Stochastic Dynamic Programming

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- 2 Curses of Dimensionality
- 3 Linear-Quadratic Setting
- Infinite Horizon





2 Curses of Dimensionality

- 3 Linear-Quadratic Setting
- 4 Infinite Horizon

Stochastic Controlled Dynamic System

A stochastic controlled dynamic system is defined by its dynamic

$$\boldsymbol{x}_{t+1} = f_t(\boldsymbol{x}_t, \boldsymbol{u}_t, \boldsymbol{w}_{t+1})$$

and initial state

 $x_0 = x_0$

The variables

- x_t is the state of the system,
- **u**_t is the control applied to the system at time t,
- w_t is an exogeneous noise.

Examples

- Stock of water in a dam:
 - x_t is the amount of water in the dam at time t,
 - **u**_t is the amount of water turbined at time t,
 - **w**_t is the inflow of water at time t.
- Boat in the ocean:
 - x_t is the position of the boat at time t,
 - **u**_t is the direction and speed chosen at time t,
 - w_t is the wind and current at time t.
- Subway network:
 - x_t is the position and speed of each train at time t,
 - **u**_t is the acceleration chosen at time t,
 - **w**_t is the delay due to passengers and incident on the network at time *t*.

Optimization Problem

We want to solve the following optimization problem

$$\min \qquad \mathbb{E} \Big[\sum_{t=0}^{T-1} L_t (\boldsymbol{x}_t, \boldsymbol{u}_t, \boldsymbol{w}_{t+1}) + K (\boldsymbol{x}_T) \Big]$$
(1a)

$$s.t. \quad \boldsymbol{x}_{t+1} = f_t (\boldsymbol{x}_t, \boldsymbol{u}_t, \boldsymbol{w}_{t+1}), \quad \boldsymbol{x}_0 = \boldsymbol{x}_0$$
(1b)

$$\boldsymbol{u}_t \in U_t (\boldsymbol{x}_t)$$
(1c)

$$\sigma(\boldsymbol{u}_t) \subset \sigma(\boldsymbol{w}_0, \cdots, \boldsymbol{w}_t)$$
(1d)

Dynamic Programming Principle

Assume that the noises w_t are independent and exogeneous.

Then, there exists an optimal solution, called a strategy, of the form $\boldsymbol{u}_t = \pi_t(\boldsymbol{x}_t)$, given by

$$\pi_t(x) = \arg\min_{u \in U_t(x)} \mathbb{E}\left[\underbrace{L_t(x, u, \boldsymbol{w}_{t+1})}_{\text{current cost}} + \underbrace{V_{t+1} \circ f_t(x, u, \boldsymbol{w}_{t+1})}_{\text{current cost}}\right],$$

where (Dynamic Programming Equation)

$$\begin{cases} V_{\mathcal{T}}(x) = K(x) \\ V_{t}(x) = \min_{u \in U_{t}(x)} \mathbb{E} \Big[L_{t}(x, u, \boldsymbol{w}_{t+1}) + V_{t+1} \circ \underbrace{f_{t}(x, u, \boldsymbol{w}_{t+1})}_{"\boldsymbol{X}_{t+1}"} \Big] \end{cases}$$

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Interpretation of Bellman Value

The Bellman's value function $V_{t_0}(x)$ can be interpreted as the value of the problem starting at time t_0 from the state x. More precisely we have

$$V_{t_0}(\mathbf{x}) = \min \qquad \mathbb{E}\left[\sum_{t=t_0}^{T-1} L_t(\mathbf{x}_t, \mathbf{u}_t, \mathbf{w}_{t+1}) + K(\mathbf{x}_T)\right] \qquad (2a)$$

s.t. $\mathbf{x}_{t+1} = f_t(\mathbf{x}_t, \mathbf{u}_t, \mathbf{w}_{t+1}), \qquad \mathbf{x}_{t_0} = \mathbf{x} \qquad (2b)$
 $\mathbf{u}_t \in U_t(\mathbf{x}_t) \qquad (2c)$
 $\sigma(\mathbf{u}_t) \subset \sigma(\mathbf{w}_0, \cdots, \mathbf{w}_t) \qquad (2d)$

Information structures in the multistage setting

Open-Loop Every decision $(u_t)_{t \in [\![0, T-1]\!]}$ is taken before any noises $(\xi_t)_{t \in [\![0, T-1]\!]}$ is known. We decide a planning, and stick to it.

