

Stochastic splitting methods for solving large-scale convex optimization problems

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Context

Fast development of optimization methods over the last decade

Why?

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Fast development of optimization methods over the last decade

Why?

- Interest in nonsmooth cost functions (*sparsity*)
- Need for optimal processing of massive datasets (*big data*)
 - ~~ large number of variables
(inverse problems, energy management)
 - ~~ large number of observations (machine learning)
- Use of more sophisticated data structures
(*graph signal processing*)

Proximity operator

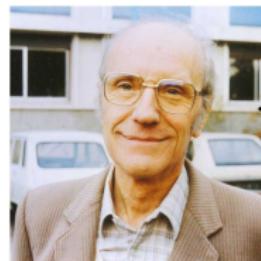
Let H be a real Hilbert space.

$\Gamma_0(H)$ is the class of lower-semicontinuous convex functions from H to $]-\infty, +\infty]$ with a nonempty domain.

Let $f \in \Gamma_0(H)$. Let $U : H \rightarrow H$ be a strongly positive self-adjoint linear operator.

The proximity operator $\text{prox}_f^U(x)$ of f at $x \in H$ relative to the metric induced by U is the unique vector $\hat{y} \in H$ such that

$$f(\hat{y}) + \frac{1}{2}\|\hat{y} - x\|_U^2 = \inf_{y \in H} f(y) + \frac{1}{2} \langle y - x \mid U(y - x) \rangle.$$



Jean-Jacques Moreau
(1923–2014)

Proximity operator

Let $f \in \Gamma_0(\mathbf{H})$. Let $\mathbf{U} : \mathbf{H} \rightarrow \mathbf{H}$ be a strongly positive self-adjoint linear operator.

The proximity operator $\text{prox}_f^{\mathbf{U}}(x)$ of f at $x \in \mathbf{H}$ relative to the metric induced by \mathbf{U} is the unique vector $\hat{y} \in \mathbf{H}$ such that

$$f(\hat{y}) + \frac{1}{2} \|\hat{y} - x\|_{\mathbf{U}}^2 = \inf_{y \in \mathbf{H}} f(y) + \frac{1}{2} \langle y - x \mid \mathbf{U}(y - x) \rangle.$$

Remark:

- If C is a nonempty closed convex subset of \mathbf{H} . Let

$$(\forall x \in \mathbf{H}) \quad \iota_C(x) = \begin{cases} 0 & \text{if } x \in C \\ +\infty & \text{otherwise.} \end{cases}$$

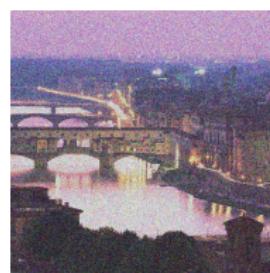
Then, $\text{prox}_{\iota_C}^{\text{Id}}$ reduces to the projection operator onto C .

Proximity operator: Bayesian interpretation

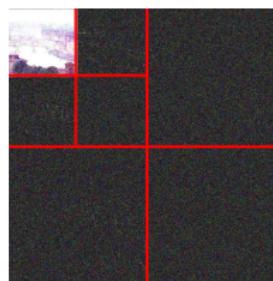
- If $H = \mathbb{R}^N$ and

$$x = \bar{y} + w$$

where \bar{y} is a realization of a random vector with probability density function $\exp(-f)$ and w is a realization of a $\mathcal{N}(0, U^{-1})$ noise, then $\text{prox}_f^U(x)$ is a Maximum A Posteriori estimate of \bar{y} .



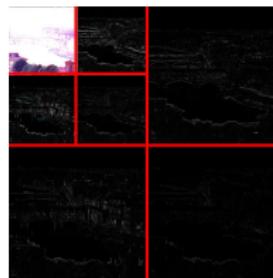
WT
→



←
 prox_f^U



IWT
←



Examples of proximity operators

- Explicit form for objective functions associated to the usual log-concave probability densities when $U = \text{Id}$. [Chaux *et al.* - 2007]
 - Laplace
 - Generalized Gaussian
 - maximum entropy
 - gamma
 - uniform
 - Weibull
 - Generalized inverse Gaussian
 - Gaussian
 - Huber
 - Smoothed Laplace
 - chi
 - triangular
 - Pearson type I
 - ...
- And many other functions ! [Combettes and Pesquet - 2011]

