

# Two Stage Stochastic Optimization for Fixing Energy Reserves

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## 1 Problem statement

We formalize the problem of fixing energy reserves in a day-ahead market as a two stage stochastic optimization problem. A decision has to be made at night of day  $J$  — which quantity of the cheapest energy production units (reserve) to be mobilized — to meet a demand that will materialize at morning of day  $J+1$ . Excess reserves are penalized; demand unsatisfied by reserves have to be covered by costly extra units. Hence, there is a trade-off to be assessed by optimization.

### 1.1 Stages

There are two stages, represented by the letter  $t$  (for time):

- $t = 0$  corresponds to night of day  $J$ ;
- $t = 1$  corresponds to morning of day  $J + 1$ .

### 1.2 Probabilistic model

We suppose that the demand, materialized on the morning of day  $J + 1$ , can take a finite number  $S$  of possible values  $D_s$ , where  $s$  denotes a *scenario* in the finite set  $\mathbb{S}$  ( $S = \text{card}(\mathbb{S})$ ).

We denote  $\pi_s$  the probability of scenario  $s$ , with

$$\forall s \in \mathbb{S}, \pi_s > 0, \sum_{s \in \mathbb{S}} \pi_s = 1. \quad (1)$$

Notice that we do not consider scenarios with zero probability.

### 1.3 Decision variables

The decision variables are the scalar  $Q_0$  and the finite sequence  $(Q_{1,s})_{s \in \mathbb{S}}$  of scalars, as follows:

- at stage  $t = 0$ , the energy reserve is  $Q_0$ ;
- at stage  $t = 1$ , a scenario  $s$  materializes and the demand  $D_s$  is observed, so that one decides of the recourse quantity  $Q_{1,s}$ .

The decision variables can be considered as indexed by a *tree* with one root (corresponding to the index 0) and as many leafs as scenarios in  $\mathbb{S}$  (each leaf corresponding to the index 1,  $s$ ):  $Q_0$  is attached to the root of the tree, and each  $Q_{1,s}$  is attached to a leaf corresponding to  $s$ .

### 1.4 Optimization problem formulation

The balance equation between supply and demand is

$$Q_0 + Q_{1,s} = D_s, \quad \forall s \in \mathbb{S}. \quad (2)$$

The energies mobilized at stages  $t = 0$  and  $t = 1$  display different features:

- at stage  $t = 0$ , the energy production has maximal capacity  $Q_0^\sharp$ , and producing  $Q_0$  costs  $\mathcal{C}_0(Q_0)$ ;
- at stage  $t = 1$ , the energy production is supposed to be unbounded, and producing  $Q_1$  costs  $\mathcal{C}_1(Q_1)$ .

We consider the stochastic optimization problem

$$\min_{Q_0, Q_{1,s}, s \in \mathbb{S}} \sum_{s \in \mathbb{S}} \pi_s [\mathcal{C}_0(Q_0) + \mathcal{C}_1(Q_{1,s})] \quad (3a)$$

$$\text{s.t.} \quad 0 \leq Q_0 \leq Q_0^\sharp \quad (3b)$$

$$0 \leq Q_{1,s} \quad \forall s \in \mathbb{S} \quad (3c)$$

$$D_s = Q_0 + Q_{1,s} \quad \forall s \in \mathbb{S} \quad (3d)$$

Here, we look for energy reserve  $Q_0$  and recourse energy  $Q_{1,s}$  so that the balance equation (3d) is satisfied (at stage  $t = 1$ ) at minimum *expected cost* in (3a). By weighing each scenario  $s$  with its probability  $\pi_s$ , the optimal solution  $(Q_0^*, (Q_{1,s}^*)_{s \in \mathbb{S}})$  performs a compromise between scenarios.

## 2 Formulation on a tree with linear costs

Here, we suppose that the costs are linear:

$$\mathcal{C}_0(Q_0) = l_0 Q_0, \quad \mathcal{C}_1(Q_1) = l_1 Q_1. \quad (4)$$

Therefore, the stochastic optimization problem (3) now becomes

$$\min_{Q_0, Q_{1,s}, s \in \mathbb{S}} \sum_{s \in \mathbb{S}} \pi_s (l_0 Q_0 + l_1 Q_{1,s}) \quad (5a)$$

$$\text{s.t.} \quad 0 \leq Q_0 \leq Q_0^\# \quad (5b)$$

$$0 \leq Q_{1,s} \quad \forall s \in \mathbb{S} \quad (5c)$$

$$D_s = Q_0 + Q_{1,s} \quad \forall s \in \mathbb{S} \quad (5d)$$

This optimization problem (5) is linear. When the number  $S$  of scenarios is not too large, we can use linear solvers.