Decision Hazard Decision u_t is taken knowing all past noises ξ_0, \ldots, ξ_t , but not knowing ξ_{t+1}, \ldots, ξ_T .

- Hazard Decision Decision \boldsymbol{u}_t is taken knowing all past noises ξ_0, \ldots, ξ_t , and the next noise $\boldsymbol{\xi}_{t+1}$ but not knowing $\boldsymbol{\xi}_{t+2}, \ldots, \boldsymbol{\xi}_T$.
- Anticipative Every decision $(\boldsymbol{u}_t)_{t \in [\![0, \mathcal{T}-1]\!]}$ is taken knowing the whole scenario $(\boldsymbol{\xi}_t)_{t \in [\![0, \mathcal{T}-1]\!]}$. There is one deterministic optimization problem by scenario.

With the same objective function this gives better and better value as the solution use more and more information.

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Information structures: comments

Open-Loop This case can happen in practice (e.g. fixed planning). There are specific methods to solve this type of optimization problem (e.g. stochastic gradient methods).

Decision Hazard The decision u_t is taken at the beginning of period [t, t + 1]. The decision is always implementable, and might be conservative as it doesnot leverage any prediction over the noise in [t, t + 1].

Hazard Decision The decision u_t is taken at the end of period [t, t + 1[. The modelization is optimistic as it assumes perfect knowledge that might not be available in practice.

Anticipative This problem is never realistic. However it is a lower bound of the real problem that can be estimated through Monte-Carlo and deterministic optimization.

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October 8, 2015 9 / 23

Independence of noise

- The Dynamic Programming equation requires only the time-independence of noises.
- This can be relaxed if we consider an extended state.
- Consider a dynamic system driven by an equation

$$\boldsymbol{y}_{t+1} = f_t(\boldsymbol{X}_t, \boldsymbol{u}_t, \boldsymbol{\varepsilon}_{t+1})$$

where the random noise ε_t is an AR1 process :

$$\boldsymbol{\varepsilon}_t = \alpha_t \boldsymbol{\varepsilon}_{t-1} + \beta_t + \boldsymbol{w}_t,$$

 $\{\mathbf{w}_t\}_{t\in\mathbb{Z}}$ being independent.

- Then y_t is called the physical state of the system and DP can be used with the information state $X_t = (y_t, \varepsilon_{t-1})$.
- Generically speaking, if the noise w_t is exogeneous (not affected by decisions u_t), then we can always apply Dynamic Programming with the state

$$(\mathbf{x}_t, \mathbf{w}_1, \ldots, \mathbf{w}_t)$$





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2 Curses of Dimensionality
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- 3 Linear-Quadratic Setting
- 4 Infinite Horizon

Dynamic Programming Algorithm

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Data: Problem parameters
Result: optimal control and value;
V_T \equiv K:
for t: T \rightarrow 0 do
      for x \in \mathbb{X}_t do
            V_t(x) = \infty;
            for u \in U_t(x) do
                 v_{u} = \mathbb{E} \left| L_{t}(x, u, \boldsymbol{w}_{t+1}) + V_{t+1} \circ f_{t}(x, u, \boldsymbol{w}_{t+1}) \right|;
                 if v_{\mu} < v then
             \begin{vmatrix} V_t(x) = v_u ; \\ \pi_t(x) = u ; \end{vmatrix}
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Algorithm 1: Dynamic Programming Algorithm (discrete case) Number of flops: $O(T \times |\mathbb{X}_t| \times |\mathbb{U}_t| \times |\mathbb{W}_t|)$.

3 curses of dimensionality

- State. If we consider 3 independent states each taking 10 values, then $|X_t| = 10^3 = 1000$. In practice DP is not applicable for states of dimension more than 5.
- Decision. The decision are often vector decisions, that is a number of independent decision, hence leading to huge |U_t(x)|.
- Expectation. In practice random information came from large data set. Without a proper statistical treatment computing an expectation is costly. Monte-Carlo approach are costly too, and unprecise.

