

# Advanced decomposition methods for discrete-time stochastic optimal control problems



## Dual Approximate Dynamic Programming

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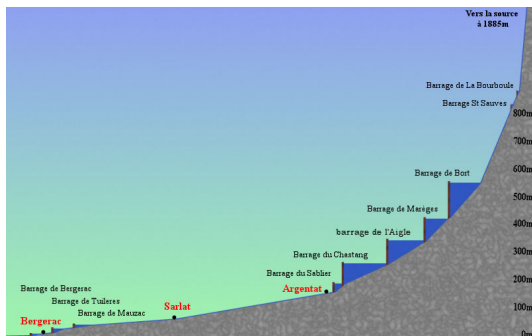
## Ultimate goal of the lecture

How to obtain “good” **strategies** (or **cost-to-go functions**) for a **large scale** stochastic optimal control problem in discrete time, for example a problem corresponding to the optimal management over a given time horizon of a system involving a large amount of **dynamical** production units.

- In order to obtain **decision strategies** (closed-loop controls), we have to use **dynamic programming** or related methods.
  - **Assumption**: Markovian case,
  - **Difficulty**: **curse of dimensionality**.
- To overcome the barrier of the dimension, we want to use **decomposition/coordination** techniques, so that we have to take into account the **information pattern** induced by the stochastic optimization problem.

## Practical applications under consideration

### Electricity production management for large hydro valleys



- *1 year time horizon:*  
compute each month  
the “values of water”  
(cost-to-go functions)
- *stochastic framework:*  
rain, market prices
- *large-scale valley:*  
4 dams and much more

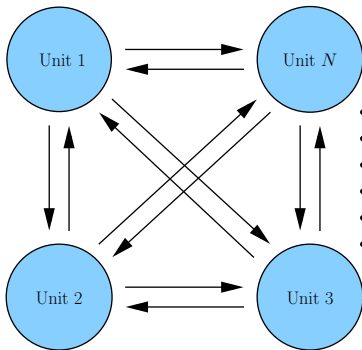
## Lecture outline

- 1 Examples and mathematical background
  - Interconnected systems
  - Optimization background
  - Standard decomposition methods
- 2 About decomposition in stochastic optimization
  - Couplings in stochastic optimization
  - Dynamic programming and decomposition
- 3 Dual approximate dynamic programming (DADP)
  - Problem formulation and price decomposition
  - Subproblems resolution and coordination
  - What has been really done?
- 4 Hydro valleys management problem
  - DADP implementation for hydro valleys
  - Numerical results for different valleys

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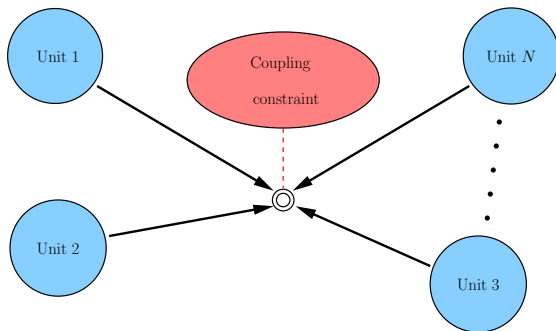
# Decomposition and coordination



Interconnected units

- The “**large system**” to be optimized consists of **interconnected** subsystems: we want to use this structure in order to formulate optimization **subproblems** of **reasonable** complexity.
- But the presence of **interactions** requires a level of **coordination**.
- Coordination must provide a **local model** of the interactions to each subproblem: it is an **iterative** process.
- The ultimate goal is to obtain the solution of the **overall problem** by concatenation of the solutions of the **subproblems**.

# Example: the “flower model”



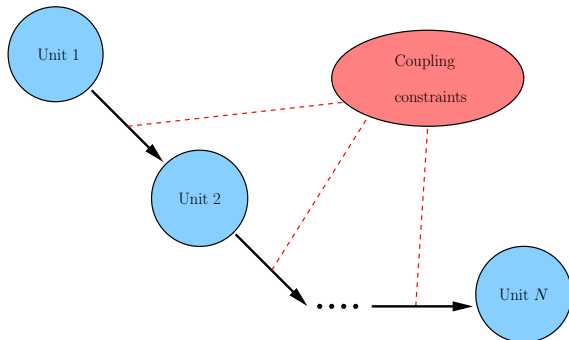
## Unit Commitment Problem

$$\min_u \sum_{i=1}^N J_i(u_i),$$

$$\text{s.t.} \quad \sum_{i=1}^N \Theta_i(u_i) = \theta.$$



# Example: the “cascade model”



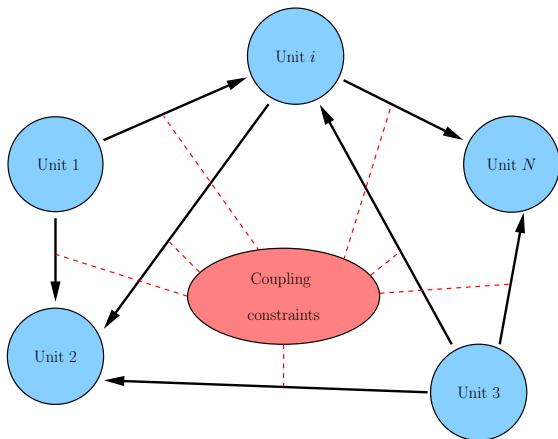
$$\min_{u,v} \sum_{i=1}^N J_i(u_i, v_i),$$

$$\text{s.t. } H_i(u_i, v_i) = v_{i+1} \quad \forall i.$$

## Dams Management Problem

Link with the flower model:  $\Theta_i(u_i, v_i) = (0, \dots, -v_i, H_i(u_i, v_i), \dots, 0)^\top$ .

# A general model



Smart Grid

$$\min_{u,v} \sum_{i=1}^N J_i \left( u_i, \sum_{j \neq i} v_{j,i} \right),$$

$$\text{s.t. } H_i \left( u_i, \sum_{j \neq i} v_{j,i} \right) = v_i.$$

## Motivation for theoretical developments

### Mathematical ingredients needed to tackle such problems.

- **Optimization** mathematical framework.
- **Duality** theory (handling of constraints).
  - ↪ Lagrangian relaxation.
- **Decomposition/coordination** methods
  - ↪ Price decomposition (Walras groping).

### Points already covered since the beginning of the week.

- **Stochastic** optimization.
- **Dynamic programming**.

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# Optimization without explicit constraint

$$\min_{u \in \mathcal{U}^{\text{ad}}} J(u) .$$

- $\mathcal{U}$ : Hilbert space with scalar product  $\langle \cdot, \cdot \rangle$ .  
Examples:  $\mathcal{U} = \mathbb{R}^n$  (vectors) or  $\mathcal{U} = L^2(\Omega, \mathcal{A}, \mathbb{P}; \mathbb{R}^n)$  (random variables).
- $\mathcal{U}^{\text{ad}}$ : closed convex subset of  $\mathcal{U}$ .
- $J : \mathcal{U} \rightarrow \mathbb{R}$ : function satisfying some properties  
(convexity, continuity, differentiability, coercivity).

**Characterization** of a solution  $u^\sharp$  (**optimality conditions**):

$$\langle \nabla J(u^\sharp), u - u^\sharp \rangle \geq 0 \quad \forall u \in \mathcal{U}^{\text{ad}} .$$

**Computation** of the solution  $u^\sharp$  (**projected gradient algorithm**):

$$u^{(k+1)} = \text{proj}_{\mathcal{U}^{\text{ad}}} \left( u^{(k)} - \rho \nabla J(u^{(k)}) \right) .$$

# Optimization with explicit constraints

$$\min_{u \in \mathcal{U}^{\text{ad}}} J(u) \quad \text{subject to} \quad \Theta(u) \in -C. \quad (\mathcal{P})$$

- $\mathcal{U}$ : Hilbert space.
- $\mathcal{U}^{\text{ad}}$ : closed convex subset of  $\mathcal{U}$ .
- $J : \mathcal{U} \rightarrow \mathbb{R}$ : cost function.
- $\mathcal{V}$ : another Hilbert space.
- $\Theta : \mathcal{U} \rightarrow \mathcal{V}$ : constraint function satisfying some properties (convexity w.r.t.  $C$ , continuity, differentiability).
- $C$ : **cone** of  $\mathcal{V}$  (examples:  $C = \{0\}$ ,  $C = \{v \geq 0\}$ ).

*An additional condition on the constraint function is needed!*

**Constraint Qualification Condition**, e.g.  $0 \in \text{int}(\Theta(\mathcal{U}^{\text{ad}}) + C)$ .

# Optimization with explicit constraints



## Karush-Kuhn-Tucker Conditions

In addition to standard conditions on  $J$  and  $\Theta$ , we assume that the constraints are **qualified**.

Then a **necessary and sufficient** condition for  $u^\# \in \mathcal{U}^{\text{ad}}$  to be a solution of Problem  $(\mathcal{P})$  is that there exists  $\lambda^\# \in \mathcal{V}$  such that:

- ①  $\langle \nabla J(u^\#) + [\Theta'(u^\#)]^\top \lambda^\#, u - u^\# \rangle \geq 0 \quad \forall u \in \mathcal{U}^{\text{ad}},$
- ②  $\Theta(u^\#) \in -C,$
- ③  $\lambda^\# \in C^*,$
- ④  $\langle \lambda^\#, \Theta(u^\#) \rangle = 0$  (**Complementary Slackness**).