**Question 1** We consider the case when  $\mathbb{S} = \{L, M, H\}$  has  $S = 3$  scenarios (low, medium, high). We want to transform the linear optimization problem (5) under a form adapted to a linear solver.

- a) **[1+1+1]** Expand the criterion (5a). Expand the inequalities (5b)–(5c) into an array of five scalar equations (one inequality per equation), and the equalities (5d) into an array of three scalar equations (one equality per equation).
- b) **[1+1+1+1+1]** Let  $x$  be the column vector  $x = (Q_0, Q_{1L}, Q_{1M}, Q_{1H})$ . Propose matrices  $A_e$ ,  $A_i$  and column vectors  $b_e$ ,  $b_i$  and  $c$  such that the linear optimization problem (5) can be written under the form

$$\min_{x \in \mathbb{R}^4} c'x \quad (6a)$$

$$\text{s.t.} \quad A_e x = b_e \quad (6b)$$

$$A_i x \leq b_i \quad (6c)$$

**Question 2** We are going to numerically solve the linear optimization problem (5).

- a) **[2]** Interpret the code below. What is the macro `linpro` doing? What is `lopt`? Copy the code into a file named `tp_q1.sce`.
- b) **[1+2]** Solve a numerical version of problem (5) with  $S = 3$  scenarios and the parameters in the code below, by executing the file `tp_q1.sce`. What is the optimal value  $Q_0^*$  of the reserve? What is the optimal value  $Q_{1,L}^*$ ? Can you explain why these values are optimal? (there is an economic explanation based on relative costs; there is a mathematical explanation based on the properties of the solutions of a linear program).
- c) **[3]** Then, increase and decrease the value of the unitary cost  $l_1$  (especially above and below  $l_0$ ). Show the numerical results that you obtain. What happens to the optimal values  $Q_0^*$ ? Explain why (make the connection with the properties of the solutions of a linear program).

```

// Formulation on a tree with linear costs.
// Numerical resolution by linear programming

// Constant initialization
S=3;// Number of random scenarios
q0m=30;// max capacity for q0

// Demand
if S==3 then
    D=[15;20;50];
    Pr=[0.2;0.6;0.2];// Probabilities of Demand
else
    D=grand(S,1,'uin',5,50);
    Pr=grand(S,1,'unf',0,1);
    Pr=Pr ./sum(Pr);// Probabilities of Demand
end

// Constants used in the cost function
l10=2;l11=5;

// a revoir pour passer a des contraintes egalité et utiliser lb et ub
c=[l10;l11 .*Pr];// cost coefficients

// inequality constraints (bounds on production)
Ai=[-eye(S+1,S+1);eye(1,S+1)];
bi=[zeros(S+1,1);q0m];

// equality constraints, i.e. production equals demand
Ae=[ones(S,1),eye(S,S)];
be=[D];

// solving by linear
// xopt should be [ 15, 0, 5, 35 ] when S=3

// scicoslab version with linpro
A=[Ae;Ai];b=[be;bi];
[xopt,lopt,fopt]=linpro(c,A,b,[],[],size(be,'*'))

```

**Question 3** We are going study the impact of the number  $S$  of scenarios on the numerical resolution of the linear optimization problem (5).

- a) [2] Take  $S = 100$ . What is the optimal value  $Q_0^*$  of the reserve? Identify the scenario  $\bar{s}$  with the lowest demand. What is the optimal value  $Q_{1,\bar{s}}^*$ ? Explain.

b) [1] For what value of  $n$  in  $S = 10^n$  can you no longer solve numerically?

### 3 Formulation on a tree with quadratic convex costs

Here, we suppose that the costs are quadratic and convex:

$$\mathcal{C}_0(Q_0) = \frac{1}{2}K_0Q_0^2 + l_0Q_0, \quad K_0 > 0, \quad \mathcal{C}_1(Q_1) = \frac{1}{2}K_1Q_1^2 + l_1Q_1, \quad K_1 > 0. \quad (7)$$

The optimization problem (3) is quadratic convex:

$$\min_{Q_0, Q_{1,s}, s \in \mathbb{S}} \sum_{s \in \mathbb{S}} \pi_s \left[ \frac{1}{2}K_0Q_0^2 + l_0Q_0 + \frac{1}{2}K_1Q_{1,s}^2 + l_1Q_{1,s} \right] \quad (8a)$$

$$\text{s.t.} \quad 0 \leq Q_0 \leq Q_0^\# \quad (8b)$$

$$0 \leq Q_{1,s} \quad \forall s \in \mathbb{S} \quad (8c)$$

$$D_s = Q_0 + Q_{1,s} \quad \forall s \in \mathbb{S} \quad (8d)$$

When the number  $S$  of scenarios is not too large, we can use quadratic solvers.