Numerical considerations

- The DP equation holds in (almost) any case.
- The algorithm shown before compute a look-up table of control for every possible state *offline*. It is impossible to do if the state is (partly) continuous.
- Alternatively, we can focus on computing offline an approximation of the value function V_t and derive the optimal control online by solving a one-step problem, solved only at the current state :

 $\pi_t(x) \in \underset{u \in U_t(x)}{\arg\min} \mathbb{E}\Big[L_t(x, u, \boldsymbol{w}_{t+1}) + V_{t+1} \circ f_t(x, u, \boldsymbol{w}_{t+1})\Big]$

- The field of Approximate DP gives methods for computing those approximate value function (decomposed on a base of functions).
- The simpler one consisting in discretizing the state, and then interpolating the value function.

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DP on a Markov Chain

- Sometimes it is easier to represent a problem as a controlled Markov Chain
- Dynamic Programming equation can be computed directly, without expliciting the control.
- Let's work out an example...





2 Curses of Dimensionality





The Linear-Quadratic setting

We assume a linear dynamic

$$\boldsymbol{x}_{t+1} = A_t \boldsymbol{x}_t + B_t \boldsymbol{u}_t + \boldsymbol{W}_{t+1}$$

associated with a quadratic cost

$$\mathbb{E}\bigg[\sum_{t=0}^{T-1}\left(\boldsymbol{X}_{t}^{\prime}Q_{t}\boldsymbol{x}_{t}+\boldsymbol{u}_{t}^{\prime}R_{t}\boldsymbol{u}_{t}\right)\bigg]+\boldsymbol{X}_{T}^{\prime}Q_{T}\boldsymbol{x}_{T}.$$

- A few more assumptions
 - x_t is of dimension n, u_t of dimension m.
 - Q_t is a symmetric semidefinite positive matrix, and R_t symmetric definite positive.
 - *w_t* is a centered (i.e. of mean 0) independent, exogeneous noise (i.e their law does not depend of the state or control), with finite second order moment.
 - The controls are non-anticipative.

Solving the LQ case

The DP equation read

$$\begin{cases} V_{\mathcal{T}}(x) = x'Q_{\mathcal{T}}x \\ V_{t}(x) = \min_{u} \mathbb{E}\left[x'Q_{t}x + u'R_{t}u + V_{t+1}(A_{t}x + B_{t}u + \mathbf{W}_{t+1})\right] \end{cases}$$

Leading to

$$V_t(x_t) = x'_t K_t x_t + \sum_{\tau=t}^{T-1} \mathbb{E} \big[\boldsymbol{w}'_{t+1} K_{t+1} \boldsymbol{w}_{t+1} \big]$$

and

$$\boldsymbol{u}_t^{\sharp} = \pi_t^{\sharp}(\boldsymbol{x}_t) = L_t \boldsymbol{x}_t \; .$$

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Solving the LQ case

We have

$$V_t(x_t) = x'_t K_t x_t + \sum_{\tau=t}^{T-1} \mathbb{E} \left[\boldsymbol{w}'_{t+1} K_{t+1} \boldsymbol{w}_{t+1} \right]$$

and

$$\boldsymbol{u}_t^{\sharp} = \pi_t^{\sharp}(\boldsymbol{x}_t) = L_t \boldsymbol{x}_t \; .$$

Where

$$L_{t} = -(B'_{t}K_{t+1}B_{t} + R_{t})^{-1}B'_{t}K_{t+1}A_{t},$$

and

$$\begin{cases} K_{T} = Q_{T} \\ K_{t} = A'_{t} \Big(K_{t+1} - K_{t+1} B_{t} \big(B'_{t} K_{t+1} B_{t} + R_{t} \big)^{-1} B'_{t} K_{t+1} \Big) A_{t} + Q_{t} \end{cases}$$