The **dual cone** of  $C$  is defined by:  $C^* = \{\lambda \in \mathcal{V}, \langle \lambda, v \rangle \geq 0 \quad \forall v \in C\}.$

# Optimization with explicit constraints



Let  $L : \mathcal{U}^{\text{ad}} \times C^* \rightarrow \mathbb{R}$  be the **Lagrangian** associated to  $(\mathcal{P})$ :

$$L(u, \lambda) = J(u) + \langle \lambda, \Theta(u) \rangle .$$

A point  $(u^\sharp, \lambda^\sharp) \in \mathcal{U}^{\text{ad}} \times C^*$  is a **saddle point** of  $L$  if

$$L(u^\sharp, \lambda) \leq L(u^\sharp, \lambda^\sharp) \leq L(u, \lambda^\sharp) \quad \forall (u, \lambda) \in \mathcal{U}^{\text{ad}} \times C^* .$$

- If  $(u^\sharp, \lambda^\sharp)$  is a **saddle point** of  $L$ , then  $u^\sharp$  is a **solution** of  $(\mathcal{P})$ .
- If  $u^\sharp$  is a **solution** of  $(\mathcal{P})$  and if the **KKT** conditions are met for some  $\lambda^\sharp$ , then  $(u^\sharp, \lambda^\sharp)$  is a **saddle point** of  $L$ .

Moreover we have that

$$L(u^\sharp, \lambda^\sharp) = J(u^\sharp) = \min_{u \in \mathcal{U}^{\text{ad}}} \max_{\lambda \in C^*} L(u, \lambda) = \max_{\lambda \in C^*} \min_{u \in \mathcal{U}^{\text{ad}}} L(u, \lambda) .$$



# Optimization with explicit constraints

IV

Dealing with the **dual problem**

$$\max_{\lambda \in C^*} \min_{u \in \mathcal{U}^{\text{ad}}} L(u, \lambda),$$

paves the way for algorithmic methods. Define the **dual function**  $\Phi$  associated to the Lagrangian  $L$  as

$$\Phi(\lambda) = \min_{u \in \mathcal{U}^{\text{ad}}} L(u, \lambda).$$

The problem of **maximizing** the dual function  $\Phi$  is equivalent to the one of solving the **dual** problem:

$$\max_{\lambda \in C^*} \Phi(\lambda) \iff \max_{\lambda \in C^*} \min_{u \in \mathcal{U}^{\text{ad}}} L(u, \lambda).$$

The **gradient** of  $\Phi$  is obtained from the minimization step in  $u$ :

$$\nabla \Phi(\lambda) = \Theta(\hat{u}_\lambda), \text{ with } \hat{u}_\lambda \text{ unique solution of } \min_{u \in \mathcal{U}^{\text{ad}}} L(u, \lambda).$$

# Optimization with explicit constraints



In order to obtain a solution of the original constrained problem, we use a **gradient algorithm** for maximizing the dual function:

$$\max_{\lambda \in C^*} \Phi(\lambda).$$

The gradient of  $\Phi$  at the current point  $\lambda^{(k)}$  of the algorithm is obtained by minimizing  $L(u, \lambda^{(k)})$  w.r.t.  $u$ .

## Uzawa's Algorithm

Choose  $\lambda^{(0)} \in C^*$ . At each iteration  $k$ ,

- ① obtain the solution  $u^{(k+1)} = \arg \min_{u \in \mathcal{U}^{\text{ad}}} J(u) + \langle \lambda^{(k)}, \Theta(u) \rangle$ ,
- ② update the multiplier  $\lambda^{(k+1)} = \text{proj}_{C^*} (\lambda^{(k)} + \rho \Theta(u^{(k+1)}))$ .

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# Specific problem structure: additive model

Consider the optimization problem with explicit constraints:

$$\min_{u \in \mathcal{U}^{\text{ad}} \subset \mathcal{U}} J(u) \quad \text{subject to} \quad \Theta(u) = \theta \in \mathcal{V}.$$

We assume that the space  $\mathcal{U}$  writes as a **Cartesian product**:

$$\mathcal{U} = \mathcal{U}_1 \times \cdots \times \mathcal{U}_N, \quad \text{so that} \quad u = (u_1, \dots, u_N) \quad \text{with} \quad u_i \in \mathcal{U}_i.$$

We moreover assume that this space decomposition is such that

- the admissible set  $\mathcal{U}^{\text{ad}}$  writes as a **Cartesian product**:

$$\mathcal{U}^{\text{ad}} = \mathcal{U}_1^{\text{ad}} \times \cdots \times \mathcal{U}_N^{\text{ad}} \quad \text{with} \quad \mathcal{U}_i^{\text{ad}} \subset \mathcal{U}_i,$$

- the functions  $J$  and  $\Theta$  write **additively**:

$$J(u) = J_1(u_1) + \cdots + J_N(u_N),$$

$$\Theta(u) = \Theta_1(u_1) + \cdots + \Theta_N(u_N).$$

# Specific problem structure: additive model

II

Then the original problem displays the so-called **additive structure**:

$$\min_{\substack{u_1 \in \mathcal{U}_1^{\text{ad}} \\ \vdots \\ u_N \in \mathcal{U}_N^{\text{ad}}}} \sum_{i=1}^N J_i(u_i) \quad \text{subject to} \quad \sum_{i=1}^N \Theta_i(u_i) - \theta = 0 .$$

Note that the **coupling** between the  $i$ 's **only** arises from the constraint  $\Theta$ . As a matter of fact,

$$\min_{\substack{u_1 \in \mathcal{U}_1^{\text{ad}} \\ \vdots \\ u_N \in \mathcal{U}_N^{\text{ad}}}} \sum_{i=1}^N J_i(u_i) \quad \Longleftrightarrow \quad \min_{u_i \in \mathcal{U}_i^{\text{ad}}} J_i(u_i) \quad \forall i = 1, \dots, N .$$

## Additive model — Price decomposition

$$\min_{u \in \mathcal{U}^{\text{ad}}} \sum_{i=1}^N J_i(u_i) \quad \text{subject to} \quad \sum_{i=1}^N \Theta_i(u_i) - \theta = 0.$$

- ① Form the **Lagrangian** of the problem. The dual problem writes:

$$\max_{\lambda \in \mathcal{V}} \min_{u \in \mathcal{U}^{\text{ad}}} \sum_{i=1}^N \left( J_i(u_i) + \langle \lambda, \Theta_i(u_i) \rangle \right) - \langle \lambda, \theta \rangle.$$

- ② Solve this problem by the **Uzawa algorithm**:

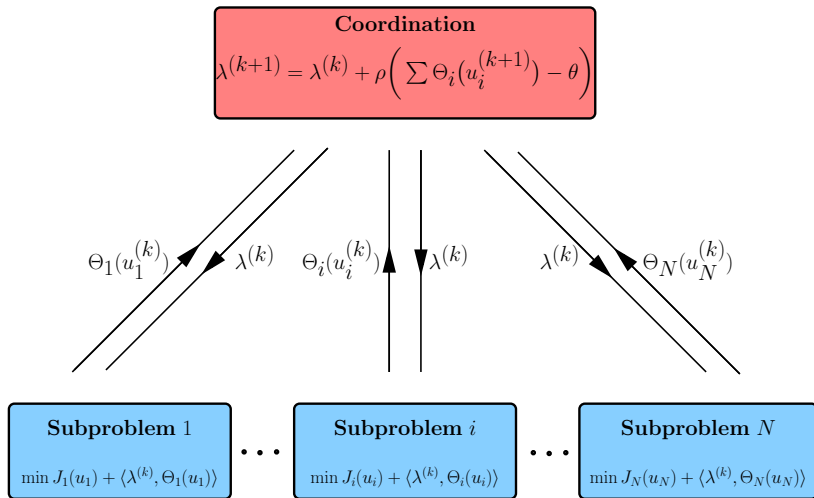
$$u_i^{(k+1)} \in \arg \min_{u_i \in \mathcal{U}_i^{\text{ad}}} J_i(u_i) + \langle \lambda^{(k)}, \Theta_i(u_i) \rangle, \quad i = 1, \dots, N,$$

$$\lambda^{(k+1)} = \lambda^{(k)} + \rho \left( \sum_{i=1}^N \Theta_i(u_i^{(k+1)}) - \theta \right).$$

$\rightsquigarrow$  *Walras groping*

## Additive model — Price decomposition

II



## Additive model — Resource allocation

$$\min_{u \in \mathcal{U}^{\text{ad}}} \sum_{i=1}^N J_i(u_i) \quad \text{subject to} \quad \sum_{i=1}^N \Theta_i(u_i) - \theta = 0 .$$

- 1 Write the constraint in a equivalent manner by introducing **new variables**  $(v_1, \dots, v_N)$  (the so-called “allocation”):

$$\sum_{i=1}^N \Theta_i(u_i) - \theta = 0 \quad \Leftrightarrow \quad \Theta_i(u_i) - v_i = 0 \quad \text{and} \quad \sum_{i=1}^N v_i - \theta = 0 .$$

Minimize the criterion w.r.t.  $u$  and  $v$ :

$$\min_{v \in \mathcal{V}^N} \sum_{i=1}^N \left( \min_{u_i \in \mathcal{U}_i^{\text{ad}}} J_i(u_i) \text{ s.t. } \Theta_i(u_i) - v_i = 0 \right) \text{ s.t. } \sum_{i=1}^N v_i - \theta = 0 .$$