**Question 4** We are going to numerically solve the quadratic convex optimization problem (8).

- a) [1] Interpret the code below. What is the macro `quapro` doing? Copy the code into a file named `tp_q2.sce`.
- b) [1] Solve a numerical version with  $S = 3$  scenarios and the parameters in the code below, by executing the files `tp_q1.sce` and `tp_q2.sce`. What is the optimal value  $Q_0^*$  of the reserve? What is the optimal value  $Q_{1,L}^*$ ?
- c) [1+1+2] Check numerically that  $(Q_0^*, Q_{1,L}^*, Q_{1,H}^*, Q_{1,M}^*)$  is an inner solution to the optimization problem (8a), that is, check numerically that the inequalities (8b) and (8c) are strict. What is the difference with the optimal solution of Question 2? Discuss the difference (make the connection with the properties of the solutions of a linear program).
- d) [2+1] Compute the derivatives of the cost functions in (7). Check numerically (giving the details of computation) that the optimal solution  $(Q_0^*, Q_{1,L}^*, Q_{1,M}^*, Q_{1,H}^*)$  satisfies the following relation between marginal costs:

$$\mathcal{C}'_0(Q_0^*) = \pi_L \mathcal{C}'_1(Q_{1,L}^*) + \pi_M \mathcal{C}'_1(Q_{1,M}^*) + \pi_H \mathcal{C}'_1(Q_{1,H}^*). \quad (9)$$

```

// Formulation on a tree with quadratic costs
// Numerical resolution by brute force quadratic programming

// just to get common data
exec('tp_q1.sce');

// the problem is now quadratic; we use a quadratic solver
A=[Ae;Ai];b=[be;bi];
KK0=10;KK1=1;
KK=diag([KK0,KK1*Pr' .*ones(1,S)]);

// scicoslab version with quapro
[xopt,lopt,fopt]=quapro(KK,c,A,b,[],[],size(be,'*'))

```

**Question 5** We are going to numerically solve the quadratic convex optimization problem (8) after changing the relative values of the (quadratic) parameters  $K_0$  and  $K_1$ . For the parameters  $K_0, l_0, l_1$ , we take the same values as those in Question 2 (hence  $l_1 > l_0$ ) and in Question 4.

- a) [1] Take  $K_1 > K_0$  with  $K_1 \approx K_0$ . What are the optimal value  $Q_0^*$  of the reserve and the optimal value  $Q_{1,L}^*$ ?
- b) [1] Take  $K_1 > K_0$  with  $K_1 \gg K_0$ . What are the optimal value  $Q_0^*$  of the reserve and the optimal value  $Q_{1,L}^*$ ?
- c) [2] Discuss.

**Question 6** We are going to study the impact of the number  $S$  of scenarios on the numerical resolution of the quadratic convex optimization problem (8).

- a) [1] Take  $S = 100$ . Solve a numerical version of problem (8) with the parameters  $K_0, l_0, K_1, l_1$  as in the code of Question 4. What is the optimal value  $Q_0^*$  of the reserve? Identify the scenario  $\bar{s}$  with the lowest demand. What is the optimal value  $Q_{1,\bar{s}}^*$ ?
- b) [1] For what value of  $n$  in  $S = 10^n$  can you no longer solve numerically?

**Question 7** This theoretical question may be ignored

(by those who want to focus on numerical results)

For this question, we temporarily ignore the inequalities (8b) and (8c) in (8). Therefore, we consider the optimization problem (8a) with equality constraints (8d), that is:

$$\min_{Q_0, Q_{1,s}, s \in \mathbb{S}} \sum_{s \in \mathbb{S}} \pi_s \left[ \frac{1}{2} K_0 Q_0^2 + l_0 Q_0 + \frac{1}{2} K_1 Q_{1,s}^2 + l_1 Q_{1,s} \right] \quad (10a)$$

$$s.t. \quad D_s = Q_0 + Q_{1,s}, \quad \forall s \in \mathbb{S}. \quad (10b)$$

a) [2] Compute the Hessian matrix of the criterion

$$J_0 : (Q_0, (Q_{1,s})_{s \in \mathbb{S}}) \in \mathbb{R} \times \mathbb{R}^S \mapsto \sum_{s \in \mathbb{S}} \pi_s \left[ \frac{1}{2} K_0 Q_0^2 + l_0 Q_0 + \frac{1}{2} K_1 Q_{1,s}^2 + l_1 Q_{1,s} \right]. \quad (11)$$

What are the dimensions of the Hessian matrix?

b) [3] Why does optimization problem (10) have a solution? (Beware of the domain)

c) [2] Why is the solution unique?