Minimization problems

- Minimization of one function

Let H be a Hilbert space. Let $f \in \Gamma_0(H)$.
We want to

$$\underset{x \in H}{\text{minimize}} \quad f(x)$$

Minimization problems

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- Standard first-order algorithmic solutions
 - ▶ gradient algorithm

$$(\forall n \in \mathbb{N}) \quad x_{n+1} = x_n - \gamma_n \nabla f(x_n)$$

$\rightsquigarrow f$ β - Lipschitz differentiable with $\beta \in]0, +\infty[$
 $0 < \inf_{n \in \mathbb{N}} \gamma_n$ and $\sup_{n \in \mathbb{N}} \gamma_n < 2\beta^{-1}$

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- ▶ proximal point algorithm

$$(\forall n \in \mathbb{N}) \quad x_{n+1} = \text{prox}_{\gamma_n f}^{\text{Id}} x_n$$

$(\gamma_n)_{n \in \mathbb{N}}$ sequence in $]0, +\infty[$ such that $\sum_{n=1}^{+\infty} \gamma_n = +\infty$

$\rightsquigarrow f$ proximable

Minimization problems

- Splitting in a sum of 2 functions

Let H be a Hilbert space. Let $g_1 \in \Gamma_0(H)$ and $g_2 \in \Gamma_0(H)$.

We want to

$$\underset{x \in H}{\text{minimize}} \quad g_1(x) + g_2(x)$$

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- ▶ forward-backward algorithm

$$(\forall n \in \mathbb{N}) \quad x_{n+1} = \text{prox}_{\gamma_n g_1}^{\text{Id}}(x_n - \gamma_n \nabla g_2(x_n))$$

g_2 β -Lipschitz differentiable with $\beta \in]0, +\infty[$

$0 < \inf_{n \in \mathbb{N}} \gamma_n$ and $\sup_{n \in \mathbb{N}} \gamma_n < 2\beta^{-1}$

Particular case: projected gradient algorithm when
 $g_1 = \iota_C$ where C nonempty closed convex subset of H

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$$(\forall n \in \mathbb{N}) \quad x_{n+1} = x_n + \lambda_n (\text{prox}_{\gamma_n g_1}^{\mathbf{U}^{-1}}(x_n - \gamma_n \mathbf{U} \nabla g_2(x_n)) - x_n)$$

g_2 β -Lipschitz differentiable with $\beta \in]0, +\infty[$

$0 < \inf_{n \in \mathbb{N}} \gamma_n$ and $\sup_{n \in \mathbb{N}} \gamma_n < 2\beta^{-1}$

$\mathbf{U}: H \rightarrow H$ strongly positive self-adjoint linear operator

$(\lambda_n)_{n \in \mathbb{N}}$ sequence in $]0, 1]$ such that $\inf_{n \in \mathbb{N}} \lambda_n > 0$

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- ▶ Douglas-Rachford algorithm

$$(\forall n \in \mathbb{N}) \quad x_{n+1} = x_n + \frac{\lambda_n}{2} \left((2\text{prox}_{\gamma g_1}^{\text{Id}} - \text{Id}) \circ (2\text{prox}_{\gamma g_2}^{\text{Id}} - \text{Id}) x_n - x_n \right)$$

$\gamma \in]0, +\infty[$, $(\lambda_n)_{n \in \mathbb{N}}$ sequence in $[0, 2]$ such that

$$\sum_{n \in \mathbb{N}} \lambda_n (2 - \lambda_n) = +\infty$$

Particular case: ADMM amounts to applying DR to the dual problem.

Minimization problems

- Splitting in a sum of $q \geq 2$ functions

Let H be a Hilbert space.

For every $k \in \{1, \dots, q\}$, let $g_k \in \Gamma_0(H)$. We want to

$$\underset{x \in H}{\text{minimize}} \quad \sum_{k=1}^q g_k(x)$$

Minimization problems

- Splitting in a sum of $q \geq 2$ functions

Let H be a Hilbert space.

For every $k \in \{1, \dots, q\}$, let G_k be a Hilbert space, let $g_k \in \Gamma_0(G_k)$, let $L_k : H \rightarrow G_k$ be a linear and bounded operator. We want to

$$\underset{x \in H}{\text{minimize}} \quad \sum_{k=1}^q g_k(L_k x)$$

Minimization problems

- Splitting in a sum of $q \geq 2$ functions

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- Algorithmic solutions

- ▶ Simultaneous direction method of multipliers-like algorithms (PPXA, PPXA+,...)
- ▶ Proximal primal-dual algorithms

Proximal primal-dual algorithm

Advantages:

- No linear operator inversion.
- Use of proximable or/and differentiable functions.
- Amenable to parallel implementations.
- Special cases: Forward-Backward and Douglas-Rachford algorithms.