## Additive model — Resource allocation

II

$$\min_{v \in \mathcal{V}^N} \sum_{i=1}^N \underbrace{\left( \min_{u_i \in \mathcal{U}_i^{\text{ad}}} J_i(u_i) \text{ s.t. } \Theta_i(u_i) - v_i = 0 \right)}_{G_i(v_i)} \text{ s.t. } \sum_{i=1}^N v_i - \theta = 0,$$

$$\min_{v \in \mathcal{V}^N} \sum_{i=1}^N G_i(v_i) \text{ s.t. } \sum_{i=1}^N v_i - \theta = 0.$$

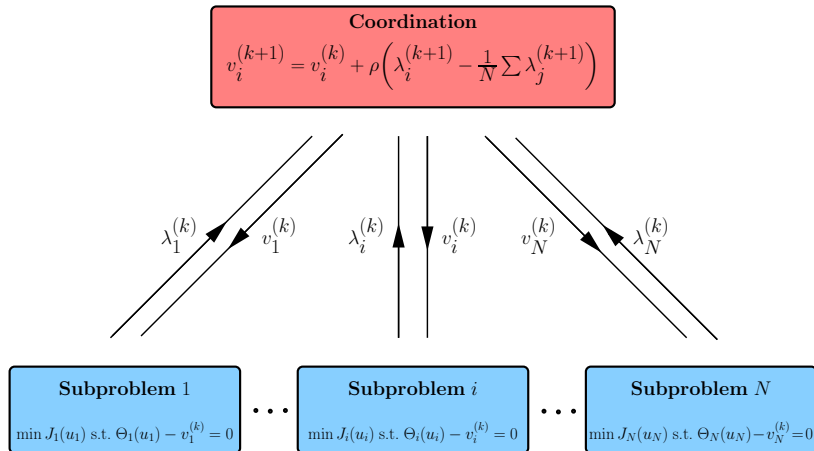
- 2 Solve the last problem using a **projected gradient method**:

$$G_i(v_i^{(k)}) = \min_{u_i \in \mathcal{U}_i^{\text{ad}}} J_i(u_i) \text{ s.t. } \Theta_i(u_i) - v_i^{(k)} = 0 \rightsquigarrow \lambda_i^{(k+1)},$$

$$v_i^{(k+1)} = v_i^{(k)} + \rho \left( \lambda_i^{(k+1)} - \frac{1}{N} \sum_{j=1}^N \lambda_j^{(k+1)} \right).$$

## Additive model — Resource allocation

III



# Additive model: conclusions

## ① Price decomposition

- **Pros:** subproblems are always feasible.
- **Cons:** admissible solution once convergence achieved.

## ② Resource allocation

- **Pros:** admissible solution at each iteration.
- **Cons:** potential existence of unfeasible subproblems.

*Straightforward extension to inequality constraints. . .*

Other methods are available, even for **non-additive structures**.

## References on decomposition/coordination methods

### Further readings on decomposition/coordination:

G. COHEN, "Optimisation des grands systèmes". *Cours du DEA Modélisation et Méthodes Mathématiques en Économie*, 2004.

G. COHEN, "Auxiliary Problem Principle and Decomposition of Optimization Problems". *Journal of Optimization Theory and Applications*, **32**, 1980.

G. COHEN & D.L. ZHU, "Decomposition coordination methods in large scale optimization problems. The nondifferentiable case and the use of augmented Lagrangians". In J.B. Cruz (Ed.): "*Advances in Large Scale Systems*", **1**, 203-266, JAI Press, Greenwich, Connecticut, 1984.

# Final remarks on decomposition methods

The theory is available for general (infinite dimensional) Hilbert spaces, and thus applies in the **stochastic framework**, that is, the case where  $\mathcal{U}$  is a space of **random variables**.

The **minimization algorithm** used for solving the subproblems is not specified in the decomposition process and is left to the user! It is however assumed that the user is able to solve the subproblem, for example in the price decomposition case:

$$\min_{u_i \in \mathcal{U}_i^{\text{ad}}} J_i(u_i) + \langle \lambda^{(k)}, \Theta_i(u_i) \rangle,$$

and to send the requested information, namely  $\Theta_i(u_i^{(k+1)})$ , to the coordination level.

**Question:** *what methods are suitable for a **stochastic problem**?*

# Final remarks on decomposition methods



Whatever the **decomposition/coordination scheme** used (price, allocation. . . ), **new variables** (depending on  $u^{(k)}$  and/or  $\lambda^{(k)}$ ) are present in the subproblems arising at iteration  $k$  of the associated algorithm.

**Example:** subproblem  $i$  in price decomposition:

$$\min_{u_i \in \mathcal{U}_i^{\text{ad}}} J_i(u_i) + \langle \lambda^{(k)}, \Theta_i(u_i) \rangle .$$

All these new variables are considered as **fixed** when solving the subproblems (they only depend on the iteration index  $k$ ). They are nothing but constant elements in the space  $\mathcal{V}$ .

**Question:** *which consequences in the **stochastic case**?*

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## Reminder of our ultimate goal

How to obtain “good” **strategies** (or **cost-to-go functions**) for a **large scale** stochastic optimal control problem in discrete time, for example a problem corresponding to the optimal management over a given time horizon of a system involving a large amount of **dynamical** production units.

- In order to obtain **decision strategies** (closed-loop controls), we have to use **dynamic programming** or related methods.
  - **Assumption**: Markovian case,
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- To overcome the barrier of the dimension, we want to use **decomposition/coordination** techniques, so that we have to take into account the **information pattern** induced by the stochastic optimization problem.



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## Stochastic optimal control problem in discrete time

We consider a **stochastic optimal control (SOC)** problem

$$\min_{\mathbf{U}, \mathbf{X}} \mathbb{E} \left( \sum_{i=1}^N \left( \sum_{t=0}^{T-1} L_t^i(\mathbf{x}_t^i, \mathbf{u}_t^i, \mathbf{w}_{t+1}) + K^i(\mathbf{x}_T^i) \right) \right),$$

subject to the constraints:

$$\mathbf{x}_0^i = f_{-1}^i(\mathbf{w}_0), \quad i = 1 \dots N,$$

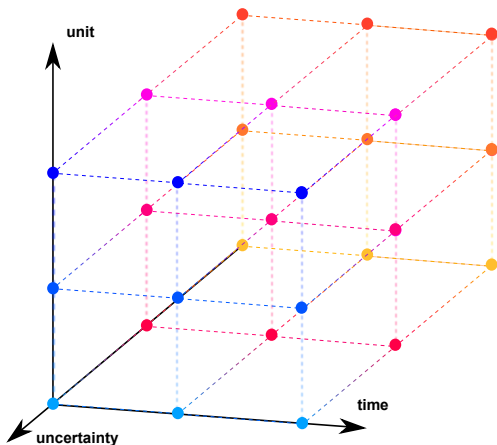
$$\mathbf{x}_{t+1}^i = f_t^i(\mathbf{x}_t^i, \mathbf{u}_t^i, \mathbf{w}_{t+1}), \quad t = 0 \dots T-1, \quad i = 1 \dots N,$$

$$\mathbf{u}_t^i = \mathbb{E}(\mathbf{u}_t^i \mid \mathcal{F}_t), \quad t = 0 \dots T-1, \quad i = 1 \dots N,$$

$$\sum_{i=1}^N \Theta_t^i(\mathbf{x}_t^i, \mathbf{u}_t^i) = 0, \quad t = 0 \dots T-1.$$

# Couplings and decompositions for SOC problems

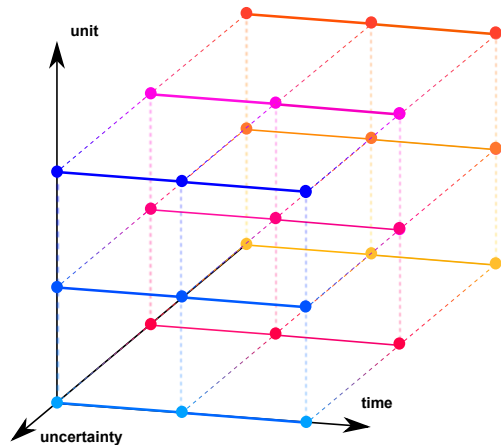
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$$\min \sum_{\omega} \sum_i \sum_t \pi_{\omega} L_t^i(\mathbf{x}_t^i, \mathbf{u}_t^i, \mathbf{w}_{t+1})$$

# Couplings and decompositions for SOC problems

II

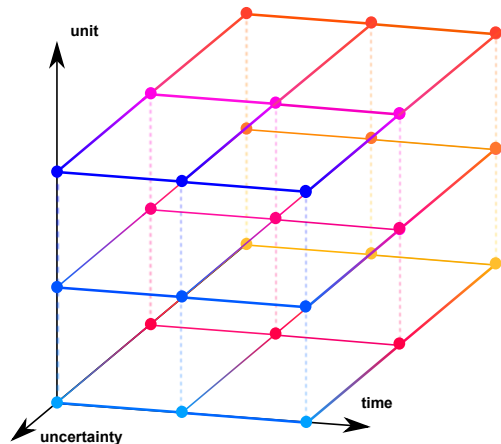


$$\min \sum_{\omega} \sum_i \sum_t \pi_{\omega} L_t^i(\mathbf{x}_t^i, \mathbf{u}_t^i, \mathbf{w}_{t+1})$$

$$\text{s.t. } \mathbf{x}_{t+1}^i - f_t^i(\mathbf{x}_t^i, \mathbf{u}_t^i, \mathbf{w}_{t+1}) = 0$$