d) [1+2] Why are the equality constraints (10b) qualified? Why does an optimal solution  $(Q_0^*, (Q_{1,s}^*)_{s \in \mathbb{S}})$  of (10) satisfy the Karush-Kuhn-Tucker (KKT) conditions (first-order optimality conditions)?

e) [2] Why is a solution of the KKT conditions an optimal solution of (10)?

f) [1] Write the Lagrangian  $\mathcal{L}_0(Q_0, (Q_{1,s})_{s \in \mathbb{S}}, (\mu_s)_{s \in \mathbb{S}})$  associated with problem (10).

g) [2] Deduce the KKT conditions. Show that there exist  $(\mu_s^*, s \in \mathbb{S})$  such that

$$C'_0(Q_0^*) - \sum_{s \in \mathbb{S}} \mu_s^* = 0 \quad \text{and} \quad \pi_s C'_1(Q_{1,s}^*) - \mu_s^* = 0, \quad \forall s \in \mathbb{S}. \quad (12)$$

h) [2] Deduce that — when  $(Q_0^*, (Q_{1,s}^*)_{s \in \mathbb{S}})$  is an inner optimal solution to problem (8a) — we have the following relation between marginal costs:

$$C'_0(Q_0^*) = \sum_{s \in \mathbb{S}} \pi_s C'_1(Q_{1,s}^*). \quad (13)$$

Give an economic interpretation of this equality.

## 4 Formulation on a fan with quadratic convex costs

When the number  $S$  of scenarios is too large, Problem (5) — be it linear or convex — cannot be solved by direct methods.

### 4.1 Dualization of non-anticipativity constraints

To bypass this problem, we use a “trick” consisting in introducing new decision variables  $(Q_{0,s})_{s \in \mathbb{S}}$ , instead of the single decision variable  $Q_0$ , and we write

$$Q_{0,s} = \sum_{s' \in \mathbb{S}} \pi_{s'} Q_{0,s'}, \quad \forall s \in \mathbb{S}. \quad (14)$$

These equalities are called the *non-anticipativity* constraints. Indeed, the equations (14) express that

$$Q_{0,s} = Q_{0,s'} , \quad \forall (s, s') \in \mathbb{S}^2 , \quad (15)$$

that is, the decision at stage  $t = 0$  does not depend on the scenario  $s$ , hence cannot anticipate the future. Later, we will treat the constraints (14) by duality.

Therefore, the stochastic optimization problem (3) now becomes

$$\min_{Q_{0,s}, Q_{1,s}, \forall s \in \mathbb{S}} \sum_{s \in \mathbb{S}} \pi_s [\mathcal{C}_0(Q_{0,s}) + \mathcal{C}_1(Q_{1,s})] \quad (16a)$$

$$\text{s.t} \quad 0 \leq Q_{0,s} \leq Q_0^\# \quad \forall s \in \mathbb{S} \quad (16b)$$

$$0 \leq Q_{1,s} \quad \forall s \in \mathbb{S} \quad (16c)$$

$$D_{1,s} = Q_{0,s} + Q_{1,s} \quad \forall s \in \mathbb{S} \quad (16d)$$

$$Q_{0,s} = \sum_{s' \in \mathbb{S}} \pi_{s'} Q_{0,s'} \quad \forall s \in \mathbb{S}$$

By the assumption (1) that there are no scenarios with zero probability ( $\pi_s > 0$ ), we replace each equality in the last equation by the equivalent one

$$\pi_s Q_{0,s} = \pi_s \sum_{s' \in \mathbb{S}} \pi_{s'} Q_{0,s'} , \quad \forall s \in \mathbb{S} . \quad (16e)$$

We attach, to each equality above a *multiplier*  $\lambda_{0,s}$ . We put

$$Q = ((Q_{0,s})_{s \in \mathbb{S}}, (Q_{1,s})_{s \in \mathbb{S}}) , \quad \lambda = (\lambda_{0,s})_{s \in \mathbb{S}} . \quad (17)$$

The corresponding *Lagrangian* is

$$\mathcal{L}(Q, \lambda) = \sum_{s \in \mathbb{S}} \pi_s \left[ \mathcal{C}_0(Q_{0,s}) + \mathcal{C}_1(Q_{1,s}) + \lambda_{0,s} (Q_{0,s} - \sum_{s' \in \mathbb{S}} \pi_{s'} Q_{0,s'}) \right] \quad (18a)$$

$$= \sum_{s \in \mathbb{S}} \pi_s \left[ \mathcal{C}_0(Q_{0,s}) + (\lambda_{0,s} - \sum_{s' \in \mathbb{S}} \pi_{s'} \lambda_{0,s'}) Q_{0,s} + \mathcal{C}_1(Q_{1,s}) \right] . \quad (18b)$$