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Bibliographical remarks:

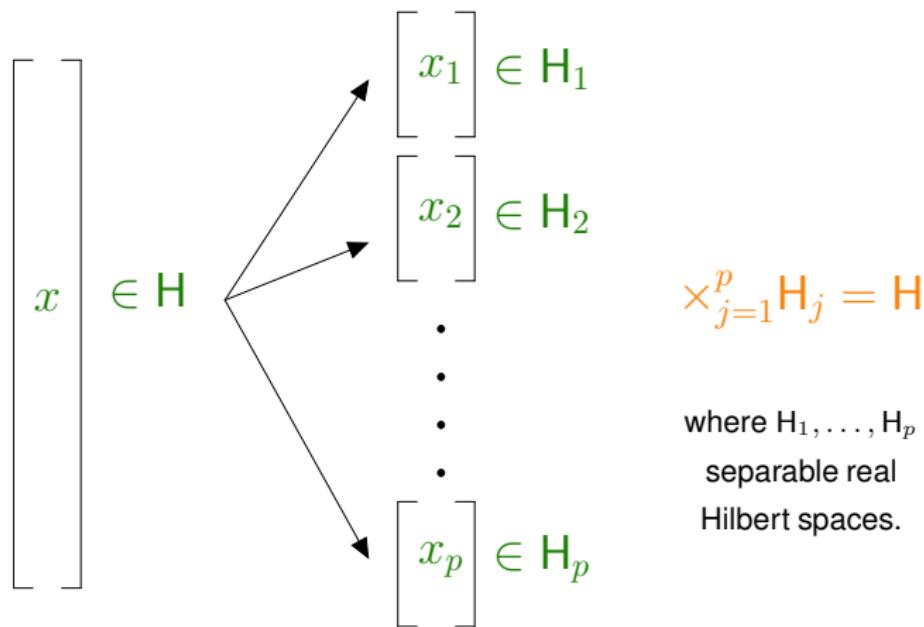
- Methods based on Forward-Backward iteration
 - ▶ type I: [Vu - 2013][Condat - 2013]
(extensions of [Esser *et al.* - 2010][Chambolle and Pock - 2011])
 - ▶ type II: [Combettes *et al.* - 2014]
(extensions of [Loris and Verhoeven - 2011][Chen *et al.* - 2014])
- Methods based on Forward-Backward-Forward iteration
[Combettes and Pesquet - 2012] [Boček and Hendrich, 2014]
- Projection based methods
[Alotaibi *et al.* - 2013]
- ...

Acceleration via block alternation

- Idea: variable splitting.

Acceleration via block alternation

- Idea: variable splitting.



Block-coordinate strategy

- ⇒ At each iteration n , update only a subset of components (~ Gauss-Seidel).

Block-coordinate strategy

⇒ At each iteration n , update only a subset of components (~ Gauss-Seidel).

Advantages:

- Reduced complexity
- Less memory requirements per iteration
- More flexibility

⇒ Useful for large-scale optimization
(see e.g. [Richtárik and Takáč, 2014])

Parallel proximal primal-dual algorithm

Optimization problem: Minimization of

$$(\forall x \in \mathbb{H}) \quad \Phi(x) = \sum_{j=1}^p (f_j(x_j) + h_j(x_j)) + \sum_{k=1}^q (g_k \square l_k)(\mathsf{L}_k x)$$

Parallel proximal primal-dual algorithm

Optimization problem: Minimization of

$$(\forall x \in \mathsf{H}) \quad \Phi(x) = \sum_{j=1}^p (\textcolor{orange}{f}_j(x_j) + \textcolor{orange}{h}_j(x_j)) + \sum_{k=1}^q (\textcolor{orange}{g}_k \square \textcolor{orange}{l}_k)(\mathsf{L}_k x)$$

where

- $(\forall j \in \{1, \dots, p\}) f_j \in \Gamma_0(\mathsf{H}_j)$
- $h_j \in \Gamma_0(\mathsf{H}_j)$ μ_j Lipschitz-differentiable with $\mu_j \in]0, +\infty[$
- $(\forall k \in \{1, \dots, q\}) g_k \in \Gamma_0(\mathsf{G}_k)$, G_k separable real Hilbert space
- $l_k \in \Gamma_0(\mathsf{G}_k)$ ν_k -strongly convex with $\nu_k \in]0, +\infty[$