# Couplings and decompositions for SOC problems

III



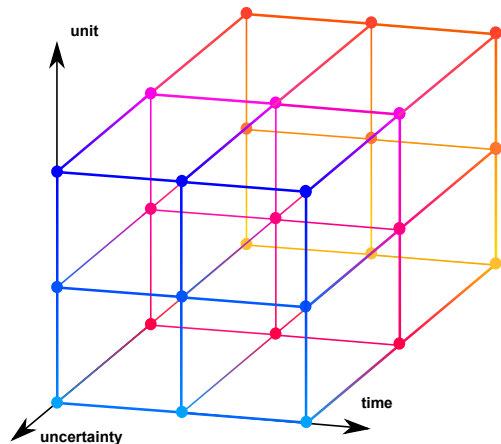
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$$\mathbf{u}_t^i - \mathbb{E}(\mathbf{u}_t^i \mid \mathcal{F}_t) = 0$$

# Couplings and decompositions for SOC problems

IV



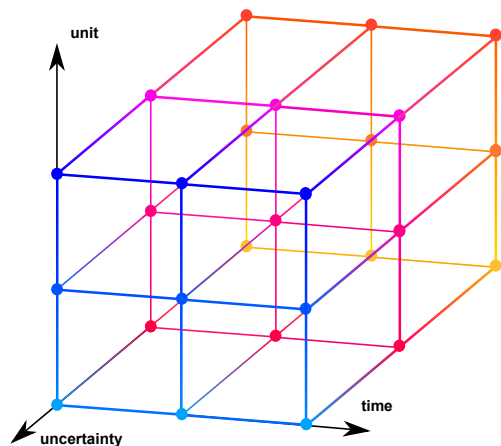
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# Couplings and decompositions for SOC problems



$$\min \sum_{\omega} \sum_i \sum_t \pi_{\omega} L_t^i(\mathbf{x}_t^i, \mathbf{u}_t^i, \mathbf{w}_{t+1})$$

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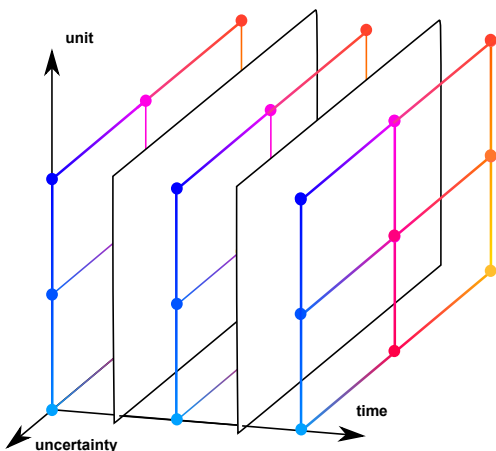
$$\mathbf{u}_t^i - \mathbb{E}(\mathbf{u}_t^i \mid \mathcal{F}_t) = 0$$

$$\sum_i \Theta_t^i(\mathbf{x}_t^i, \mathbf{u}_t^i) = 0$$

3 additive structures!

# Couplings and decompositions for SOC problems

## VI



$$\min \sum_{\omega} \sum_i \sum_t \pi_{\omega} L_t^i(\mathbf{x}_t^i, \mathbf{u}_t^i, \mathbf{w}_{t+1})$$

$$\text{s.t. } \mathbf{x}_{t+1}^i - f_t^i(\mathbf{x}_t^i, \mathbf{u}_t^i, \mathbf{w}_{t+1}) = 0$$

$$\mathbf{u}_t^i - \mathbb{E}(\mathbf{u}_t^i | \mathcal{F}_t) = 0$$

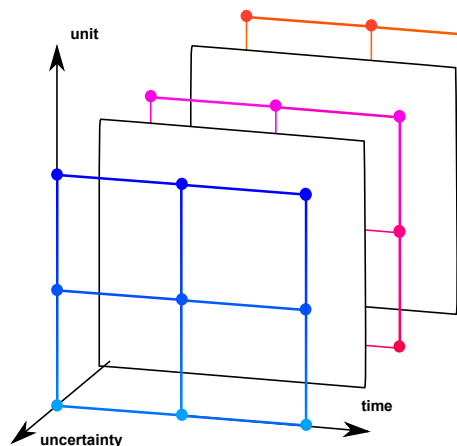
$$\sum_i \Theta_t^i(\mathbf{x}_t^i, \mathbf{u}_t^i) = 0$$

Time decomposition



# Couplings and decompositions for SOC problems

## VII



$$\min \sum_{\omega} \sum_i \sum_t \pi_{\omega} L_t^i(\mathbf{x}_t^i, \mathbf{u}_t^i, \mathbf{w}_{t+1})$$

$$\text{s.t. } \mathbf{x}_{t+1}^i - f_t^i(\mathbf{x}_t^i, \mathbf{u}_t^i, \mathbf{w}_{t+1}) = 0$$

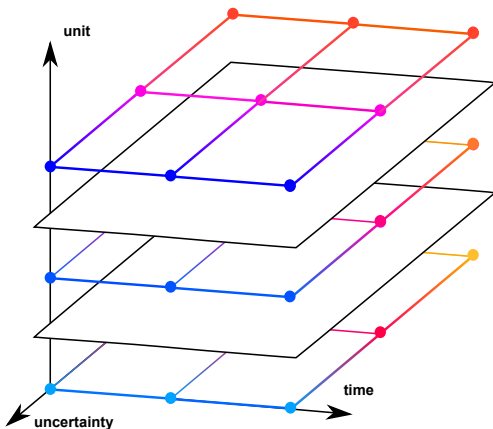
$$\mathbf{u}_t^i - \mathbb{E}(\mathbf{u}_t^i \mid \mathcal{F}_t) = 0$$

$$\sum_i \Theta_t^i(\mathbf{x}_t^i, \mathbf{u}_t^i) = 0$$

Progressive hedging

# Couplings and decompositions for SOC problems

## VIII



$$\min \sum_{\omega} \sum_i \sum_t \pi_{\omega} L_t^i(\mathbf{x}_t^i, \mathbf{u}_t^i, \mathbf{w}_{t+1})$$

$$\text{s.t. } \mathbf{x}_{t+1}^i - f_t^i(\mathbf{x}_t^i, \mathbf{u}_t^i, \mathbf{w}_{t+1}) = 0$$

$$\mathbf{u}_t^i - \mathbb{E}(\mathbf{u}_t^i | \mathcal{F}_t) = 0$$

$$\sum_i \Theta_t^i(\mathbf{x}_t^i, \mathbf{u}_t^i) = 0$$

Purpose of DADP

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# Mixing spatial decomposition and dynamic programming

Consider the SOC problem (the **spatial structure** is highlighted)

$$\min_{\mathbf{U}, \mathbf{X}} \sum_{i=1}^N \left( \mathbb{E} \left( \sum_{t=0}^{T-1} L_t^i(\mathbf{X}_t^i, \mathbf{U}_t^i, \mathbf{W}_{t+1}) + K^i(\mathbf{X}_T^i) \right) \right),$$

subject to the constraints:

$$\mathbf{X}_{t+1}^i = f_t^i(\mathbf{X}_t^i, \mathbf{U}_t^i, \mathbf{W}_{t+1}), \quad t = 0 \dots T-1, \quad i = 1 \dots N,$$

$$\mathbf{U}_t^i = \mathbb{E}(\mathbf{U}_t^i \mid \mathcal{F}_t), \quad t = 0 \dots T-1, \quad i = 1 \dots N,$$

$$\sum_{i=1}^N \Theta_t^i(\mathbf{X}_t^i, \mathbf{U}_t^i) = 0, \quad t = 0 \dots T-1 \quad \rightsquigarrow \quad \mathbf{\Lambda}_t,$$

and assume that the random variables  $\mathbf{W}_t$  are **independent** (**white noise** assumption).

## Dynamic programming yields centralized controls

Under the **white noise** assumption, it is possible to use **dynamic programming (DP)** in order to solve the SOC problem.

The true **optimal** control  $U_t^i$  of unit  $i$  is a feedback of the **whole** system state, that is, a function of all  $X_t^i$ 's:

$$U_t^i = \gamma_t^i(X_t^1, \dots, X_t^N).$$

Of course, a **straightforward use** of **DP** is prohibited for  $N$  large (curse of dimensionality), and decomposition is needed!

But the feedback  $\gamma_t^i$  structurally induces a coupling between all the units! Moreover, a **naive decomposition** of the problem should lead to **decentralized feedbacks**:

$$U_t^i = \hat{\gamma}_t^i(X_t^i),$$

which, in most cases, are **far from being optimal**...