**Question 8** This theoretical question may be ignored

(by those who want to focus on numerical results)

When the costs are quadratic and convex as in (7), show that

a) [3] the criterion

$$J : Q \in \mathbb{R}^S \times \mathbb{R}^S \mapsto \sum_{s \in \mathbb{S}} \pi_s \left[ \frac{1}{2} K_0 Q_{0,s}^2 + l_0 Q_{0,s} + \frac{1}{2} K_1 Q_{1,s}^2 + l_1 Q_{1,s} \right] \quad (19)$$

in (16a) is  $a$ -strongly convex in  $((Q_{0,s})_{s \in \mathbb{S}}, (Q_{1,s})_{s \in \mathbb{S}})$ , and provide a possible  $a > 0$ ,



- b) [2] the domain defined by the constraints in the optimization problem (16) is closed,
- c) [2] the domain defined by the constraints in the optimization problem (16) is convex,
- d) [2+1] the optimization problem (16) has a solution, and it is unique, denoted by  $Q^*$ ,
- e) [1] there exists a multiplier  $\lambda^*$  such that  $(Q^*, \lambda^*)$  is a saddle point of the Lagrangian  $\mathcal{L}$  in (18).

## 4.2 Uzawa algorithm

We consider the following optimization problem

$$\min_{u \in \mathbb{U}^{\text{ad}}} J(u) \tag{20a}$$

$$s.t. \quad \Theta(u) = 0, \tag{20b}$$

under the assumptions that

- the set  $\mathbb{U}^{\text{ad}}$  is a closed convex subset of a Euclidian space  $\mathbb{R}^N$ ,
- the criterion  $J : \mathbb{R}^N \rightarrow \mathbb{R}$  is an  $a$ -strongly convex ( $a > 0$ ) and differentiable function,
- the constraint mapping  $\theta : \mathbb{R}^N \rightarrow \mathbb{R}^M$  is affine,  $\kappa$ -Lipschitz ( $\kappa > 0$ ),
- the Lagrangian  $\mathcal{L}(u, \lambda) = J(u) + \langle \lambda, \Theta(u) \rangle$  admits a saddle-point over  $\mathbb{U}^{\text{ad}} \times \mathbb{R}^M$ .

Then the following algorithm — called *dual gradient* algorithm, or *Uzawa algorithm* — converges toward the optimal solution of Problem (20), when  $0 < \rho < 2a/\kappa^2$ .

**Data:** Initial multiplier  $\lambda^{(0)}$ , step  $\rho$

**Result:** optimal decision and multiplier;

**repeat**

$u^{(k)} = \arg \min_{u \in \mathbb{U}^{\text{ad}}} \mathcal{L}(u, \lambda^{(k)})$  ;

$\lambda^{(k+1)} = \lambda^{(k)} + \rho \Theta(u^{(k)})$  ;

**until**  $\Theta(u^{(k)}) = 0$ ;

**Algorithm .1:** Dual Gradient Algorithm

### 4.3 Numerical resolution by Uzawa algorithm (quadratic convex case)

When the costs are quadratic and convex as in (7), the optimization problem (16) becomes:

$$\min_{Q_{0,s}, Q_{1,s}, \forall s \in \mathbb{S}} \sum_{s \in \mathbb{S}} \pi_s \left[ \frac{1}{2} K_0 Q_{0,s}^2 + l_0 Q_{0,s} + \frac{1}{2} K_1 Q_{1,s}^2 + l_1 Q_{1,s} \right] \quad (21a)$$

$$\text{s.t.} \quad 0 \leq Q_{0,s} \leq Q_0^\# \quad \forall s \in \mathbb{S} \quad (21b)$$

$$0 \leq Q_{1,s} \quad \forall s \in \mathbb{S} \quad (21c)$$

$$D_{1,s} = Q_{0,s} + Q_{1,s} \quad \forall s \in \mathbb{S} \quad (21d)$$

$$Q_{0,s} = \sum_{s' \in \mathbb{S}} \pi_{s'} Q_{0,s'} \quad \forall s \in \mathbb{S} \quad (21e)$$

**Question 9** This theoretical question may be ignored

(by those who want to focus on numerical results)

When the costs are quadratic and convex as in (7), identify in the optimization problem (21) the corresponding elements in the Uzawa algorithm .1:

- a) [1] decision variable  $u$ ,
- b) [1] decision set  $\mathbb{R}^N$  (for what  $N$  ?),
- c) [2] affine constraints mapping  $\Theta : \mathbb{R}^N \rightarrow \mathbb{R}^M$ , corresponding to the constraints (21e) (why is it  $\kappa$ -Lipschitz, and for which  $\kappa$  ?).
- d) [3] Explain why the Uzawa algorithm converges towards the optimal solution of Problem (21).