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- $l_k \in \Gamma_0(\mathsf{G}_k)$ ν_k -strongly convex with $\nu_k \in]0, +\infty[$
- $g_k \square l_k$ inf-convolution of g_k and l_k :

$$(\forall v_k \in \mathsf{G}_k) \quad (g_k \square l_k)(v_k) = \inf_{v'_k \in \mathsf{G}_k} g_k(v'_k) + h_k(v_k - v'_k)$$

$$g_k \square \iota_{\{0\}} = g_k$$

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- $g_k \square l_k$ inf-convolution of g_k and l_k
- $(\forall x \in \mathsf{H}) \mathsf{L}_k x = \sum_{j=1}^p \mathsf{L}_{k,j} x_j$, where
 $(\forall j \in \{1, \dots, p\}) \mathsf{L}_{k,j}: \mathsf{H}_j \rightarrow \mathsf{G}_k$ bounded linear operator with

$$\mathbb{L}_k = \{j \in \{1, \dots, p\} \mid \mathsf{L}_{k,j} \neq 0\} \neq \emptyset$$

$$\mathbb{L}_j^* = \{k \in \{1, \dots, q\} \mid \mathsf{L}_{k,j} \neq 0\} \neq \emptyset.$$

Block parallel proximal primal-dual problem

Primal problem : Find an element of the set F of solutions to

$$\underset{x_1 \in H_1, \dots, x_p \in H_p}{\text{minimize}} \quad \sum_{j=1}^p (f_j(x_j) + h_j(x_j)) + \sum_{k=1}^q (g_k \square l_k) \left(\sum_{j=1}^p L_{k,j} x_j \right)$$

Dual problem : Find an element of the set F^* of solutions to

$$\underset{v_1 \in G_1, \dots, v_q \in G_q}{\text{minimize}} \quad \sum_{j=1}^p (f_j^* \square h_j^*) \left(- \sum_{k=1}^q L_{k,j}^* v_k \right) + \sum_{k=1}^q (g_k^*(v_k) + l_k^*(v_k))$$

For every $j \in \{1, \dots, p\}$, $f_j^* \in \Gamma_0(H_j)$ is the conjugate of f_j such that

$$(\forall y_j \in H_j) \quad f_j^*(y_j) = \sup_{x_j \in H_j} \langle x_j \mid y_j \rangle - f_j(x_j).$$

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We assume that there exists $(\bar{x}_1, \dots, \bar{x}_p) \in H_1 \times \dots \times H_p$ such that

$$(\forall j \in \{1, \dots, p\}) \quad 0 \in \partial f_j(\bar{x}_j) + \nabla h_j(\bar{x}_j) \\ + \sum_{k=1}^q L_{k,j}^* \partial(g_k \square l_k) \left(\sum_{j'=1}^p L_{k,j'} \bar{x}_{j'} \right).$$

Random block-coordinate primal-dual algorithm

```
for n = 0, 1, ...
  for k = 1, ..., q
    uk,n = εp+k,n (proxgk*Uk-1 (vk,n + Uk( ∑j ∈ Lk Lk,j xj,n - ∇lk*(vk,n) + dk,n) ) + bk,n)
    vk,n+1 = vk,n + λn εp+k,n (uk,n - vk,n),
    for j = 1, ..., p
      yj,n = εj,n (proxfjWj-1 (xj,n - Wj( ∑k ∈ Lj* Lk,j* (2uk,n - vk,n) + ∇hj(xj,n) + cj,n) ) + aj,n)
      xj,n+1 = xj,n + λn εj,n (yj,n - xj,n)
```

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      xj,n+1 = xj,n + λn εj,n (yj,n - xj,n)
```

where

- $(\varepsilon_n)_{n \in \mathbb{N}}$ identically distributed \mathbb{D} -valued random variables with
 $\mathbb{D} = \{0, 1\}^{p+q} \setminus \{\mathbf{0}\}$
~~ binary variables signaling the blocks to be activated

Random block-coordinate primal-dual algorithm