## Price decomposition in the stochastic case

Apply **price decomposition** to the SOC problem by dualizing the spatial coupling constraint. Then a dual multiplier  $\Lambda_t^{(k)}$  appears in each subproblem  $i$  at each iteration  $k$ :

$$\min_{\mathbf{u}^i, \mathbf{x}^i} \mathbb{E} \left( \sum_{t=0}^{T-1} (L_t^i(\mathbf{x}_t^i, \mathbf{u}_t^i, \mathbf{w}_{t+1}) + \Lambda_t^{(k)} \cdot \Theta_t^i(\mathbf{x}_t^i, \mathbf{u}_t^i)) + K^i(\mathbf{x}_T^i) \right).$$

The  $\Lambda_t^{(k)}$ 's are **fixed random variables** at this step of the algorithm. Subproblem  $i$  encompasses **two** noise variables, namely  $\mathbf{w}_{t+1}$  and  $\Lambda_t^{(k)}$ , but the  $\Lambda_t^{(k)}$ 's may be **correlated** in time, in which case the **white noise** assumption fails!

Otherwise stated, the original state  $\mathbf{x}_t^i$  is not a “good” state for subproblem  $i$ : the feature which seemed to have been won by decomposition is actually lost again by dynamic programming.

## Summary

- On the one hand, it seems that dynamic programming **cannot be decomposed** in a straightforward manner.
- On the other hand, applying a decomposition scheme to a SOC problem introduces **coordination instruments** in the subproblems, e.g. the multipliers  $\Lambda_t^{(k)}$  in the case of price decomposition. They correspond to additional fixed random variables whose **time structure** is **unknown**,<sup>1</sup> so that dynamic programming cannot be used for solving the subproblems

**Question:** how to **handle** these coordination instruments in order to be able to use dynamic programming and to obtain (at least an **approximation** of) the overall optimum of the SOC problem?

---

<sup>1</sup>One can only say that  $\Lambda_t^{(k)}$  is **measurable** with respect to  $(W_0, \dots, W_t)$ .

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## Optimization problem

We recall the SOC problem under consideration.

$$\min_{\mathbf{U}, \mathbf{X}} \sum_{i=1}^N \left( \mathbb{E} \left( \sum_{t=0}^{T-1} L_t^i(\mathbf{X}_t^i, \mathbf{U}_t^i, \mathbf{W}_{t+1}) + K^i(\mathbf{X}_T^i) \right) \right), \quad (1a)$$

subject to **dynamics** constraints

$$\mathbf{X}_0^i = f_{-1}^i(\mathbf{W}_0), \quad (1b)$$

$$\mathbf{X}_{t+1}^i = f_t^i(\mathbf{X}_t^i, \mathbf{U}_t^i, \mathbf{W}_{t+1}), \quad (1c)$$

to **measurability** constraints

$$\mathbf{U}_t^i \preceq \sigma(\mathbf{W}_0, \dots, \mathbf{W}_t), \quad (1d)$$

and to **spatial** constraints

$$\sum_{i=1}^N \Theta_t^i(\mathbf{X}_t^i, \mathbf{U}_t^i) = 0. \quad \text{Coupling constraints} \quad (1e)$$

# Assumptions

## Assumption 1 (White noise)

Noises  $\mathbf{W}_0, \dots, \mathbf{W}_T$  are independent over time.

Note that we also assume **full noise observation**:

$$\mathbf{U}_t^i \preceq \sigma(\mathbf{W}_0, \dots, \mathbf{W}_t).$$

As a consequence of these assumptions, there is **no optimality loss** to seek the controls  $\mathbf{U}_t^i$  as a function of the state at time  $t$  rather than a function of the past noises:

$$\mathbf{U}_t^i = \gamma_i^t(\mathbf{X}_t^1, \dots, \mathbf{X}_t^N).$$

We are in the so-called **Markovian** case, and **DP** applies.

## Lagrangian formulation

We dualize the **coupling constraints** and obtain the **Lagrangian**

$$\mathcal{L}(\mathbf{X}, \mathbf{U}, \boldsymbol{\Lambda}) = \sum_{i=1}^N \left( \mathbb{E} \left( \sum_{t=0}^{T-1} L_t^i(\mathbf{X}_t^i, \mathbf{U}_t^i, \mathbf{W}_{t+1}) + K^i(\mathbf{X}_T^i) + \sum_{t=0}^{T-1} \boldsymbol{\Lambda}_t \cdot \Theta_t^i(\mathbf{X}_t^i, \mathbf{U}_t^i) \right) \right),$$

where the  $\boldsymbol{\Lambda}_t$ 's are  $\sigma(\mathbf{W}_0, \dots, \mathbf{W}_t)$  - measurable random variables.

We assume that a **saddle point** of  $\mathcal{L}$  exists,<sup>2</sup> so that

$$\min_{\mathbf{X}, \mathbf{X}} \max_{\boldsymbol{\Lambda}} \mathcal{L}(\mathbf{X}, \mathbf{U}, \boldsymbol{\Lambda}) = \max_{\boldsymbol{\Lambda}} \min_{\mathbf{U}, \mathbf{X}} \mathcal{L}(\mathbf{X}, \mathbf{U}, \boldsymbol{\Lambda}).$$

---

<sup>2</sup>Such an assumption requires going beyond the **Hilbert setting**...

## Uzawa algorithm

At **iteration**  $k$  of the algorithm,<sup>3</sup>

- 1 **Solve** subproblem  $i$ ,  $i = 1, \dots, N$ , with  $\Lambda^{(k)}$  fixed:

$$\min_{\mathbf{u}^i, \mathbf{x}^i} \mathbb{E} \left( \sum_{t=0}^{T-1} \left( L_t^i(\mathbf{x}_t^i, \mathbf{u}_t^i, \mathbf{w}_{t+1}) + \Lambda_t^{(k)} \cdot \Theta_t^i(\mathbf{x}_t^i, \mathbf{u}_t^i) \right) + K^i(\mathbf{x}_T^i) \right),$$

subject to

$$\mathbf{x}_{t+1}^i = f_t^i(\mathbf{x}_t^i, \mathbf{u}_t^i, \mathbf{w}_{t+1}),$$

$$\mathbf{u}_t^i \preceq \sigma(\mathbf{w}_0, \dots, \mathbf{w}_t).$$

The subproblem solution is denoted  $(\mathbf{u}^{i,(k+1)}, \mathbf{x}^{i,(k+1)})$ .

- 2 **Update** the multipliers  $\Lambda_t$ :

$$\Lambda_t^{(k+1)} = \Lambda_t^{(k)} + \rho_t \left( \sum_{i=1}^N \Theta_t^i(\mathbf{x}_t^{i,(k+1)}, \mathbf{u}_t^{i,(k+1)}) \right).$$

---

<sup>3</sup>The convergence of this algorithm is **not guaranteed** in this context...

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## Main idea of the approximation

As already seen,  $\Lambda_t^{(k)}$  depends on  $(W_0, \dots, W_t)$ , so that solving a subproblem by **DP** is as complex as solving the initial problem.

In order to overcome the difficulty, we **choose** at each time  $t$  a random variable  $Y_t$  which is measurable w.r.t. the past noises  $(W_0, \dots, W_t)$ . We call  $Y = (Y_0, \dots, Y_{T-1})$  the **information process** associated to the constraint.

The **core idea** is to replace the multiplier  $\Lambda_t^{(k)}$  at iteration  $k$  by its **conditional expectation w.r.t.  $Y_t$** , that is,  $\mathbb{E}(\Lambda_t^{(k)} | Y_t)$ .

This will lead to a “good” approximation if

$Y_t$  is **“sufficiently” correlated** to the random variable  $\Lambda_t$ .

*Note that we require that the information process is not influenced by controls.*

## Subproblem approximation

Following this idea, we **replace** subproblem  $i$  in Uzawa algorithm by:

$$\min_{\mathbf{U}^i, \mathbf{X}^i} \mathbb{E} \left( \sum_{t=0}^{T-1} \left( L_t^i(\mathbf{X}_t^i, \mathbf{U}_t^i, \mathbf{W}_{t+1}) + \mathbb{E}(\boldsymbol{\Lambda}_t^{(k)} \mid \mathbf{Y}_t) \cdot \Theta_t^i(\mathbf{X}_t^i, \mathbf{U}_t^i) \right) + K^i(\mathbf{X}_T^i) \right),$$

$$\begin{aligned} \text{subject to } \mathbf{X}_{t+1}^i &= f_t^i(\mathbf{X}_t^i, \mathbf{U}_t^i, \mathbf{W}_{t+1}), \\ \mathbf{U}_t^i &\preceq \sigma(\mathbf{W}_0, \dots, \mathbf{W}_t). \end{aligned}$$

The conditional expectation  $\mathbb{E}(\boldsymbol{\Lambda}_t^{(k)} \mid \mathbf{Y}_t)$  corresponds to a given function  $\mu_t$  of the variable  $\mathbf{Y}_t$ , so that subproblem  $i$  now involves the white noise process  $\mathbf{W}$  and the information process  $\mathbf{Y}$ . If the process  $\mathbf{Y}$  follows a **Markovian** dynamics, e.g.

$$\mathbf{Y}_{t+1} = h_t(\mathbf{Y}_t, \mathbf{W}_{t+1}),$$

then  $(\mathbf{X}_t^i, \mathbf{Y}_t)$  is a **valid state** for subproblem  $i$  and **DP** applies.



