times n}$  and  $k \in \mathbb{R}^n$  define the value function  $\varphi : \mathbb{R}^m \rightarrow \overline{\mathbb{R}}$  by

$$\forall b \in \mathbb{R}^m, \quad \varphi(b) = \inf_{\substack{x \in \mathbb{R}^n \\ Ax = b \\ x \geq 0}} \langle x, k \rangle$$

Then **epi  $\varphi$  is a polyhedron**



## Proof that $\text{epi } \varphi$ is a polyhedron

$$A \in \mathbb{R}^{m \times n}, k \in \mathbb{R}^n$$

Let  $b \in \mathbb{R}^m$ , we assume that  $-\infty < \varphi(b)$

$$\varphi(b) \leq t$$

$$\iff \inf_{\substack{x \in \mathbb{R}^n \\ Ax = b \\ x \geq 0}} \langle x, k \rangle \leq t$$

$$\iff \min_{\substack{x \in \mathbb{R}^n \\ Ax = b \\ x \geq 0}} \langle x, k \rangle \leq t \quad (\text{as bounded feasible LPs are attained})$$

$$\iff \exists x \in \mathbb{R}^n \text{ s.t. } Ax = b, x \geq 0, \langle x, k \rangle - t \leq 0$$

$$\iff \text{epi } \varphi = \pi_{(b,t)} \left\{ (b, t, x) \in \mathbb{R}^m \times \mathbb{R} \times \mathbb{R}^n : \begin{cases} Ax = b \\ x \geq 0 \\ \langle x, k \rangle - t \leq 0 \end{cases} \right\}$$

Thus  $\text{epi } \varphi$  is the projection of a polyhedron

So,  $\text{epi } \varphi$  is a polyhedron

Rockafellar [1970, Theorem 19.3]

# Strong duality for LP. Step (c) Fenchel-Moreau theorem

## Definition

A function  $f : \mathbb{R}^m \rightarrow \overline{\mathbb{R}}$  is said to be **closed convex** Rockafellar [1974] if **EITHER** [ $f$  is proper **AND** epi  $f$  is a closed convex set] **OR**  $f \equiv +\infty$  **OR**  $f \equiv -\infty$

## Theorem

[Fenchel-Moreau Theorem]

A function  $f : \mathbb{R}^m \rightarrow \overline{\mathbb{R}}$  is closed convex if and only if  $f^{**'} = f$

So we have **strong duality**

$$\varphi^{**'} = \underbrace{\varphi}$$

as Steps (a) and (b) imply  
that  $\varphi$  is a closed function

# About closed convex functions: the case of valley functions

## Definition

Let  $C \subset \mathbb{R}^n$  be a closed convex set

Let  $f: \mathbb{R}^n \rightarrow \bar{\mathbb{R}}$  be a function

We say that  $f$  is a **valley function** if

$$f(u) = \begin{cases} -\infty & \text{if } u \in C \\ +\infty & \text{otherwise} \end{cases}$$

## Remark

*Valley functions have a closed convex epigraph*

**BUT** *are not closed convex functions*

*(except the cases  $f \equiv -\infty$  or  $f \equiv +\infty$ )*

## Example when strong duality is not achieved for LP

$$-\infty = \left( \begin{array}{cc} \sup & \lambda_1 \\ \lambda \in \mathbb{R}^2 & \\ \lambda_1 + \lambda_2 \leq -1 & \\ -\lambda_1 - \lambda_2 \leq 0 & \end{array} \right)$$

$$= \varphi^{**'}((1, 0)) < \varphi((1, 0)) =$$

$$\left( \begin{array}{cc} \inf & -x_1 \\ x \in \mathbb{R}^2 & \\ x_1 - x_2 = 1 & \\ x_1 - x_2 = 0 & \\ x \geq 0 & \end{array} \right) = +\infty$$

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# Summary of the perturbation-duality scheme for LP

Rockafellar [1974]

1. We **perturb** a minimization problem

$$\forall b \in \mathbb{R}^m, \varphi(b) = \inf_{\substack{x \in \mathbb{R}^n \\ Ax = b \\ x \geq 0}} \langle x, k \rangle$$

2. We pair the primal space  $\mathbb{R}^m$  and a dual space  $\mathbb{R}^m$

$$\langle \cdot, \cdot \rangle: \mathbb{R}^m \times \mathbb{R}^m \rightarrow \mathbb{R}$$

3. We **biconjugate** the perturbation function  $\varphi$

$$\left( \begin{array}{c} \sup \\ \lambda \in \mathbb{R}^m \\ \lambda^T A \leq k \end{array} \langle b, \lambda \rangle \right) = \underbrace{\varphi^{**'}(b) \leq \varphi(b), \forall b \in \mathbb{R}^m}_{\text{weak duality is guaranteed}}$$

4. Under suitable assumptions, **strong duality** by polyhedral property of the epigraph of the perturbation function

# Indicator function of a subset

For any subset  $X \subset \mathcal{X}$ , its indicator function  $\iota_X$  is

$$\iota_X(x) = \begin{cases} 0 & \text{if } x \in X \\ +\infty & \text{if } x \notin X \end{cases}$$

# Summary (bis) of the perturbation-duality scheme for LP

1. We **perturb** a minimization problem

$$\varphi(b) = \inf_{x \in \mathbb{R}^n} \mathcal{R}(x, b), \quad \forall b \in \mathbb{R}^m$$

where the **Rockafellian**  $\mathcal{R}: \mathbb{R}^n \times \mathbb{R}^m \rightarrow \bar{\mathbb{R}}$  is defined by  $\mathcal{R}(x, b) = \langle x, k \rangle + \iota_{\mathbb{R}_+^n}(u) + \iota_{\{0\}}(Ax - b)$

2. We pair the primal space  $\mathbb{R}^m$  and a dual space  $\mathbb{R}^m$

$$\langle \cdot, \cdot \rangle: \mathbb{R}^m \times \mathbb{R}^m \rightarrow \mathbb{R}$$

3. We **biconjugate** the perturbation function  $\varphi$

$$\left( \sup_{\substack{\lambda \in \mathbb{R}^m \\ \lambda^T A \leq k}} \langle b, \lambda \rangle \right) = \underbrace{\varphi^{**'}(b)}_{\text{weak duality is guaranteed}} \leq \varphi(b), \quad \forall b \in \mathbb{R}^m$$

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# The Fenchel conjugacy

## Definition

Two vector spaces  $\mathcal{U}$  and  $\mathcal{V}$ , paired by a bilinear form  $\langle \cdot, \cdot \rangle$  (in the sense of convex analysis), give rise to the classic **Fenchel conjugacy** between  $\overline{\mathbb{R}}^{\mathcal{U}}$  and  $\overline{\mathbb{R}}^{\mathcal{V}}$