```
for n = 0, 1, ...
  for k = 1, ..., q
    u_{k,n} = \varepsilon_{p+k,n} \left( \text{prox}_{g_k^*}^{U_k^{-1}} \left( v_{k,n} + U_k \left( \sum_{j \in \mathbb{L}_k} L_{k,j} x_{j,n} - \nabla l_k^*(v_{k,n}) + d_{k,n} \right) \right) + b_{k,n} \right)
    v_{k,n+1} = v_{k,n} + \lambda_n \varepsilon_{p+k,n} (u_{k,n} - v_{k,n}),
    for j = 1, ..., p
      y_{j,n} = \varepsilon_{j,n} \left( \text{prox}_{f_j}^{W_j^{-1}} \left( x_{j,n} - W_j \left( \sum_{k \in \mathbb{L}_j^*} L_{k,j}^* (2u_{k,n} - v_{k,n}) + \nabla h_j(x_{j,n}) + c_{j,n} \right) + a_{j,n} \right) \right)
      x_{j,n+1} = x_{j,n} + \lambda_n \varepsilon_{j,n} (y_{j,n} - x_{j,n})
```

where

- $(\varepsilon_n)_{n \in \mathbb{N}} \rightsquigarrow$ binary variables signaling the blocks to be activated
- $x_0, (a_n)_{n \in \mathbb{N}}$, and $(c_n)_{n \in \mathbb{N}}$ H-valued random variables, $v_0, (b_n)_{n \in \mathbb{N}}$, and $(d_n)_{n \in \mathbb{N}}$ G-valued random variables with $G = G_1 \times \dots \times G_q$
 $\rightsquigarrow (a_n)_{n \in \mathbb{N}}, (b_n)_{n \in \mathbb{N}}, (c_n)_{n \in \mathbb{N}}$, and $(d_n)_{n \in \mathbb{N}}$: error terms

Random block-coordinate primal-dual algorithm

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where

- (ε_n)_{n ∈ N} ↪ binary variables signaling the blocks to be activated
- (a_n)_{n ∈ N}, (b_n)_{n ∈ N}, (c_n)_{n ∈ N}, and (d_n)_{n ∈ N}: error terms
- (forall j ∈ {1, ..., p}) W_j: H_j → H_j and (forall k ∈ {1, ..., q}) U_k: G_k → G_k strongly positive self-adjoint preconditioning linear operators such that

$$1 - \left(\sum_{j=1}^p \sum_{k=1}^q \|U_k^{1/2} L_{k,j} W_j^{1/2}\|^2 \right)^{1/2} > \frac{1}{2} \max\{(\|W_j\| \mu_j)_{1 \leq j \leq p}, (\|U_k\| \nu_k)_{1 \leq k \leq q}\}.$$

Random block-coordinate primal-dual algorithm

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where

- (ε_n)_{n ∈ N} ↪ binary variables signaling the blocks to be activated
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- (forall j ∈ {1, ..., p}) W_j: H_j → H_j and (forall k ∈ {1, ..., q}) U_k: G_k → G_k strongly positive self-adjoint preconditioning linear operators