## Possible choices for the information process

- 1 **Perfect memory:**  $\mathbf{Y}_t = (\mathbf{W}_0, \dots, \mathbf{W}_t)$ .
  - $\mathbb{E}(\boldsymbol{\Lambda}_t^{(k)} \mid \mathbf{Y}_t) = \boldsymbol{\Lambda}_t^{(k)}$ : **no approximation!**
  - A **valid state** for each subproblem is  $(\mathbf{W}_0, \dots, \mathbf{W}_t)$ .
- 2 **Minimal information:**  $\mathbf{Y}_t \equiv \text{cste}$ .
  - $\boldsymbol{\Lambda}_t^{(k)}$  is approximated by its **expectation**  $\mathbb{E}(\boldsymbol{\Lambda}_t^{(k)})$ .
  - The information variable does not deliver online information.
  - A **valid state** for subproblem  $i$  is  $\mathbf{X}_t^i$ .
- 3 **Static information:**  $\mathbf{Y}_t = h_t(\mathbf{W}_t)$ .
  - A part of  $\mathbf{W}_t$  mostly “explains” the optimal multiplier.
  - A **valid state** for subproblem  $i$  is  $\mathbf{X}_t^i$ .
- 4 **Dynamic information:**  $\mathbf{Y}_{t+1} = h_t(\mathbf{Y}_t, \mathbf{W}_{t+1})$ .
  - A number of possibilities: mimicking the state of another unit, adding a hidden dynamics...
  - A **valid state** for subproblem  $i$  is  $(\mathbf{X}_t^i, \mathbf{Y}_t)$ .

## Dynamic programming equation

In the case of **dynamic information**, the **DP** equation associated to subproblem  $i$  writes:

$$V_T^i(x, y) = K^i(x),$$

$$V_t^i(x, y) = \min_u \mathbb{E} \left( \left( L_t^i(x, u, \mathbf{W}_{t+1}) \right. \right. \\ \left. \left. + \mathbb{E}(\boldsymbol{\Lambda}_t^{(k)} \mid \mathbf{Y}_t = y) \cdot \Theta_t^i(x, u) \right. \right. \\ \left. \left. + V_{t+1}^i(\mathbf{X}_{t+1}^i, \mathbf{Y}_{t+1}) \right) \right),$$

subject to the dynamics:

$$\mathbf{X}_{t+1}^i = f_t^i(x, u, \mathbf{W}_{t+1}),$$

$$\mathbf{Y}_{t+1} = h_t(y, \mathbf{W}_{t+1}).$$

## About the coordination: standard way

The task of coordination is performed thanks to scenarios.

- A set of **noise scenarios** is drawn once for all. **Trajectories** of the information process  $\mathbf{Y}$  are **simulated** along the scenarios.
- At iteration  $k$ , the **optimal trajectories** of the state process  $\mathbf{X}^{i,(k+1)}$  and of the control process  $\mathbf{U}^{i,(k+1)}$  are **simulated** along the noise scenarios, for all subsystems.

- The **dual multipliers** are **updated** along the noise scenarios according to the formula:

$$\boldsymbol{\Lambda}_t^{(k+1)} = \boldsymbol{\Lambda}_t^{(k)} + \rho_t \left( \sum_{i=1}^N \Theta_t^i(\mathbf{x}_t^{i,(k+1)}, \mathbf{u}_t^{i,(k+1)}) \right).$$

- The **conditional expectations**  $\mathbb{E}(\boldsymbol{\Lambda}_t^{(k+1)} | \mathbf{Y}_t)$  are obtained by **regression** of the trajectories of  $\boldsymbol{\Lambda}_t^{(k+1)}$  on those of  $\mathbf{Y}_t$ .

## About the coordination: information based way

II

One may perform the coordination by dealing with **functions** of  $\mathbf{Y}_t$ .

- Compute the **optimal trajectories** of the state process  $\mathbf{x}^{i,(k+1)}$  and of the control process  $\mathbf{u}^{i,(k+1)}$  along the noise scenarios.
- Compute the **conditional expectation** of the gradient:

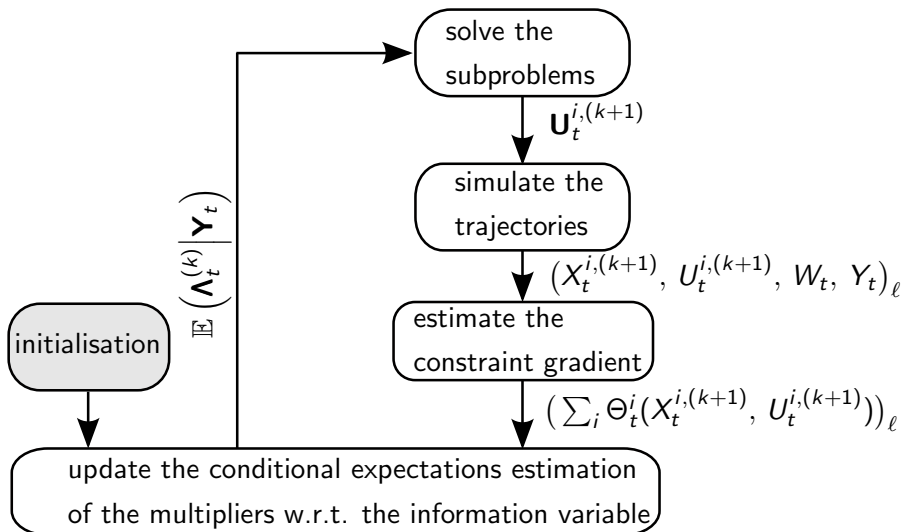
$$\mathbb{E} \left( \sum_{i=1}^N \Theta_t^i(\mathbf{x}_t^{i,(k+1)}, \mathbf{u}_t^{i,(k+1)}) \mid \mathbf{Y}_t \right).$$

- Update the **conditional expectation** of the multipliers:

$$\begin{aligned} \mathbb{E}(\boldsymbol{\Lambda}_t^{(k+1)} \mid \mathbf{Y}_t) &= \mathbb{E}(\boldsymbol{\Lambda}_t^{(k)} \mid \mathbf{Y}_t) \\ &+ \rho_t \mathbb{E} \left( \sum_{i=1}^N \Theta_t^i(\mathbf{x}_t^{i,(k+1)}, \mathbf{u}_t^{i,(k+1)}) \mid \mathbf{Y}_t \right). \end{aligned}$$

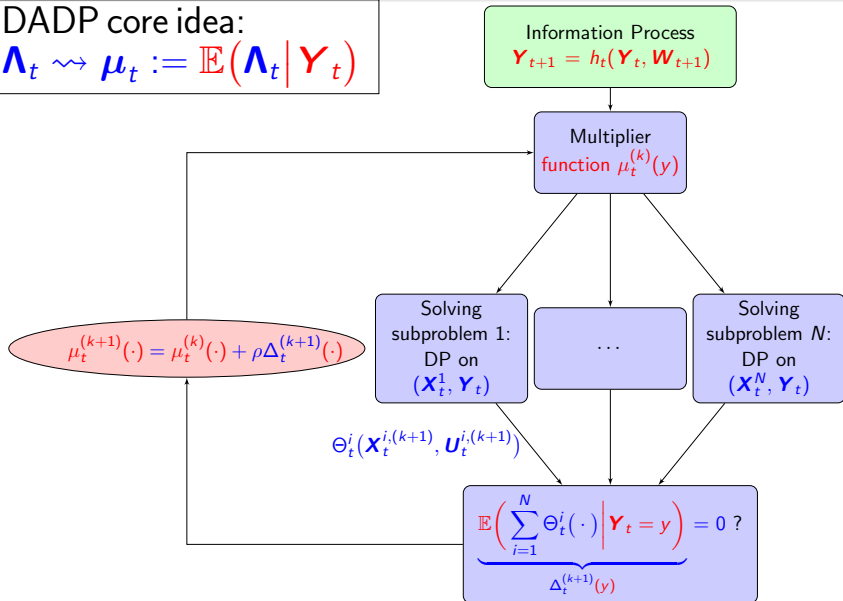
**Many numerical advantages if the support of  $\mathbf{Y}_t$  is finite.**

# DADP flowchart



DADP core idea:

$$\Lambda_t \rightsquigarrow \mu_t := \mathbb{E}(\Lambda_t | Y_t)$$



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## Interpretation of DADP

The **approximation** made on the **dual process** gives us a tractable way of computing strategies for the subsystems. Let us examine precisely the consequences in terms of **constraints**.

Consider a **relaxed** problem derived from (1):

$$\min_{\mathbf{u}, \mathbf{x}} \mathbb{E} \left( \sum_{i=1}^N \left( \sum_{t=0}^{T-1} L_t^i(\mathbf{x}_t^i, \mathbf{u}_t^i, \mathbf{w}_{t+1}) + K^i(\mathbf{x}_T^i) \right) \right), \quad (2a)$$

subject to the **modified coupling** constraints:

$$\mathbb{E} \left( \sum_{i=1}^N \Theta_t^i(\mathbf{x}_t^i, \mathbf{u}_t^i) \mid \mathbf{y}_t \right) = 0. \quad (2b)$$



# Interpretations of DADP



## Proposition 1

Suppose the Lagrangian associated with Problem (2) has a **saddle point**. Then the **DADP** algorithm can be interpreted as the **Uzawa algorithm applied** to Problem (2).