With any function  $f: \mathcal{U} \rightarrow \overline{\mathbb{R}}$ , we associate the function  $f^*: \mathcal{V} \rightarrow \overline{\mathbb{R}}$  defined by

$$f^*(v) = \sup_{u \in \mathcal{U}} \{ \langle u, v \rangle - f(u) \}, \quad \forall v \in \mathcal{V}$$

# The biconjugate function is a minorant of the function

## Definition

Let  $f: \mathcal{U} \rightarrow \overline{\mathbb{R}}$  be a function

Its **biconjugate**  $f^{**'}: \mathcal{U} \rightarrow \overline{\mathbb{R}}$  is defined by

$$f^{**'}(u) = \sup_{v \in \mathcal{V}} \{ \langle u, v \rangle - f^*(v) \}$$

The inequality below is instrumental in obtaining weak duality

## Proposition

For any function  $f: \mathcal{U} \rightarrow \overline{\mathbb{R}}$ , we have that

$$f^{**'} \leq f$$

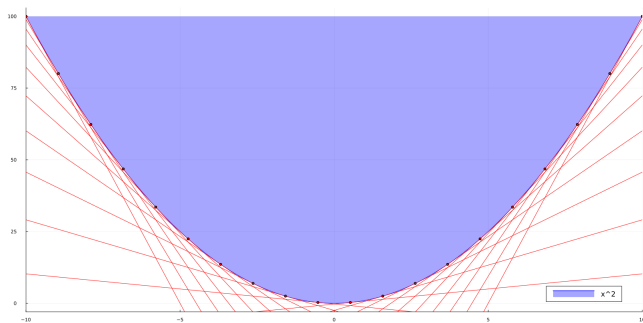
# Fenchel-Moreau Theorem

The equality below is instrumental in obtaining strong duality

## Theorem

[Fenchel-Moreau]

The function  $f: \mathcal{U} \rightarrow \overline{\mathbb{R}}$  is **closed convex** if and only if  $f^{**'} = f$



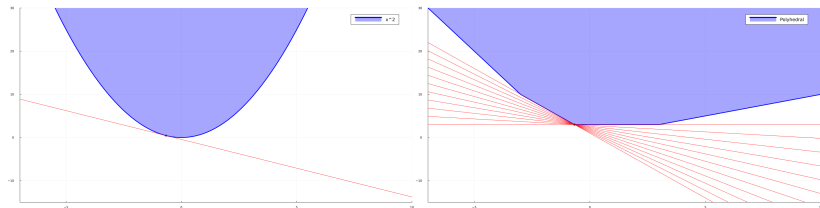
# Moreau-Rockafellar subdifferential

## Definition

Let  $f : \mathcal{U} \rightarrow \overline{\mathbb{R}}$  be a function

Its **subdifferential**  $\partial f(u) \subset \mathcal{V}$  at any  $u \in \mathcal{U}$  such that  $f(u) \in \mathbb{R}$ , is defined by

$$v \in \partial f(u) \iff \langle u', v \rangle - f(u') \leq \langle u, v \rangle - f(u), \quad \forall u' \in \mathcal{U}$$



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  - ▶ a dual problem
  - ▶ weak duality
4. Deduce conditions for **strong duality** by means of either global or local properties of the perturbation function



# Perturbation duality scheme Rockafellar [1974]

sets	optimization set $\mathcal{X}$	primal space $\mathcal{U}$	pairing $\mathcal{U} \overset{\langle \cdot, \cdot \rangle}{\leftrightarrow} \mathcal{V}$	dual space $\mathcal{V}$
variables	<b>decision</b> $x \in \mathcal{X}$	<b>perturbation</b> $u \in \mathcal{U}$	$\langle u, v \rangle$ $\in \mathbb{R}$	sensitivity $v \in \mathcal{V}$
bivariate functions		<b>Rockafellian</b> $\mathcal{R}: \mathcal{X} \times \mathcal{U} \rightarrow \overline{\mathbb{R}}$		<b>Lagrangian</b> $\mathcal{L}: \mathcal{X} \times \mathcal{V} \rightarrow \overline{\mathbb{R}}$
definition				$\mathcal{L}(x, v) =$ $\inf_{u \in \mathcal{U}} \{ \mathcal{R}(x, u) - \langle u, v \rangle \}$
property				$-\mathcal{L}(x, \cdot) = (\mathcal{R}(x, \cdot))^*$
property				$-\mathcal{L}(x, \cdot)$ is $\star'$ -convex (hence $\mathcal{L}(x, \cdot)$ is concave usc)
univariate functions		<b>perturbation function</b> $\varphi: \mathcal{U} \rightarrow \overline{\mathbb{R}}$		<b>dual function</b> $\psi: \mathcal{V} \rightarrow \overline{\mathbb{R}}$
definition		$\varphi(u) = \inf_{x \in \mathcal{X}} \mathcal{R}(x, u)$		$\psi(v) = \inf_{x \in \mathcal{X}} \mathcal{L}(x, v)$
property				$-\psi = \varphi^*$

## Weak duality

# Perturbation/Rockafellian (Step 1)

Data: set  $\mathcal{X}$ , function  $h: \mathcal{X} \rightarrow \overline{\mathbb{R}}$  and

original minimization problem  $\inf_{x \in \mathcal{X}} h(x)$

- ▶ Embedding/perturbation scheme given by a vector space  $\mathcal{U}$ , and a Rockafellian  $\mathcal{R}: \mathcal{X} \times \mathcal{U} \rightarrow \overline{\mathbb{R}}$  such that

$$h(x) = \mathcal{R}(x, 0), \quad \forall x \in \mathcal{X}$$

- ▶ The perturbation function  $\varphi: \mathcal{U} \rightarrow \overline{\mathbb{R}}$  is defined by

$$\varphi(u) = \inf_{x \in \mathcal{X}} \mathcal{R}(x, u)$$

original minimization problem  $\varphi(0) = \inf_{x \in \mathcal{X}} h(x)$

# Duality/Fenchel conjugacy (Steps 2,3)

- ▶ Dual vector space  $\mathcal{V}$  paired to  $\mathcal{U}$  by a bilinear form  $\langle \cdot, \cdot \rangle$

We obtain weak duality

$$\begin{aligned}\varphi^{**'}(0) &= \overbrace{\sup_{v \in \mathcal{V}} \{-\varphi^*(v)\}}^{\text{dual problem}} \\ &\leq \\ \varphi(0) &= \underbrace{\inf_{x \in \mathcal{X}} h(x)}_{\text{original problem}}\end{aligned}$$