- (forall n ∈ N) λ_n ∈]0, 1] such that inf_{n ∈ N} λ_n > 0.

Random block-coordinate primal-dual algorithm

Theorem

Let (Ω, \mathcal{F}, P) be the underlying probability space.

Set $(\forall n \in \mathbb{N}) \mathcal{X}_n = \sigma(x_{n'}, v_{n'})_{0 \leq n' \leq n}$. Assume that

- ① $\sum_{n \in \mathbb{N}} \sqrt{\mathbb{E}(\|a_n\|^2 | \mathcal{X}_n)} < +\infty$, $\sum_{n \in \mathbb{N}} \sqrt{\mathbb{E}(\|b_n\|^2 | \mathcal{X}_n)} < +\infty$,
 $\sum_{n \in \mathbb{N}} \sqrt{\mathbb{E}(\|c_n\|^2 | \mathcal{X}_n)} < +\infty$, and $\sum_{n \in \mathbb{N}} \sqrt{\mathbb{E}(\|d_n\|^2 | \mathcal{X}_n)} < +\infty$
P-a.s.
- ② The variables $(\varepsilon_n)_{n \in \mathbb{N}}$ are identically distributed such that
 $(\forall j \in \{1, \dots, p\}) P[\varepsilon_{j,0} = 1] > 0$.
- ③ For every $n \in \mathbb{N}$, ε_n and \mathcal{X}_n are independent.
- ④ For every $k \in \{1, \dots, q\}$ and $n \in \mathbb{N}$,
$$\bigcup_{j \in \mathbb{L}_k} \{\omega \in \Omega \mid \varepsilon_{j,n}(\omega) = 1\} \subset \{\omega \in \Omega \mid \varepsilon_{p+k,n}(\omega) = 1\}.$$

Then, $(x_n)_{n \in \mathbb{N}}$ converges weakly P-a.s. to an F -valued random variable,
and $(v_n)_{n \in \mathbb{N}}$ converges weakly P-a.s. to an F^* -valued random variable.

Proof: based on properties of quasi-Fejér stochastic sequences [Combettes and Pesquet, 2014].

Random block-coordinate primal-dual algorithm

Theorem

Let (Ω, \mathcal{F}, P) be the underlying probability space.

Set $(\forall n \in \mathbb{N}) \mathcal{X}_n = \sigma(x_{n'}, v_{n'})_{0 \leq n' \leq n}$. Assume that

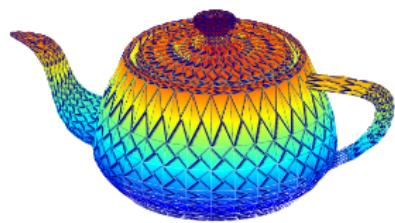
- ① $\sum_{n \in \mathbb{N}} \sqrt{\mathbb{E}(\|a_n\|^2 | \mathcal{X}_n)} < +\infty$, $\sum_{n \in \mathbb{N}} \sqrt{\mathbb{E}(\|b_n\|^2 | \mathcal{X}_n)} < +\infty$,
 $\sum_{n \in \mathbb{N}} \sqrt{\mathbb{E}(\|c_n\|^2 | \mathcal{X}_n)} < +\infty$, and $\sum_{n \in \mathbb{N}} \sqrt{\mathbb{E}(\|d_n\|^2 | \mathcal{X}_n)} < +\infty$
P-a.s.