**Question 10** We are going to numerically solve the quadratic convex optimization problem (21) by Uzawa algorithm.

- a) [3] Detail what the code below is doing ; explain how the code implements the Uzawa algorithm. Why do we use the macro `quapro`? What is the dual function? Copy the code into a file named `tp_q3.sce`.
- b) [2] Solve a numerical version with  $S = 3$  scenarios. What is the solution  $(Q_0^*, (Q_{1,s}^*)_{s \in \mathbb{S}})$  given by the algorithm? Do you have that  $D_{1,s} = Q_0^* + Q_{1,s}^*$ , for all  $s \in \mathbb{S}$  ?
- c) [1] Then, try with  $S = 100$ . What is the value  $Q_0^*$  of the reserve given by the algorithm? Identify the scenario  $\bar{s}$  with the lowest demand. What is the value  $Q_{1,\bar{s}}^*$  given by the algorithm?
- d) [1] For what value of  $n$  in  $S = 10^n$  can you no longer solve numerically?

```

// formulation on a fan
// Constant initialization

// Constant initialization
S=3;// Number of random scenarios
q0m=30;// max capacity for q0

// Demand
if S==3 then
    D=[15;20;50];
    Pr=[0.2;0.6;0.2];// Probabilities of Demand
else
    D=grand(S,1,'uin',5,50);
    Pr=grand(S,1,'unf',0,1);
    Pr=Pr ./sum(Pr);// Probabilities of Demand
end

// Constants used in the cost function
l10=2;l11=5;
KK0=10;KK1=1;

// Uzawa iterations when the dualized constraints are the S constraints
// Q0(ss) = sum(Pr.*Q0);

Q0=zeros(S,1);
Q1=zeros(S,1);
f0=zeros(S,1);

rho=5;

lambda=zeros(S,1);
// iterations of the Uzawa algorithm
for it=0:30 do
    // decomposed minimizations (loop over scenarios ss)
    for ss=1:S do
        // inequality constraints (bounds)
        // bounds on (Q0s,Q1s)
        l1=[l10*Pr(ss);l11*Pr(ss)];// cost coefficients
        Ai=[-eye(2,2);eye(1,2)];
        bi=[zeros(2,1);q0m];
        // equality constraints i.e production equals demand
        Ae=[1,1];
    end
end

```

```

be=[D(ss)];
A=[Ae;Ai];b=[be;bi];
KK=Pr(ss)*diag([KK0,KK1]);
//
cc=11+Pr(ss)*[lambda(ss)-sum(Pr .*lambda);0];
[xopt,lopt,fopt]=quapro(KK,cc,A,b,[],[],size(be,'*'));
Q0(ss)=xopt(1);
Q1(ss)=xopt(2);
f0(ss)=fopt;
printf("Dual function %f\n",sum(f0));
end
Q0bar=sum(Pr .*Q0);
lambda=lambda+rho*(Pr .*(Q0-Q0bar));
end

```

## 5 Formulation on a fan with linear costs

Here, we suppose that the costs are linear, as in (4).

### 5.1 Difficulties in applying Uzawa algorithm with linear costs

The optimization problem (16) becomes:

$$\min_{Q_{0,s}, Q_{1,s}, \forall s \in \mathbb{S}} \sum_{s \in \mathbb{S}} \pi_s [l_0 Q_{0,s} + l_1 Q_{1,s}] \quad (22a)$$

$$\text{s.t.} \quad 0 \leq Q_{0,s} \leq Q_0^\# \quad \forall s \in \mathbb{S} \quad (22b)$$

$$0 \leq Q_{1,s} \quad \forall s \in \mathbb{S} \quad (22c)$$

$$D_{1,s} = Q_{0,s} + Q_{1,s} \quad \forall s \in \mathbb{S} \quad (22d)$$

$$Q_{0,s} = \sum_{s' \in \mathbb{S}} \pi_{s'} Q_{0,s'} \quad \forall s \in \mathbb{S} \quad (22e)$$

**Question 11** We are going to numerically solve the linear optimization problem (22).

- a) [2] Detail what the code below is doing. Why do we use the macro `linpro`? Copy the code into a file named `tp-q4.sce`.
- b) [3] Solve a numerical version with  $S = 3$  scenarios. What do you observe regarding convergence of the Uzawa algorithm? Can you explain why?

```

// formulation on a fan with linear cost
// Uzawa does not work