## Weak duality and Lagrangian

# Lagrangian

- ▶ Lagrangian  $\mathcal{L}: \mathcal{X} \times \mathcal{V} \rightarrow \overline{\mathbb{R}}$  defined by

$$\mathcal{L}(x, v) = \inf_{u \in \mathcal{U}} \left\{ \underbrace{\mathcal{R}(x, u)}_{\text{Rockafellian}} - \langle u, v \rangle \right\}, \quad \forall (x, v) \in \mathcal{X} \times \mathcal{V}$$

- ▶ As  $\mathcal{L}(x, v) \leq \mathcal{R}(x, 0) - \langle 0, v \rangle = h(x)$ , we get that

$$\sup_{v \in \mathcal{V}} \mathcal{L}(x, v) \leq h(x)$$

hence that

original minimization problem

$$\inf_{x \in \mathcal{X}} \sup_{v \in \mathcal{V}} \mathcal{L}(x, v) \leq \inf_{x \in \mathcal{X}} h(x)$$

# Dual function

- ▶ The **dual function**  $\psi: \mathcal{V} \rightarrow \overline{\mathbb{R}}$  is defined by

$$\psi(v) = \inf_{x \in \mathcal{X}} \mathcal{L}(x, v), \quad \forall v \in \mathcal{V}$$

- ▶ and the **dual problem** is

$$\varphi^{**}(0) = \sup_{v \in \mathcal{V}} \{ \langle 0, v \rangle - \varphi^*(v) \} = \overbrace{\sup_{v \in \mathcal{V}} \psi(v)}^{\text{dual problem}}$$

$$\begin{aligned} \text{as } -\varphi^*(v) &= -\left( \inf_{x \in \mathcal{X}} \mathcal{R}(x, \cdot) \right)^*(v) \\ &= -\sup_{x \in \mathcal{X}} \left\{ \sup_{u \in \mathcal{U}} \{ \langle u, v \rangle - \mathcal{R}(x, u) \} \right\} \\ &= -\sup_{x \in \mathcal{X}} \left\{ \underbrace{-\inf_{u \in \mathcal{U}} \{ -\langle u, v \rangle + \mathcal{R}(x, u) \}}_{\text{Lagrangian}} \right\} \\ &= \inf_{x \in \mathcal{X}} \mathcal{L}(x, v) = \psi(v) \end{aligned}$$

# Weak duality with Lagrangian

$$\begin{aligned} & \varphi^{**'}(0) \\ &= \sup_{v \in \mathcal{V}} \{-\varphi^*(v)\} \\ &= \underbrace{\sup_{v \in \mathcal{V}} \inf_{x \in \mathcal{X}} \mathcal{L}(x, v)}_{\text{dual problem}} \\ &\leq \\ & \inf_{x \in \mathcal{X}} \sup_{v \in \mathcal{V}} \mathcal{L}(x, v) \\ &\leq \inf_{x \in \mathcal{X}} h(x) \\ &= \varphi(0) \end{aligned}$$



## Strong duality

# Strong duality

$$\begin{array}{ccc} \text{dual problem} & & \text{original problem} \\ \underbrace{\varphi^{**'}(0)} & \leq & \underbrace{\varphi(0)} \\ \text{weak duality} & & \end{array}$$

## Definition

Strong duality  $\iff \varphi^{**'}(0) = \varphi(0) \iff \varphi$  is  $\star$ -convex at 0

## Paths to strong duality in the convex case

- ▶ Suppose that the **Rockafellian**  $\mathcal{R}: \mathcal{X} \times \mathcal{U} \rightarrow \overline{\mathbb{R}}$  is a **jointly convex** function
- ▶ Then, the **perturbation function**  $\varphi: \mathcal{U} \rightarrow \overline{\mathbb{R}}$  is **convex** as the **marginal function**  $\varphi(u) = \inf_{x \in \mathcal{X}} \mathcal{R}(x, u)$
- ▶ If, in addition,
  - ▶ either (global property) the function  $\varphi$  is **proper and lower semicontinuous**, and then  $\varphi^{**'} = \varphi$  by the Fenchel-Moreau Theorem,
  - ▶ or (local property) the subdifferential  $\partial\varphi(0) \neq \emptyset$ , and then the function  $\varphi$  is  $\star$ -convex at 0,

and we get **strong duality**  $\underbrace{\varphi^{**'}(0)}_{\text{dual problem}} = \underbrace{\varphi(0)}_{\text{original problem}}$

# Outline of the presentation

PDS for linear programming (LP)

**PDS for convex programming**

Background on duality in convex analysis

PDS with Fenchel duality

**Examples of PDS for convex programs**

Generalized perturbation duality scheme

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Conclusion

## Lagrangian duality (the case of inequality constraints)

# Classic Lagrangian duality (the case of inequality constraints)

- ▶ Optimization set  $\mathcal{X}$
- ▶ Objective function  $h: \mathcal{X} \rightarrow ] - \infty, +\infty]$
- ▶ Mapping  $\theta = (\theta_1, \dots, \theta_p) : \mathcal{X} \rightarrow \mathbb{R}^p$ , and  $\bar{u} \in \mathbb{R}^p$

We consider the optimization problem

$$\min_{\theta(x) \leq \bar{u}} h(x) = \min_{\substack{\theta_1(x) \leq \bar{u}_1 \\ \theta_p(x) \leq \bar{u}_p}} h(x)$$

# Perturbation and Rockafellian

- ▶ Perturbation space  $\mathcal{U} = \mathbb{R}^p$
- ▶ Rockafellian  $\mathcal{R}: \mathcal{X} \times \mathbb{R}^p \rightarrow \overline{\mathbb{R}}$

$$\mathcal{R}(x, u) = h(x) + \iota_{\{\theta(x) - \bar{u} \leq u\}} = h(x) + \sum_{j=1}^p \iota_{\{\theta_j(x) - \bar{u}_j \leq u_j\}}$$

# Duality, Lagrangian and dual function

- ▶ Dual space  $\mathcal{V} = \mathbb{R}^P$
- ▶ We deduce the **Lagrangian**  $\mathcal{L}: \mathcal{X} \times \mathbb{R}^P \rightarrow \overline{\mathbb{R}}$

$$\mathcal{L}(x, v) = h(x) + \langle \theta(x) - \bar{u}, v \rangle = h(x) + \sum_{j=1}^P v_j (\theta_j(x) - \bar{u})$$