- ② The variables $(\varepsilon_n)_{n \in \mathbb{N}}$ are identically distributed such that
 $(\forall j \in \{1, \dots, p\}) \mathbb{P}[\varepsilon_{j,0} = 1] > 0$.
- ③ For every $n \in \mathbb{N}$, ε_n and \mathcal{X}_n are independent.
- ④ For every $k \in \{1, \dots, q\}$ and $n \in \mathbb{N}$,

$$\varepsilon_{p+k,n} = \max_{1 \leq j \leq p} \{\varepsilon_{j,n} \mid j \in \mathbb{L}_k\}.$$

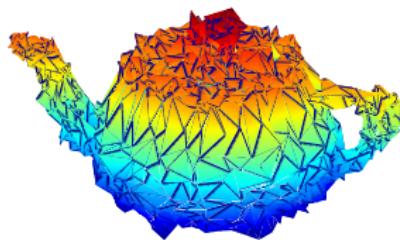
Then, $(x_n)_{n \in \mathbb{N}}$ converges weakly P-a.s. to an \mathcal{F} -valued random variable,
and $(v_n)_{n \in \mathbb{N}}$ converges weakly P-a.s. to an \mathcal{F}^* -valued random variable.

Proof: based on properties of quasi-Fejér stochastic sequences [Combettes and Pesquet, 2014].

Mesh denoising problem



Initial mesh \bar{x}



Observed mesh z

Undirected nonreflexive graph

OBJECTIVE:

Estimate $\bar{x} = (\bar{x}^{(i)})_{1 \leq i \leq M}$ from noisy observations $z = (z^{(i)})_{1 \leq i \leq M}$ where, for every $i \in \{1, \dots, M\}$, $\bar{x}^{(i)} \in \mathbb{R}^3$ is the vector of 3D coordinates of the i -th vertex of a mesh

$$\rightsquigarrow \mathbf{H} = (\mathbb{R}^3)^M$$

Mesh denoising problem

OBJECTIVE:

Estimate $\bar{x} = (\bar{x}^{(i)})_{1 \leq i \leq M}$ from noisy observations $z = (z^{(i)})_{1 \leq i \leq M}$ where, for every $i \in \{1, \dots, M\}$, $\bar{x}^{(i)} \in \mathbb{R}^3$ is the vector of 3D coordinates of the i -th vertex of a mesh

COST FUNCTION:

$$\Phi(x) = \sum_{j=1}^M \psi_j(x^{(j)} - z^{(j)}) + \iota_{C_j}(x^{(j)}) + \eta_j \|(x^{(j)} - x^{(i)})_{i \in \mathcal{N}_j}\|_{1,2}$$

where, for every $j \in \{1, \dots, M\}$,

- $\psi_j: \mathbb{R}^3 \rightarrow \mathbb{R}$: $\ell_2 - \ell_1$ Huber function
 - ▶ robust data fidelity measure
- C_j : nonempty convex subset of \mathbb{R}^3
 - ▶ box constraint
- \mathcal{N}_j : neighborhood of j -th vertex
- $(\eta_j)_{1 \leq j \leq M}$: nonnegative regularization constants.

Mesh denoising problem

OBJECTIVE:

Estimate $\bar{x} = (\bar{x}^{(i)})_{1 \leq i \leq M}$ from noisy observations $z = (z^{(i)})_{1 \leq i \leq M}$ where, for every $i \in \{1, \dots, M\}$, $\bar{x}^{(i)} \in \mathbb{R}^3$ is the vector of 3D coordinates of the i -th vertex of a mesh