```

```

S=3; // Number of random scenarios
q0m=30; // max capacity for q0
D=[15;20;50];
Pr=[0.2;0.6;0.2]; // Probabilities of Demand
rho=0.1;
// Constants used in the cost function
l10=2;l11=5;

// Uzawa iterations when the dualized constraints are the S constraints
// Q0(ss) = sum(Pr.*Q0);

Q0=zeros(S,1);
Q1=zeros(S,1);
f0=zeros(S,1);

lambda=zeros(S,1);
for it=0:30 do
    // decomposed minimizations
    for ss=1:S do
        // inequality constraints (bounds)
        // bounds on (Q0s,Q1s)
        l1=[l10*Pr(ss);l11*Pr(ss)]; // cost coefficients
        Ai=[-eye(2,2);eye(1,2)];
        bi=[zeros(2,1);q0m];
        // equality constraints i.e production equals demand
        Ae=[1,1];
        be=[D(ss)];
        A=[Ae;Ai];b=[be;bi];
        //
        cc=l1+Pr(ss)*[lambda(ss)-sum(Pr .*lambda);0];
        [xopt,lopt,fopt]=linpro(cc,A,b,[],[],size(be,'*'));
        Q0(ss)=xopt(1);
        Q1(ss)=xopt(2);
        f0(ss)=fopt;
        printf("Dual function %f\n",sum(f0));
    end
    Q0bar=sum(Pr .*Q0);
    lambda=lambda+rho*(Pr.*(Q0-Q0bar));
end

```

## 5.2 Dualization of non-anticipativity constraints

To bypass the difficulty in applying Uzawa algorithm with linear costs, we use a “trick” consisting in introducing new decision variables  $\overline{Q_0}$  and  $(Q_{0,s})_{s \in \mathbb{S}}$ , instead of the single decision variable  $Q_0$ , and we write

$$Q_{0,s} = \overline{Q_0}, \quad \forall s \in \mathbb{S}. \quad (23)$$

These equalities are another form of the *non-anticipativity* constraints. Indeed, the equations (23) express that the decision at stage  $t = 0$  cannot anticipate the future, hence cannot depend on the scenario  $s$ . We will treat these constraints by duality.

Therefore, the stochastic optimization problem (3) now becomes

$$\min_{\overline{Q_0}, Q_{0,s}, Q_{1,s}, \forall s \in \mathbb{S}} \sum_{s \in \mathbb{S}} \pi_s [l_0 Q_{0,s} + l_1 Q_{1,s}] \quad (24a)$$

$$\text{s.t.} \quad 0 \leq Q_{0,s} \leq Q_0^\# \quad \forall s \in \mathbb{S} \quad (24b)$$

$$0 \leq Q_{1,s} \quad \forall s \in \mathbb{S} \quad (24c)$$

$$D_{1,s} = Q_{0,s} + Q_{1,s} \quad \forall s \in \mathbb{S} \quad (24d)$$

$$Q_{0,s} = \overline{Q_0} \quad \forall s \in \mathbb{S} \quad (24e)$$

By the assumption (1) that there are no scenarios with zero probability ( $\pi_s > 0$ ), we replace each equality in (24e) by the equivalent one

$$\pi_s Q_{0,s} = \pi_s \overline{Q_0}, \quad \forall s \in \mathbb{S}. \quad (25)$$

We attach, to each equality above a *multiplier*  $\lambda_{0,s}$ . We put

$$Q_0 = (Q_{0,s})_{s \in \mathbb{S}}, \quad Q_1 = (Q_{1,s})_{s \in \mathbb{S}}, \quad \lambda = (\lambda_{0,s})_{s \in \mathbb{S}}. \quad (26)$$

## 5.3 Augmented Lagrangian and obstacles to decomposition

### 5.4 Progressive Hedging algorithm (quadratic solver)

**Question 12** *We are going to numerically solve the linear optimization problem (24) by the Progressive Hedging algorithm.*

- a) [3] *Detail what the code below is doing. Why do we use the macro `quapro`? Explain the two roles of the new parameter `rr`. Copy the code into a file named `tp_q5.sce`.*
- b) [2] *Solve a numerical version with  $S = 3$  scenarios. What do you observe regarding convergence of the Uzawa algorithm? Can you explain why?*

```
// formulation on a fan
// with linear cost and augmented lagrangian
```

```

S=3;// Number of random scenarios
q0m=30;// max capacity for q0
D=[15;20;50];
Pr=[0.2;0.6;0.2];// Probabilities of Demand
// Constants used in the cost function
l10=2;l11=5;
// Constant used both as a quadratic term and as a gradient step
rr=0.1;

// Uzawa iterations when the dualized constraints are the S constraints
// Q0s = sum(Pr.*Q0)

Q0=zeros(S,1);
Q1=zeros(S,1);
f0=zeros(S,1);

lambda=zeros(S,1);
Q0bar=0;
for it=0:300 do
    // decomposed minimizations
    // we alternate minimization
    for ss=1:S do
        // inequality constraints (bounds)
        // bounds on (Q0s,Q1s)
        l1=[l10*Pr(ss);l11*Pr(ss)];// cost coefficients
        Ai=[-eye(2,2);eye(1,2)];
        bi=[zeros(2,1);q0m];
        // equality constraints i.e production equals demand
        Ae=[1,1];
        be=[D(ss)];
        A=[Ae;Ai];b=[be;bi];
        KK=rr*Pr(ss)*diag([1,0]);
        cc=l1+[Pr(ss)*lambda(ss);0];
        cc=cc+[-rr*Q0bar*Pr(ss)];
        [xopt,lopt,fopt]=quapro(KK,cc,A,b,[],[],size(be,'*'));
        Q0(ss)=xopt(1);
        Q1(ss)=xopt(2);
        f0(ss)=fopt;
    end
    // updates of Q0bar
    Q0bar=sum(Pr .*Q0);
    lambda=lambda+rr*(Pr .*(Q0-Q0bar));
end