- ▶ We deduce the **dual function**  $\psi: \mathbb{R}^P \rightarrow \overline{\mathbb{R}}$

$$\psi(v) = \inf_{x \in \mathcal{X}} \mathcal{L}(x, v) = \inf_{x \in \mathcal{X}} \left\{ h(x) + \sum_{j=1}^P v_j (\theta_j(x) - \bar{u}) \right\}$$

which is concave upper semicontinuous,  
as the supremum of affine functions



# Paths to strong duality in the convex case

- ▶ Suppose that
  - ▶ the optimization set  $\mathcal{X}$  is a vector space
  - ▶ the objective function  $h: \mathcal{X} \rightarrow ]-\infty, +\infty]$  is convex
  - ▶ each component of the mapping  $\theta = (\theta_1, \dots, \theta_p): \mathcal{X} \rightarrow \mathbb{R}^p$  is a convex function
- ▶ Then, the **perturbation function**  $\varphi: \mathbb{R}^p \rightarrow \overline{\mathbb{R}}$  is a **convex** function as the marginal

$$\varphi(u) = \inf_{x \in \mathcal{X}} \left\{ h(x) + \sum_{j=1}^p v_j (\theta_j(x) - u) \right\}$$

- ▶ If, in addition,
  - ▶ either the function  $\varphi$  is **proper and lower semicontinuous**
  - ▶ or its subdifferential  $\partial\varphi(0) \neq \emptyset$

then we get **strong duality**

## Fenchel-Rockafellar duality

# The Fenchel-Rockafellar dual problem

## Proposition

adapted from Rockafellar [1970, Corollary 31.2.1]

Let  $f, g: \mathbb{R}^n \rightarrow \overline{\mathbb{R}}$  be proper convex functions  
and let  $L: \mathbb{R}^n \rightarrow \mathbb{R}^m$  be a linear mapping

$$\sup_{v \in \mathbb{R}^m} \{-g^*(v) - f^*(L^T v)\} \leq \inf_{u \in \mathbb{R}^n} \{f(u) + g(Lu)\}$$

Furthermore, equality is achieved if either

- ▶  $\exists u \in \text{ri}(\text{dom} f)$  s.t.  $Lu \in \text{ri}(\text{dom} g)$
- ▶  $\exists v \in \text{ri}(\text{dom} g^*)$  s.t.  $L^T v \in \text{ri}(\text{dom} f^*)$

# Perturbation and Rockafellian TODO

- ▶ Perturbation space  $\mathcal{U} = \mathbb{R}^p$
- ▶ Rockafellian  $\mathcal{R}: \mathcal{X} \times \mathbb{R}^p \rightarrow \overline{\mathbb{R}}$

$$\mathcal{R}(x, u) = h(x) + \iota_{\{\theta(x) - \bar{u} \leq u\}} = h(x) + \sum_{j=1}^p \iota_{\{\theta_j(x) - \bar{u}_j \leq u_j\}}$$

# Duality, Lagrangian and dual function TODO

- ▶ Dual space  $\mathcal{V} = \mathbb{R}^p$
- ▶ We deduce the **Lagrangian**  $\mathcal{L}: \mathcal{X} \times \mathbb{R}^p \rightarrow \overline{\mathbb{R}}$

$$\mathcal{L}(x, v) = h(x) + \langle \theta(x) - \bar{u}, v \rangle = h(x) + \sum_{j=1}^p v_j (\theta_j(x) - \bar{u})$$

- ▶ We deduce the **dual function**  $\psi: \mathbb{R}^p \rightarrow \overline{\mathbb{R}}$

$$\psi(v) = \inf_{x \in \mathcal{X}} \mathcal{L}(x, v) = \inf_{x \in \mathcal{X}} \left\{ h(x) + \sum_{j=1}^p v_j (\theta_j(x) - \bar{u}) \right\}$$

which is concave upper semicontinuous,  
as the supremum of affine functions

## Application to regularized problem

For a linear mapping  $L: \mathbb{R}^n \rightarrow \mathbb{R}^m$  and a proper convex function  $f: \mathbb{R}^m \rightarrow \overline{\mathbb{R}}$  and suitable assumptions

$$\sup_{v \in \mathbb{R}^m} -\|L^T v\|^2 - f^*(v) = \inf_{u \in \mathbb{R}^n} f(Lu) + \frac{1}{2} \|u\|^2$$

Can be useful for computation if  $m < n$  and  $f^*$  easy to compute

# Semidefinite programming dual problem

adapted from Calafiore and El Ghaoui [2014, Chapter 11]

Let  $\mathbb{S}^n$  be the set of  $n \times n$  symmetric matrices

Let  $\mathbb{S}_+^n \subset \mathbb{S}^n$  be the set of  $n \times n$  semidefinite matrices

## Proposition

Let  $K, A_1, \dots, A_m \in \mathbb{S}^n$  and  $b \in \mathbb{R}^m$

Then, we have

$$\sup_{\substack{v \in \mathbb{R}^m \\ K - \sum_{j=1}^m v_j A_j \succeq 0}} \langle b, v \rangle \leq \inf_{\substack{X \in \mathbb{S}^n \\ \text{trace } A_j X = b_j, j=1, \dots, m \\ X \succeq 0}} \text{trace } KX$$

Furthermore, equality is achieved  
if some Slater's condition is satisfied

# Perturbation and Rockafellian TODO

- ▶ Perturbation space  $\mathcal{U} = \mathbb{R}^p$
- ▶ Rockafellian  $\mathcal{R}: \mathcal{X} \times \mathbb{R}^p \rightarrow \overline{\mathbb{R}}$

$$\mathcal{R}(x, u) = h(x) + \iota_{\{\theta(x) - \bar{u} \leq u\}} = h(x) + \sum_{j=1}^p \iota_{\{\theta_j(x) - \bar{u}_j \leq u_j\}}$$



# Duality, Lagrangian and dual function TODO

- ▶ Dual space  $\mathcal{V} = \mathbb{R}^p$
- ▶ We deduce the **Lagrangian**  $\mathcal{L}: \mathcal{X} \times \mathbb{R}^p \rightarrow \overline{\mathbb{R}}$

$$\mathcal{L}(x, v) = h(x) + \langle \theta(x) - \bar{u}, v \rangle = h(x) + \sum_{j=1}^p v_j (\theta_j(x) - \bar{u})$$

- ▶ We deduce the **dual function**  $\psi: \mathbb{R}^p \rightarrow \overline{\mathbb{R}}$

$$\psi(v) = \inf_{x \in \mathcal{X}} \mathcal{L}(x, v) = \inf_{x \in \mathcal{X}} \left\{ h(x) + \sum_{j=1}^p v_j (\theta_j(x) - \bar{u}) \right\}$$

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# Introducing generalized convexity