**Proof.** Since the duality term  $\mathbb{E}(\mathbb{E}(\Lambda_t^{(k)} | \mathbf{Y}_t) \cdot \Theta_t^i(\mathbf{X}_t^i, \mathbf{U}_t^i))$  which appears in the cost function of subproblem  $i$  in DADP can be written:

$$\mathbb{E}(\mathbb{E}(\Lambda_t^{(k)} | \mathbf{Y}_t) \cdot \Theta_t^i(\mathbf{X}_t^i, \mathbf{U}_t^i)) = \mathbb{E}(\Lambda_t^{(k)} \cdot \mathbb{E}(\Theta_t^i(\mathbf{X}_t^i, \mathbf{U}_t^i) | \mathbf{Y}_t)) ,$$

the global constraint **really** handled by DADP is:

$$\mathbb{E} \left( \sum_{i=1}^N \Theta_t^i(\mathbf{X}_t^i, \mathbf{U}_t^i) \mid \mathbf{Y}_t \right) = 0 . \quad \square$$

DADP thus consists in replacing an **almost-sure** constraint by its **conditional expectation** w.r.t. the **information variable**  $\mathbf{Y}_t$ .

## Interpretations of DADP



- DADP as an **approximation** of the optimal multiplier

$$\Lambda_t \rightsquigarrow \mathbb{E}(\Lambda_t \mid \mathbf{Y}_t) .$$

- DADP as a **decision-rule** approach for the dual problem

$$\max_{\Lambda} \min_{\mathbf{U}, \mathbf{X}} \mathcal{L}(\mathbf{X}, \mathbf{U}, \Lambda) \rightsquigarrow \max_{\Lambda_t \preceq \mathbf{Y}_t} \min_{\mathbf{U}, \mathbf{X}} \mathcal{L}(\mathbf{X}, \mathbf{U}, \lambda) .$$

- DADP as a **constraint relaxation** for the primal problem

$$\sum_{i=1}^N \Theta_t^i(\mathbf{x}_t^i, \mathbf{u}_t^i) = 0 \rightsquigarrow \mathbb{E}\left(\sum_{i=1}^N \Theta_t^i(\mathbf{x}_t^i, \mathbf{u}_t^i) \mid \mathbf{Y}_t\right) = 0 .$$

Thanks to the last interpretation, the optimal value given by DADP is a **guaranteed lower bound** for the original problem value.

## Summary

To summarize, **DADP** leads to solve the approximated problem

$$\min_{\mathbf{U}, \mathbf{X}} \mathbb{E} \left( \sum_{i=1}^N \sum_{t=0}^{T-1} \left( L_t^i(\mathbf{X}_t^i, \mathbf{U}_t^i, \mathbf{W}_t) + K^i(\mathbf{X}_T^i) \right) \right) \quad \text{s.t.} \quad \mathbb{E} \left( \sum_{i=1}^N \Theta_t^i(\mathbf{X}_t^i, \mathbf{U}_t^i) \mid \mathbf{Y}_t \right) = 0,$$

whereas the true problem is

$$\min_{\mathbf{U}, \mathbf{X}} \mathbb{E} \left( \sum_{i=1}^N \sum_{t=0}^{T-1} \left( L_t^i(\mathbf{X}_t^i, \mathbf{U}_t^i, \mathbf{W}_t) + K^i(\mathbf{X}_T^i) \right) \right) \quad \text{s.t.} \quad \sum_{i=1}^N \Theta_t^i(\mathbf{X}_t^i, \mathbf{U}_t^i) = 0.$$

### Questions:

- ★ What is the suitable theoretical framework of the algorithm?
  - ◇ Existence of a multiplier ?
  - ◇ Convergence of the algorithm ?
- ★ Does the approximate solution converge to the true solution?
- ★ How to obtain a feasible solution from the approximate solution?

## Some questions

### ★ What is the suitable theoretical framework of the algorithm?

The convergence of Uzawa's algorithm is granted provided that:

- the problem is posed in Hilbert spaces,
- and a saddle point exists.

It thus seems natural to place ourselves in a Hilbert space. But it is known (works by Rockafellar and Wets) that a saddle point doesn't exist in Hilbert spaces for such problems. . . (*See V. Leclère thesis.*)

### ★ Does the approximate solution converge to the true solution?

Epiconvergence results are available w.r.t. the information given by  $Y_t$ . But epiconvergence raises technical problems when addressed to stochastic optimization problems. (*See V. Leclère thesis.*)

### ★ How to obtain a feasible solution from the approximate solution?

Use an appropriate heuristic! (*See J.-C. Alais thesis.*)

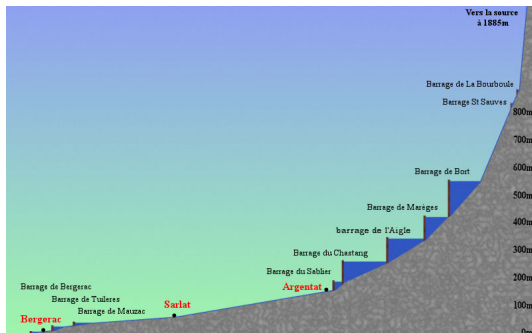
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# Some French hydro valleys



## Motivation

### Electricity production management for large hydro valleys

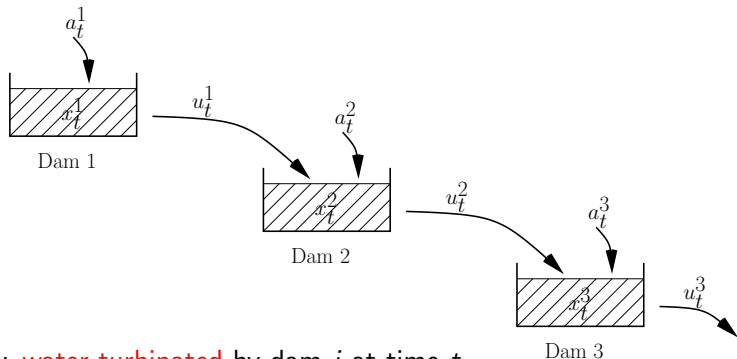


- *1 year time horizon:*  
compute each month  
the “values of water”  
(cost-to-go functions)
- *stochastic framework:*  
rain, market prices
- *large-scale valley:*  
4 dams and more

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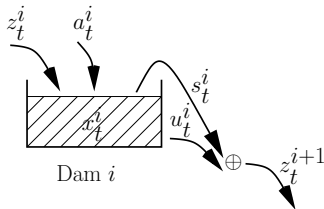
## Operating scheme



$u_t^i$  : water turbinated by dam  $i$  at time  $t$ ,  
 $x_t^i$  : water volume of dam  $i$  at time  $t$ ,  
 $a_t^i$  : water inflow at dam  $i$  at time  $t$ ,  
 $p_t^i$  : water price at dam  $i$  at time  $t$ ,

**Randomness:**  $w_t^i = (a_t^i, p_t^i)$  ,  $w_t = (w_t^1, \dots, w_t^N)$ .

# Dynamics and cost functions



## Dam dynamics

$$\begin{aligned}
 x_{t+1}^i &= f_t^i(x_t^i, u_t^i, w_t^i, z_t^i), \\
 &= x_t^i - u_t^i + a_t^i + z_t^i - s_t^i, \\
 z_t^{i+1} &\text{ being the outflow of dam } i: \\
 z_t^{i+1} &= g_t^i(x_t^i, u_t^i, w_t^i, z_t^i), \\
 &= u_t^i + \underbrace{\max\{0, x_t^i - u_t^i + a_t^i + z_t^i - \bar{x}^i\}}_{s_t^i}.
 \end{aligned}$$

We assume that  $u_t^i$  is chosen once  $w_t^i$  is observed (HD information structure), so that  $\underline{u}^i \leq u_t^i \leq \min\{\bar{u}^i, x_t^i + a_t^i + z_t^i - \underline{x}^i\}$ .

Gain at time  $t < T$ :  $L_t^i(x_t^i, u_t^i, w_t^i, z_t^i) = p_t^i u_t^i - \epsilon (u_t^i)^2$ .

Final gain at time  $T$ :  $K^i(x_T^i) = -a^i \min\{0, x_T^i - \hat{x}^i\}^2$ .

## Stochastic optimization problem

The **global optimization** problem reads:

$$\max_{(\mathbf{X}, \mathbf{U}, \mathbf{Z})} \mathbb{E} \left( \sum_{i=1}^N \left( \sum_{t=0}^{T-1} L_t^i(\mathbf{X}_t^i, \mathbf{U}_t^i, \mathbf{W}_t^i, \mathbf{Z}_t^i) + K^i(\mathbf{X}_T^i) \right) \right),$$

subject to:

$$\mathbf{X}_{t+1}^i = f_t^i(\mathbf{X}_t^i, \mathbf{U}_t^i, \mathbf{W}_t^i, \mathbf{Z}_t^i), \quad \forall i, \quad \forall t,$$

$$\mathbf{U}_t^i \preceq \sigma(\mathbf{W}_0, \dots, \mathbf{W}_t), \quad \forall i, \quad \forall t,$$

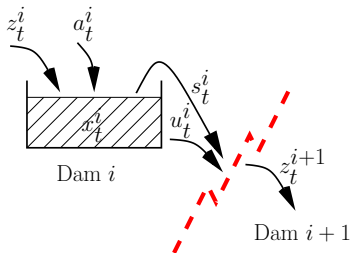
$$\mathbf{Z}_t^{i+1} = g_t^i(\mathbf{X}_t^i, \mathbf{U}_t^i, \mathbf{W}_t^i, \mathbf{Z}_t^i), \quad \forall i, \quad \forall t.$$

$\rightsquigarrow$  **Additive structure** (“cascade” model).