2 Curses of Dimensionality

3 Linear-Quadratic Setting



Introducing the Bellman operators

We define the Bellman operator associated to our optimisation problem

$$T_t(\boldsymbol{J}): \boldsymbol{x} \mapsto \min_{\boldsymbol{u} \in U_t(\boldsymbol{x})} \mathbb{E} \Big[L_t(\boldsymbol{x}, \boldsymbol{u}, \boldsymbol{w}_{t+1}) + \boldsymbol{J} \circ f_t(\boldsymbol{x}, \boldsymbol{u}, \boldsymbol{w}_{t+1}) \Big] .$$

The Dynamic Programming equation can then be written

$$\begin{cases} V_{\mathcal{T}} = K \\ V_t = T_t \Big(V_{t+1} \Big) \end{cases}$$

We also construct the policy-dependent Bellman operator

$$T_t^{\pi}(J): x \mapsto \mathbb{E}\left[L_t(x, \pi(x), \boldsymbol{w}_{t+1}) + J \circ f_t(x, \pi(x), \boldsymbol{w}_{t+1})\right]$$

Discounted fixed cost case

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We now consider the following specific case problem, where $(\mathbf{w}_t)_{t\in\mathbb{N}}$ is i.i.d.

nin
$$\mathbb{E}\left[\sum_{t=0}^{I} \alpha^{t} L(\mathbf{x}_{t}, \mathbf{u}_{t}, \mathbf{w}_{t+1})\right]$$
 (3)

s.t.
$$\mathbf{x}_{t+1} = f(\mathbf{x}_t, \mathbf{u}_t, \mathbf{w}_{t+1}), \quad \mathbf{x}_0 = \mathbf{x}_0$$
 (4)
 $\mathbf{u}_t \in U(\mathbf{x}_t)$ (5)
 $\sigma(\mathbf{u}_t) \subset \sigma(\mathbf{w}_0, \cdots, \mathbf{w}_t)$ (6)

where $\alpha \in]0,1]$. Note that the constraint and cost structure doesnot depend on *t*.

The Bellman operator is given by

$$T(J): x \mapsto \min_{u \in U(x)} \mathbb{E} \Big[L(x, u, \boldsymbol{w}_{t+1}) + \alpha J \circ f(x, u, \boldsymbol{w}_{t+1}) \Big]$$

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Infinite horizon problems

There is different ways of considering the above problem in an "infinite horizon" setting.

- Discounted case. This is the case where α < 1. It is especially easy to treat if the cost *L* is bounded.
- Stochastic shortest path. In this case α = 1 but there is a "cemetary state" such that once reached the system remains there with null cost. Moreover, we assume that the system always reach the cemetary state in a finite time.
- Average cost per stage problems. This approach is mainly taken if the infinite time cost isn't finite (for example $\alpha = 1$ and L > 0). We consider

$$\lim_{T\to\infty}\frac{1}{T} \quad \mathbb{E}\Big[\sum_{t=0}^{T-1}L(\boldsymbol{x}_t,\boldsymbol{u}_t,\boldsymbol{w}_{t+1})\Big]$$

An overview of typical infinite horizon results

Here are the main results that can be shown in infinite horizon problems (under the right set of assumptions)

- the sequence of value function V_{n+1} = T(V_n), converges toward the value function of the infinite horizon problem: lim_{n→∞} V_n = V[‡].
- The optimal value of the infinite horizon problem is a fixed point of the Bellman operator: $V^{\sharp} = T(V^{\sharp})$.
- If π is such that $V^{\sharp} = T^{\pi} V^{\sharp}$ then the stationnary policy π is optimal.

Value iteration algorithm

```
Data: Initial value V^{(0)}

Result: optimal policy and value;

repeat

for x \in \mathbb{X} do

V^{(k+1)}(x) = T(V^{(k)})(x)

until \|V^{(k+1)} - V^{(k)}\|_{\infty} < \varepsilon;
```

Algorithm 2: Value iteration algorithm

- Each step takes $O(|\mathbb{X}| \times |\mathbb{U}| \times |\Omega|)$ flops.
- The error $|V_n(x) V^{\sharp}(x)|$ is bounded by $C\alpha^n$.

Policy iteration algorithm



The policy iteration algorithm terminate in a finite number of step.