<b>Fenchel conjugate</b> $f^*(v) = \sup_{u \in \mathbb{R}^m} \langle u, v \rangle - f(u)$	<b><math>c</math>-conjugate</b> $g^c(v) = \sup_{u \in U} c(u, v) \dot{+} (-g(u))$
<b>Fenchel biconjugate</b> $f^{**'}(u) = \sup_{v \in \mathbb{R}^m} \langle u, v \rangle - f^*(v)$	<b><math>c</math>-biconjugate</b> $g^{cc'}(u) = \sup_{v \in V} c(u, v) \dot{+} (-g^c(v))$
<b><math>\star</math> - convex functions</b> $\iff f = f^{**'}$	<b><math>c</math>-convex functions</b> $\iff g = g^{cc'}$

with the Moreau lower and upper additions

$$(+\infty) \dot{+} (-\infty) = (-\infty) \dot{+} (+\infty) = -\infty$$

$$(+\infty) \dot{+} (-\infty) = (-\infty) \dot{+} (+\infty) = +\infty$$

# Generalized perturbation-duality scheme Balder [1977]

sets	optimization set $\mathcal{X}$	primal set $\mathcal{U}$	coupling $\mathcal{U} \overset{c}{\leftrightarrow} \mathcal{V}$	dual set $\mathcal{V}$
variables	decision $x \in \mathcal{X}$	perturbation $u \in \mathcal{U}$	$c(u, v) \in \mathbb{R}$	sensitivity $v \in \mathcal{V}$
bivariate functions		Rockafellian $\mathcal{R}: \mathcal{X} \times \mathcal{U} \rightarrow \overline{\mathbb{R}}$		Lagrangian $\mathcal{L}: \mathcal{X} \times \mathcal{V} \rightarrow \overline{\mathbb{R}}$
definition				$\mathcal{L}(x, v) = \inf_{u \in \mathcal{U}} \{ \mathcal{R}(x, u) + (-c(u, v)) \}$
property				$-\mathcal{L}(x, \cdot) = (\mathcal{R}(x, \cdot))^c$
property				$-\mathcal{L}(x, \cdot)$ is $c'$ -convex
univariate functions		perturbation function $\varphi: \mathcal{U} \rightarrow \overline{\mathbb{R}}$		dual function $\psi: \mathcal{V} \rightarrow \overline{\mathbb{R}}$
definition		$\varphi(u) = \inf_{x \in \mathcal{X}} \mathcal{R}(x, u)$		$\psi(v) = \inf_{x \in \mathcal{X}} \mathcal{L}(x, v)$
property				$-\psi = \varphi^c$

- ▶ Anchor  $\bar{u} \in \mathcal{U}$  and dual maximization problem (weak duality)  
 $\varphi^{cc'}(\bar{u}) = \sup_{v \in \mathcal{V}} \{ c(\bar{u}, v) + \psi(v) \} \leq \inf_{x \in \mathcal{X}} h(x) = \varphi(\bar{u})$
- ▶ Strong duality iff  $\varphi$  is  $c$ -convex at  $\bar{u}$  iff  $\varphi^{cc'}(\bar{u}) = \varphi(\bar{u})$

## Case of evaluation couplings (developped later)

- ▶ Given a primal set  $\mathcal{U}$  and a function set  $\mathcal{F} \subset \{F: \mathcal{U} \rightarrow \overline{\mathbb{R}}\}$ , the **evaluation coupling**  $c_{\mathcal{F}}: \mathcal{U} \times \mathcal{F} \rightarrow \overline{\mathbb{R}}$  is defined by

$$c_{\mathcal{F}}(u, F) = F(u), \quad \forall u \in \mathcal{U}, F \in \mathcal{F}$$

- ▶ For a given (perturbation) function  $\varphi: \mathcal{U} \rightarrow \overline{\mathbb{R}}$ , **weak duality** is always achieved

$$\varphi^{c_{\mathcal{F}}c_{\mathcal{F}'}} \leq \varphi$$

- ▶ Sufficient condition for **strong duality**

$$\varphi \in \mathcal{F} \implies \varphi^{c_{\mathcal{F}}c_{\mathcal{F}'}} = \varphi$$

- ▶ Two **trivial cases** of strong duality

1.  $\mathcal{F} = \{F: \mathcal{U} \rightarrow \overline{\mathbb{R}}\} = \overline{\mathbb{R}}^{\mathcal{U}}$
2.  $\mathcal{F} = \{\varphi\}$

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## Pure integer linear program in standard form

$$\begin{aligned} & \inf \quad \langle x, k \rangle \\ & x \in \mathbb{Z}^n \\ & Ax = b \\ & x \geq 0 \end{aligned}$$

# Usual continuous dual in PILP

$$\begin{array}{ccc} \text{dual (continuous) problem} & & \text{primal (integer) problem} \\ \underbrace{\sup_{\lambda \in \mathbb{R}^m} \langle b, \lambda \rangle}_{A^T \lambda \leq k} & \leq & \underbrace{\inf_{\substack{x \in \mathbb{Z}^n \\ Ax = b \\ x \geq 0}} \langle x, k \rangle} \end{array}$$

weak duality

- ▶ Right-hand side  $b$  perturbation and scalar product coupling
- ▶ Usually strong duality is not achieved
- ▶ Can we design tighter dual problems?  
(Useful for Branch-and-bound like methods)



# Changing the perturbation/changing the coupling

$$\varphi(b) = \inf_{\substack{x \in \mathbb{Z}_+^n \\ Ax = b}} \langle x, k \rangle$$

$$\langle \cdot, \cdot \rangle : \mathbb{R}^m \times \mathbb{R}^m \rightarrow \mathbb{R}$$

$$\varphi(b) = \dots$$

$$\langle \cdot, \cdot \rangle : \mathbb{R}^m \times \mathbb{R}^m \rightarrow \mathbb{R}$$

$$\varphi(b) = \inf_{\substack{x \in \mathbb{Z}_+^n \\ Ax = b}} \langle x, k \rangle$$

$$c : \mathbb{R}^m \times \mathcal{F} \rightarrow \overline{\mathbb{R}}$$

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# (Geoffrion) Lagrangian relaxation Geoffrion [1974]

1. We partially perturb

$$\forall b^1 \in \mathbb{R}^{m_1}, \varphi(b^1) = \inf_x \langle x, k \rangle$$
$$\begin{aligned} A^1 x &= b^1 \\ A^2 x &= b^2 \\ x &\geq 0 \\ x &\in \mathbb{Z}^n \end{aligned}$$

2. We pair the primal space  $\mathbb{R}^{m_1}$  and a dual space  $\mathbb{R}^{m_1}$

$$\langle, \rangle : \mathbb{R}^{m_1} \times \mathbb{R}^{m_1} \rightarrow \mathbb{R}$$