COST FUNCTION:

$$\Phi(x) = \sum_{j=1}^M \psi_j(x^{(j)} - z^{(j)}) + \iota_{C_j}(x^{(j)}) + \eta_j \|(x^{(j)} - x^{(i)})_{i \in \mathcal{N}_j}\|_{1,2}$$

IMPLEMENTATION DETAILS:

a block \equiv a vertex $\Rightarrow p = M$

Algorithm type I:

$$(\forall j \in \{1, \dots, M\})$$

- $h_j = \psi_j(\cdot - z^{(j)})$

- $f_j = \iota_{C_j}$

$$q = M$$

$$(\forall k \in \{1, \dots, M\})(\forall x \in \mathsf{H})$$

- $g_k(\mathsf{L}_k x) = \|(x^{(k)} - x^{(i)})_{i \in \mathcal{N}_k}\|_{1,2}$

- $l_k = \iota_{\{0\}}$

Algorithm type II:

$$(\forall j \in \{1, \dots, M\})$$

- $h_j = \psi_j(\cdot - z^{(j)})$

$$q = 2M$$

$$(\forall k \in \{1, \dots, M\})(\forall x \in \mathsf{H})$$

- $g_k(\mathsf{L}_k x) = \|(x^{(k)} - x^{(i)})_{i \in \mathcal{N}_k}\|_{1,2}$

- $g_{M+k}(\mathsf{L}_{M+k} x) = \iota_{C_k}(x^{(k)})$

- $l_k = \iota_{\{0\}}$

Simulation results (algorithm type II)

- positions of the original mesh are corrupted through an i.i.d zero-mean Gaussian mixture noise model.
- a limited number r of variables can be handled at each iteration, where

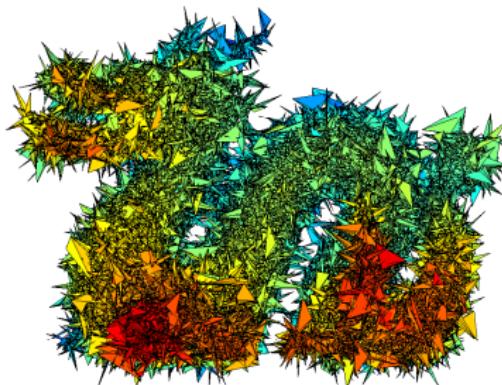
$$\sum_{j=1}^p \varepsilon_{j,n} = r \leq p.$$

- mesh decomposed into p/r non-overlapping sets.



Original mesh,

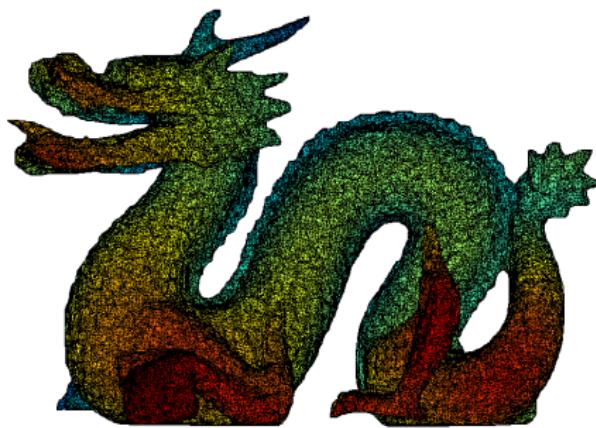
$M = 100250$.



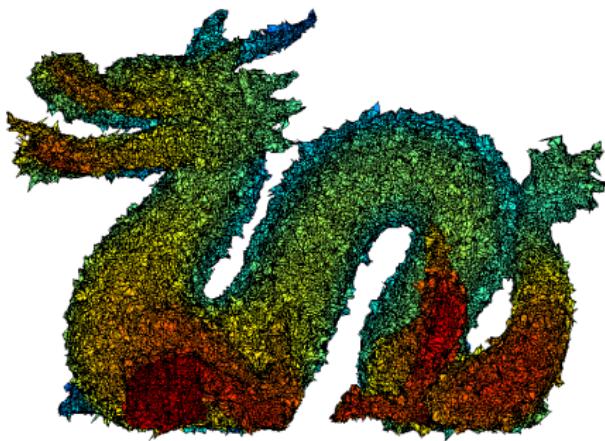
Noisy mesh,

$\text{MSE} = 2.89 \times 10^{-6}$.

Simulation results (algorithm type II)

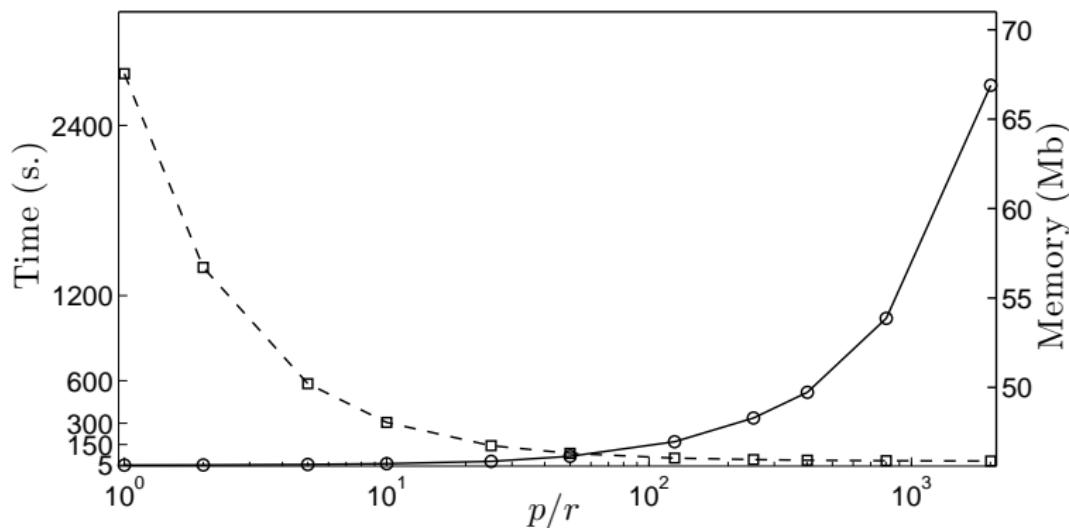


Proposed reconstruction,
 $MSE = 8.09 \times 10^{-8}$.



Laplacian smoothing,
 $MSE = 5.23 \times 10^{-7}$.

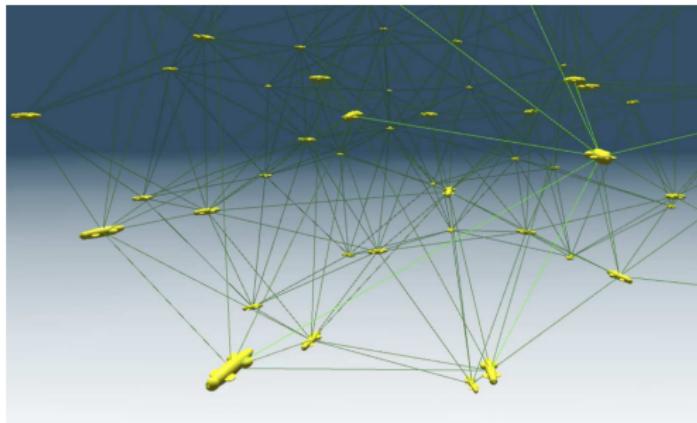
Complexity (algorithm type II)



- dashed line: required memory
- continuous line: reconstruction time

Available extensions

Asynchronous distributed algorithms
(stochastic, primal-dual, proximal, defined on a hypergraph)
[Pesquet and Repetti - 2014][Bianchi *et al.* - 2014]



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