```

```
end
```

```
// solution (Q0bar,Q1)
Q1=max(0,(D-Q0bar));
// to be compared with tp_q1
```

## 5.5 Progressive Hedging algorithm (linear solver)

### A Additional code for “Formulation on a tree with linear costs”

```
// Formulation on a tree with linear costs.
// Numerical resolution by linear programming
// using nsp linprog (glpk).

// Constant initialization
S=3; // Number of random scenarios
q0m=30; // max capacity for q0

// Demand
if S==3 then
    D=[15;20;50];
    Pr=[0.2;0.6;0.2]; // Probabilities of Demand
else
    D=grand(S,1,'uin',5,50);
    Pr=grand(S,1,'unf',0,1);
    Pr=Pr ./sum(Pr); // Probabilities of Demand
end

// Constants used in the cost function
l10=2;l11=5;

// a revoir pour passer a des contraintes egalite et utiliser lb et ub
c=[l10;l11 .*Pr]; // cost coefficients

// inequality constraints (bounds on production)
Ai=[-eye(S+1,S+1);eye(1,S+1)];
bi=[zeros(S+1,1);q0m];

// equality constraints, i.e. production equals demand
Ae=[ones(S,1),eye(S,S)];
be=[D];
```



```

// solving by linear
// xopt should be [ 15, 0, 5, 35 ] when S=3

if exists('%nsp') then
    // in nsp linprog is glpk
    [xopt,fopt,flag]=linprog(c,Ai,bi,Ae,be);
    // we can also use quapro if available with a 0 quadratic term
    if exists('quapro','callable') then
        A=[Ae;Ai];b=[be;bi];
        [xopt_q,lagr_q,fopt_q]=quapro(zeros(S+1,S+1),c,A,b,[],[],size(be,'*'))
    end
else
    // scicoslab version with linpro
    A=[Ae;Ai];b=[be;bi];
    [xopt,lopt,fopt]=linpro(c,A,b,[],[],size(be,'*'))
end

```

## B Additional code for “Formulation on a tree with quadratic convex costs”

```

// Formulation on a tree with quadratic costs
// Numerical resolution by brute force quadratic programming
// using nsp quapro or optim

if exists('%nsp') then
    load_toolbox('quapro');
end

// just to get common data
exec('tp_q1.sce');

// the problem is now quadratic; we use a quadratic solver if available
A=[Ae;Ai];b=[be;bi];
KK0=10;KK1=1;
KK=diag([KK0,KK1*Pr' .*ones(1,S)]);

if exists('%nsp') then
    // we can also use quapro if available with a 0 quadratic term
    if exists('quapro','callable') then
        [xopt_q,lagr_q,fopt_q]=quapro(KK,c,A,b,[],[],size(be,'*'))
    end
end

```

```

    end
else
    // scicoslab version with quapro
    [xopt,lopt,fopt]=quapro(KK,c,A,b,[],[],size(be,'*'))
end

// BEWARE: using a global optim if quapro is not available
// we eliminate the equality constraint to use optim on q0

function [f,g,ind]=costf(q0,ind)
    q1=max(0,(D-q0));
    // expression of the cost
    f=sum(Pr.*(l10*q0+l11*q1))+(1/2)*[q0;q1]'*KK*[q0;q1]
    // expression of the cost derivative
    g=sum(Pr.*(l10-l11*(D-q0 >= 0)))+KK(1,1)*q0+-sum((D-q0 >= 0).*(KK(2:$,2:$)*q1))
endfunction

if exists('%nsp') then
    Q0=optim(costf,0,xinf = 0,xsup = q0m)
else
    [fo,Q0]=optim(costf,'b',0,q0m,0)
end

// solution (Q0,Q1)
Q1=max(0,(D-Q0))

```