3. We biconjugate the perturbation function  $\varphi$

$$\varphi^{**'}(b^1) = \sup_{\lambda \in \mathbb{R}^{m_1}} \underbrace{\inf_{\substack{A^2 x = b^2 \\ x \geq 0 \\ x \in \mathbb{Z}^n}} \langle x, k \rangle + \langle b^1 - A^1 x, \lambda \rangle}_{g(\lambda)}$$

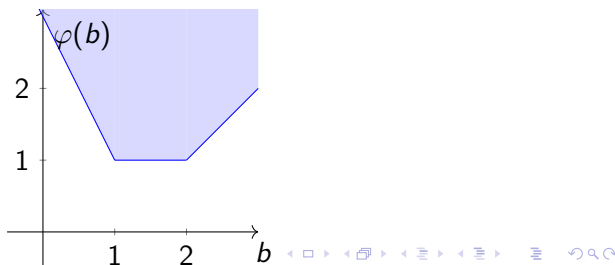
# Example of perturbation function's epigraph for LP

## Example

Let  $\varphi : \mathbb{R} \rightarrow \mathbb{R}$  be defined as

$$\begin{aligned} \varphi(b) = & \inf_{x \in \mathbb{R}^3} && x_1 + x_2 + x_3 \\ & && x_1 + x_2 + 3x_3 = 1 \\ & && x_1 + 2x_2 + 4x_3 = b \\ & && x \geq 0 \end{aligned}$$

Then  $\varphi(b) = \max\{3 - 2b, 1, b - 1\}$ ,  $\forall b \in \mathbb{R}$



# Condition for tighter gap than continuous dual problem

adapted from Conforti, Cornuéjols, and Zambelli [2014, Corollary 8.4.]

## Proposition

If  $A^2$  and  $b^2$  are **rational** then

$$\sup_{\substack{\lambda \in \mathbb{R}^m \\ A^T \lambda \leq k}} \langle b, \lambda \rangle \leq \varphi^{**'}(b^1) = \sup_{\lambda \in \mathbb{R}^{m_1}} g(\lambda)$$

where

$$g(\lambda) = \inf_{\substack{A^2 x = b^2 \\ x \geq 0 \\ x \in \mathbb{Z}^n}} \langle x, k \rangle + \langle b^1 - A^1 x, \lambda \rangle$$

$$\text{and } A = \begin{pmatrix} A^1 \\ A^2 \end{pmatrix}, \quad b = \begin{pmatrix} b^1 \\ b^2 \end{pmatrix}, \quad m = m_1 + m_2$$

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# Evaluation coupling

## Definition

Let  $\mathcal{F} \subset \{F: \mathbb{R}^m \rightarrow \overline{\mathbb{R}}\}$  be a set of functions

We call  $c_{\mathcal{F}}: \mathbb{R}^m \times \mathcal{F} \rightarrow \overline{\mathbb{R}}$  defined by

$$c_{\mathcal{F}}(b, F) = F(b), \quad \forall b \in \mathbb{R}^m, \forall F \in \mathcal{F}$$

the **evaluation coupling** of  $\mathcal{F}$

## Remark

- ▶ Here the **dual variables are functions**
- ▶ If  $\mathcal{F} = \{F: \mathbb{R}^m \rightarrow \overline{\mathbb{R}} \mid \text{is linear}\}$ , then  $c_{\mathcal{F}} = \langle \cdot, \cdot \rangle$

# Resulting evaluation dual problem

also see Tind and Wolsey [1981, Sect. 6]

Consider the perturbation function defined by

$$\forall b \in \mathbb{R}^m, \varphi(b) = \inf_x \langle x, k \rangle$$
$$\begin{array}{l} Ax = b \\ x \geq 0 \\ x \in \mathbb{Z}^n \end{array}$$

## Proposition

Let  $\mathcal{F} \subset \{F: \mathbb{R}^m \rightarrow \overline{\mathbb{R}}\}$  be a set of functions

Then, for any  $b \in \mathbb{R}^m$

$$\varphi^{c_{\mathcal{F}} c_{\mathcal{F}'}}(b) = \sup_{F \in \mathcal{F}} \left\{ F(b) + \inf_{x \in \mathbb{Z}_+^n} \{ \langle x, k \rangle - F(Ax) \} \right\}$$

$$\varphi^{c_{\mathcal{F}} c_{\mathcal{F}'}}(b) \leq \varphi(b)$$

weak duality



## Proof. Compute first conjugate

$$\begin{aligned}\varphi^c(F) &= \sup_{b \in \mathbb{R}^m} \{c(b, F) \dagger (-\varphi(b))\} \\ &= \sup_{b \in \mathbb{R}^m} \{c(b, F) \dagger (-\inf_{\substack{x \in \mathbb{Z}_+^n \\ Ax=b}} \langle x, k \rangle)\} \\ &= \sup_{b \in \mathbb{R}^m} \{c(b, F) \dagger \sup_{\substack{x \in \mathbb{Z}_+^n \\ Ax=b}} -\langle x, k \rangle\} \\ &= \sup_{x \in \mathbb{Z}_+^n} \{-\langle x, k \rangle + \sup_{\substack{b \in \mathbb{R}^m \\ Ax=b}} c(b, F)\} \\ &= \sup_{x \in \mathbb{Z}_+^n} \{-\langle x, k \rangle + c(Ax, F)\} \\ &= -\inf_{x \in \mathbb{Z}_+^n} \{\langle x, k \rangle - F(Ax)\}\end{aligned}$$

# Revisiting the Fenchel coupling with evaluation coupling

- perturbation space  $\mathbb{R}^m$   $\times$  dual space  $\mathbb{R}^m \rightarrow \mathbb{R}$
- ▶ Bilinear coupling  $\langle \cdot, \cdot \rangle : \mathbb{R}^m \times \mathbb{R}^m \rightarrow \mathbb{R}$
  - ▶  $\mathbb{R}^m$  can be identified to the functional space  $\Lambda = \{F: \mathbb{R}^m \rightarrow \mathbb{R} \mid F \text{ is linear}\}$

$$\lambda \in \mathbb{R}^m \leftrightarrow F \in \Lambda$$

$$\langle b, \lambda \rangle \leftrightarrow c_\Lambda(b, F)$$

- ▶ Thus, the resulting dual problem

$$\begin{aligned} \varphi^{c_\Lambda c_{\Lambda'}}(b) &= \sup_{\lambda \in \mathbb{R}^m} \left\{ \langle b, \lambda \rangle + \underbrace{\inf_{x \in \mathbb{Z}_+^n} \{ \langle x, k \rangle - \langle Ax, \lambda \rangle \}}_{\iota_{A^T \lambda \leq k}} \right\} \\ &= \sup_{\substack{\lambda \in \mathbb{R}^m \\ A^T \lambda \leq k}} \langle b, \lambda \rangle \end{aligned}$$

# Inclusion of functional sets

Consider the perturbation function defined by

$$\forall b \in \mathbb{R}^m, \varphi(b) = \inf_x \langle x, k \rangle$$
$$Ax = b$$
$$x \geq 0$$
$$x \in \mathbb{Z}^n$$

## Proposition

If  $\mathcal{F}^1 \subset \dots \subset \mathcal{F}^J \subset \{F: \mathbb{R}^m \rightarrow \bar{\mathbb{R}}\}$ , then

$$\varphi^{c_{\mathcal{F}^1} c_{\mathcal{F}^1}'} \leq \dots \leq \varphi^{c_{\mathcal{F}^J} c_{\mathcal{F}^J}'} \leq \varphi$$

The **larger** the set of dual functions, the **tighter** the gap!

# Characterization of strong duality for evaluation coupling

## Proposition

Tind and Wolsey [1981, Proposition 6.8]

Let  $\mathcal{F} \subset \{F: \mathbb{R}^m \rightarrow \overline{\mathbb{R}}\}$  be a set of functions and  $f: \mathbb{R}^m \rightarrow \overline{\mathbb{R}}$

Then

$$f^{c_{\mathcal{F}} c_{\mathcal{F}'}} = f \iff \exists \{f_i\}_{i \in I} \subset \mathcal{F} \text{ s.t. } f = \sup_{i \in I} f_i$$

## Remark

Cases when the *equality* is trivially true

- ▶  $\mathcal{F}$  is *too "general"*:  $\mathcal{F} = \{F: \mathbb{R}^m \rightarrow \overline{\mathbb{R}}\}$
- ▶  $\mathcal{F}$  is *too "specific"*:  $\mathcal{F} = \{f\}$

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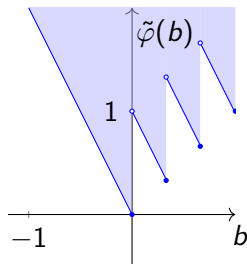
## Example

Let  $\varphi: \mathbb{R} \rightarrow \mathbb{R}$  be defined as

$$\forall b \in \mathbb{R}, \varphi(b) = \inf_{\substack{x \in \mathbb{Z}^2 \\ x_1 - x_2 = b \\ x \geq 0}} x_1 + 2x_2$$

Then  $\varphi$  coincides on its domain with

$$\tilde{\varphi}(b) = \max\{-2b, \lceil 3b \rceil - 2b\}, \forall b \in \mathbb{R}$$



# Summary of the perturbation-duality scheme for PILP

1. We **perturb** a minimization problem

$$\forall b \in \mathbb{R}^m, \varphi(b) = \inf_{\substack{x \\ Ax = b \\ x \in \mathbb{Z}_+^n}} \langle x, k \rangle$$

2. We pair the primal space  $\mathbb{R}^m$  and a function set  $\mathcal{F}$

$$c_{\mathcal{F}}: \mathbb{R}^m \times \mathcal{F} \rightarrow \bar{\mathbb{R}}$$
$$c_{\mathcal{F}}(b, F) = F(b)$$

3. We biconjugate the perturbation function  $\varphi$

weak duality is guaranteed

$$\overbrace{\varphi^{cc'}(b) \leq \varphi(b), \forall b \in \mathbb{R}^m}$$
$$\varphi^{cc'}(b) = \sup_{F \in \mathcal{F}} \left\{ F(b) + \inf_{x \in \mathbb{Z}_+^n} \{ \langle x, k \rangle - F(Ax) \} \right\}$$

# Strong duality for the subadditive dual problem

Let  $\mathcal{S}$  be the set of **subadditive functions**

$$\mathcal{S} = \{F: \mathbb{R}^m \rightarrow \overline{\mathbb{R}} \text{ s.t. } F(b^1 + b^2) \leq F(b^1) \dot{+} F(b^2), \forall b^1, b^2\}$$

$$(-\infty) \dot{+} (+\infty) = (+\infty) \dot{+} (-\infty) = +\infty$$

## Proposition

Let  $A \in \mathbb{R}^{m \times n}$  and  $k \in \mathbb{R}^n$

Let  $c_{\mathcal{S}}: \mathbb{R}^m \times \mathcal{S} \rightarrow \overline{\mathbb{R}}$  be the evaluation coupling of  $\mathcal{S}$

Then, for all  $b \in \mathbb{R}^m$

$$\varphi^{c_{\mathcal{S}} c_{\mathcal{S}'}}(b) = \varphi(b)$$



## Proof of subadditive strong duality. $\varphi$ is subadditive

(i) Let  $b^1, b^2 \in \mathbb{R}^m$

(ii) ▶ Assume there are  $x^1, x^2 \in \mathbb{Z}_+^n$  s.t.

$$Ax^1 = b^1, \quad Ax^2 = b^2$$

Then  $A(x^1 + x^2) = b^1 + b^2$

▶ For all such  $x^1, x^2$  and by definition of  $\varphi$

$$\varphi(b^1 + b^2) \leq \langle k, x^1 + x^2 \rangle = \langle k, x^1 \rangle + \langle k, x^2 \rangle$$

▶ Going to the infimum in  $x^1$ , then in  $x^2$

$$\varphi(b^1 + b^2) \leq \varphi(b^1) + \varphi(b^2)$$

(iii) If there is no such  $x^1, x^2 \in \mathbb{Z}_+^n$

Then  $\varphi(b^1) = +\infty$  or  $\varphi(b^2) = +\infty$

Thus, by definition of  $\dot{+}$

$$\varphi(b^1 + b^2) \leq \varphi(b^1) \dot{+} \varphi(b^2)$$

## Previous work on PILP duality

- ▶ **Surveys** of previous works  
Tind and Wolsey [1981], Güzelsoy, Ralphs, and Cochran [2010]
- ▶ Some pioneer papers on strong duality in **the rational case** (when  $A \in \mathbb{Q}^{m \times n}$ ,  $k \in \mathbb{Q}^n$ )  
Johnson [1973], Jeroslow [1979], Wolsey [1981], Blair and Jeroslow [1982]
- ▶ **Computation** of optimal dual functions  
Wolsey [1981], Klabjan [2007]

### Subadditive dual problem in Jeroslow [1979]

$$\begin{aligned} & \sup && F(\bar{b}) \\ & F: \mathbb{R}^m \rightarrow \mathbb{R} \\ & F(A_j) \leq k_j \\ & F(0) \leq 0 \\ & F \text{ is subadditive} \end{aligned}$$

# Outline of the presentation

PDS for linear programming (LP)

PDS for convex programming

PDS for pure integer linear programming (PILP)

(Geoffrion) Lagrangian relaxation

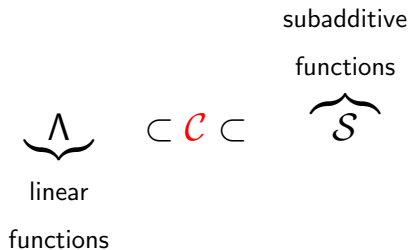
Evaluation coupling in PDS

Subadditive dual problem for PILP

**Chvátal strong duality for rational PILP**

Conclusion

# Somewhere in between the linear and the subadditive



# Definition of Chvátal functions

## Definition

The class of Chvátal functions  $\mathcal{C}$  is the smallest class of functions  $\mathcal{F} \subset \{F: \mathbb{R}^m \rightarrow \mathbb{R}\}$  such that

$$b \in \mathbb{R}^m \mapsto \langle b, \lambda \rangle \in \mathcal{F}, \quad \forall \lambda \in \mathbb{Q}^m \quad (\text{linear functions})$$

$$\alpha F_1 + \beta F_2 \in \mathcal{F}, \quad \forall F_1, F_2 \in \mathcal{F}, \quad \alpha, \beta \in \mathbb{Q}_+ \quad (\text{conic combination})$$

$$\lceil F \rceil \in \mathcal{F}, \quad \forall F \in \mathcal{F} \quad (\text{round-up})$$

## Examples in 1D

- ▶  $b \mapsto \frac{3}{4}b$
- ▶  $b \mapsto \lceil b \rceil$
- ▶  $b \mapsto \frac{3}{4}b + \frac{7}{10}\lceil b \rceil$
- ▶  $b \mapsto 15b + \frac{39}{22}\lceil \frac{3}{4}b + \frac{7}{10}\lceil b \rceil \rceil + \lceil 16b \rceil$

# Strong duality with Chvátal functions

adapted from Blair and Jeroslow [1982]

We define a perturbation function

$$\forall b \in \mathbb{R}^m, \varphi(b) = \inf_{\substack{x \\ Ax = b \\ x \in \mathbb{Z}_+^n}} \langle x, k \rangle$$

## Proposition

We remind that  $c_C$  is the evaluation coupling of the Chvátal functions  
If  $A \in \mathbb{Q}^{m \times n}$  and  $k \in \mathbb{Q}^n$  then

$$\varphi^{c_C c_C'}(b) = \varphi(b), \quad \forall b \in \text{dom} \varphi$$

## Remark

*The perturbation function  $\varphi$  is defined on  $\mathbb{R}^m$  but  $\text{dom} \varphi \subset \mathbb{Q}^m$*

# Outline of the presentation

PDS for linear programming (LP)

PDS for convex programming

PDS for pure integer linear programming (PILP)

Conclusion

# Steps of the perturbation-duality scheme

Rockafellar [1974]

1. **Perturb** a minimization problem with a perturbation (primal) variable belonging to a vector space, and a Rockafellian function
2. **Pair** the (primal) perturbation space with a dual space by means of a bilinear form  $\langle , \rangle$
3. **Biconjugate** the perturbation function, and get
  - ▶ a dual problem
  - ▶ weak duality
4. Deduce conditions for **strong duality** by means of either global or local properties of the perturbation function



## Branching out: rank restricted Chvátal functions

# Rank restricted Chvátal functions

Let  $F \in \mathcal{C}$  be a Chvátal function

## Definition

The **rank** of  $F$  is defined *informally* as the smallest number of  $\lceil \cdot \rceil$  needed to encode  $F$   
We denote by  $\mathcal{C}^r \subset \mathcal{C}$  the Chvátal function of rank **not greater** than  $r \in \mathbb{N}$

- ▶ Inclusion of function sets

$$\underbrace{\mathcal{C}^0}_{\text{linear functions with rational } \lambda} \subset \mathcal{C}^1 \subset \dots \subset \mathcal{C}^m \subset \dots \subset \mathcal{S}$$

- ▶ Weak duality chain

$$\varphi^{\mathcal{C}^0 \mathcal{C}^0'} \leq \varphi^{\mathcal{C}^1 \mathcal{C}^1'} \leq \dots \varphi^{\mathcal{C}^r \mathcal{C}^r'} \leq \dots \leq \varphi^{\mathcal{C}^s \mathcal{C}^s'} \leq \varphi$$

## Example of partially perturbed restricted Chvátal scheme

- ▶ We define a perturbation function

$$\forall b \in \mathbb{R}^{m_1}, \varphi(b^1) = \inf_x \langle x, k \rangle$$
$$\begin{aligned} A^1 x &= b^1 \\ A^2 x &= b^2 \\ x &\geq 0 \\ x &\in \mathbb{Z}^n \end{aligned}$$

- ▶ We define a coupling between primal and dual space

$$c : \mathbb{R}^{m_1} \times \mathbb{Q}^{m_1} \times \mathbb{Q}_+ \rightarrow \mathbb{R}$$
$$c(b^1, (\lambda, \alpha)) = \lambda^T b^1 + \alpha \left[ \beta^T b^1 \right],$$
$$\forall b^1 \in \mathbb{R}^{m_1}, \forall (\lambda, \alpha) \in \mathbb{Q}^{m_1} \times \mathbb{Q}_+$$

- ▶ We biconjugate the perturbation functions

$$\underbrace{\varphi^{cc'}(b^1) \leq \varphi(b^1), \forall b \in \mathbb{R}^{m_1}}_{\text{weak duality}}$$

## Resulting dual problems yields a tighter gap

$$\varphi^{cc'}(b) = \sup_{(\lambda, \alpha) \in \mathbb{Q}^m \times \mathbb{Q}_+} \underbrace{\left\{ \lambda^T b^1 + \alpha \left[ \beta^T b^1 \right] + \inf_{\substack{A^2 x = b^2 \\ x \geq 0 \\ x \in \mathbb{Z}^n}} \left\{ k^T x + \lambda^T A^1 x + \alpha \left[ \beta^T A^1 x \right] \right\} \right\}}_{\tilde{g}(\lambda, \alpha)}$$

We have a tighter gap

$$\underbrace{\sup_{\lambda \in \mathbb{Q}^{m_1}} g(\lambda)}_{\text{Lagrangian relaxation}} \leq \sup_{(\lambda, \alpha) \in \mathbb{Q}^{m_1} \times \mathbb{Q}_+} \tilde{g}(\lambda, \alpha) \leq \underbrace{\varphi(b)}_{\text{original problem}}$$

# Thank you for your attention !

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