**Assumption.** Noises  $\mathbf{W}_0, \dots, \mathbf{W}_{T-1}$  are *independent over time*.

## Price decomposition

- Dualize the coupling constraints  $\mathbf{Z}_t^{i+1} = g_t^i(\mathbf{X}_t^i, \mathbf{U}_t^i, \mathbf{W}_t^i, \mathbf{Z}_t^i)$ . Denote by  $\Lambda_t^{i+1}$  the associated multiplier (random variable).
- Minimize the dual problem (using a gradient-like algorithm).



- At iteration  $k$ , the duality term:

$$\Lambda_t^{i+1,(k)} \cdot (\mathbf{Z}_t^{i+1} - g_t^i(\mathbf{X}_t^i, \mathbf{U}_t^i, \mathbf{W}_t^i, \mathbf{Z}_t^i)),$$

is the difference of two terms:

- $\Lambda_t^{i+1,(k)} \cdot \mathbf{Z}_t^{i+1} \rightsquigarrow$  dam  $i+1$ ,
  - $\Lambda_t^{i+1,(k)} \cdot g_t^i(\dots) \rightsquigarrow$  dam  $i$ .
- Dam by dam decomposition for the maximization in  $(\mathbf{X}, \mathbf{U}, \mathbf{Z})$  at  $\Lambda_t^{i+1,(k)}$  fixed.

## DADP implementation

DADP **approximation**:

- replace the constraint  $\mathbf{Z}_t^{i+1} - g_t^i(\mathbf{X}_t^i, \mathbf{U}_t^i, \mathbf{W}_t^i, \mathbf{Z}_t^i) = 0$  by its **conditional expectation** with respect to  $\mathbf{Y}_t^i$ :

$$\mathbb{E}(\mathbf{Z}_t^{i+1} - g_t^i(\mathbf{X}_t^i, \mathbf{U}_t^i, \mathbf{W}_t^i, \mathbf{Z}_t^i) \mid \mathbf{Y}_t^i) = 0,$$

- where  $(\mathbf{Y}_0^i, \dots, \mathbf{Y}_{T-1}^i)$  is a “well-chosen” **information process**.

The expression of the  $i$ -th subproblem becomes:

$$\begin{aligned} \max_{\mathbf{U}^i, \mathbf{Z}^i, \mathbf{X}^i} \mathbb{E} \left( \sum_{t=0}^{T-1} \left( L_t^i(\mathbf{X}_t^i, \mathbf{U}_t^i, \mathbf{W}_t^i, \mathbf{Z}_t^i) + \mathbb{E}(\boldsymbol{\Lambda}_t^{i,(k)} \mid \mathbf{Y}_t^{i-1}) \cdot \mathbf{Z}_t^i \right. \right. \\ \left. \left. - \mathbb{E}(\boldsymbol{\Lambda}_t^{i+1,(k)} \mid \mathbf{Y}_t^i) \cdot g_t^i(\mathbf{X}_t^i, \mathbf{U}_t^i, \mathbf{W}_t^i, \mathbf{Z}_t^i) \right) \right. \\ \left. + K^i(\mathbf{X}_T^i) \right). \end{aligned}$$

## A crude relaxation: $\mathbf{Y}_t^i \equiv \text{cste}$

- 1 The multipliers  $\boldsymbol{\Lambda}_t^{i,(k)}$  appear only in the subproblems by means of their expectations  $\mathbb{E}(\boldsymbol{\Lambda}_t^{i,(k)})$ , so that each subproblem involves the **1-dimensional** state variable  $\mathbf{X}_t^i$ .
- 2 For the gradient algorithm, the coordination task reduces to:

$$\mathbb{E}(\boldsymbol{\Lambda}_t^{i,(k+1)}) = \mathbb{E}(\boldsymbol{\Lambda}_t^{i,(k)}) - \rho_t \mathbb{E}\left(\mathbf{Z}_t^{i+1,(k)} - g_t^i(\mathbf{X}_t^{i,(k)}, \mathbf{U}_t^{i,(k)}, \mathbf{W}_t^i, \mathbf{Z}_t^{i,(k)})\right).$$

- 3 The constraints taken into account by DADP are in fact

$$\mathbb{E}\left(\mathbf{Z}_t^{i+1} - g_t^i(\mathbf{X}_t^i, \mathbf{U}_t^i, \mathbf{W}_t^i, \mathbf{Z}_t^i)\right) = 0.$$

The solutions do not satisfy the initial almost sure constraints: need to use a **heuristic method** to regain admissibility.

## How to regain admissible policies?

We have computed  $N$  local Bellman functions  $V_t^i$  at each time  $t$ , each depending on a single state variable  $x^i$ , whereas we want a unique global Bellman function  $V_t$  depending on  $(x^1, \dots, x^N)$ .

**Value function approximation:** form the following functions:

$$\widehat{V}_t(x^1, \dots, x^N) = \sum_{i=1}^N V_t^i(x^i).$$

For any  $(x_t, w_t)$  at time  $t$ , solve the one-step DP problem:

$$\max_{u, z} \sum_{i=1}^N L_t^i(x_t^i, u^i, w_t^i, z^i) + \widehat{V}_{t+1}(x_{t+1}^1, \dots, x_{t+1}^N),$$

$$\text{s.t. } x_{t+1}^i = f_t^i(x_t^i, u^i, w_t^i, z^i) \quad \text{and} \quad z^{i+1} = g_t^i(x_t^i, u^i, w_t^i, z^i).$$

$\rightsquigarrow$  control value  $u_t^\# = (u_t^{1,\#}, \dots, u_t^{N,\#})$  to be used at  $(x_t, w_t)$ .

# Full optimization and simulation process

## • Optimization

- Apply DADP and compute the cost-to-go functions  $V_t^i$ .
- Form the approximated global Bellman functions  $\hat{V}_t$ .

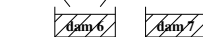
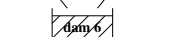
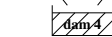
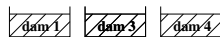
## • Simulation

- Draw a large number of noise scenarios.
- Compute the control values along each scenario by solving the one-step DP problems involving the  $\hat{V}_t$ 's, thus satisfying all the constraints of the initial problem:
  - ↪ payoff value for each scenario,
  - ↪ state and control trajectories.
- Evaluate the quality of the solution: mean payoff, . . .



- 1 Examples and mathematical background
  - Interconnected systems
  - Optimization background
  - Standard decomposition methods
- 2 About decomposition in stochastic optimization
  - Couplings in stochastic optimization
  - Dynamic programming and decomposition
- 3 Dual approximate dynamic programming (DADP)
  - Problem formulation and price decomposition
  - Subproblems resolution and coordination
  - What has been really done?
- 4 Hydro valleys management problem
  - DADP implementation for hydro valleys
  - Numerical results for different valleys

## Academic case studies of increasing complexity



Discretization

 $T \rightsquigarrow 12$  $X \rightsquigarrow 41$  $U \rightsquigarrow 6$  $W \rightsquigarrow 10$ 

4-Dams

6-Dams

8-Dams

10-Dams

## Optimal values and computational times

Valley	4-Dams	6-Dams	8-Dams	10-Dams
DP CPU time	$1.6 \cdot 10^3$ '	$\sim 10^8$ '	$\sim \infty$	$\sim \infty$
DP value	3743	N.A.	N.A.	N.A.
SDDP <sub>c</sub> value	3742	7026	11834	17069
SDDP <sub>c</sub> CPU time	5'	7'	9'	50'
Valley	4-Dams	6-Dams	8-Dams	10-Dams

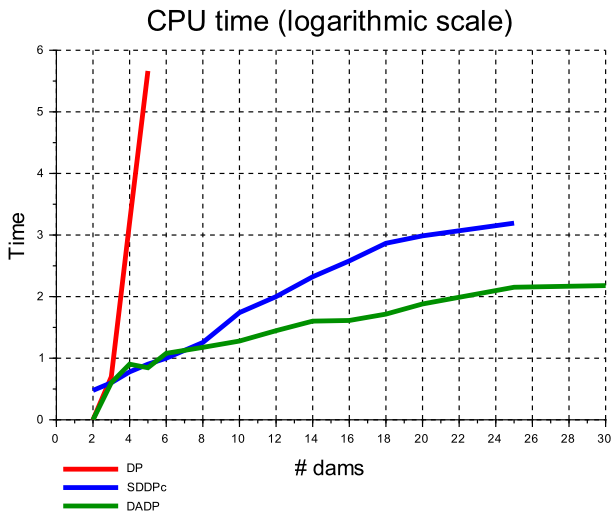
Table: Results obtained by DP and SDDP<sub>c</sub><sup>4</sup>

Valley	4-Dams	6-Dams	8-Dams	10-Dams
DADP CPU time	7'	12'	17'	24'
DADP value	3667	6816	11573	16760
Gap with SDDP <sub>c</sub>	-2.0%	-3.0%	-2.2%	-1.8%

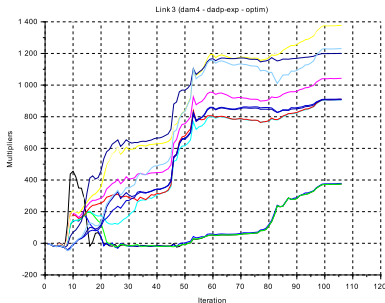
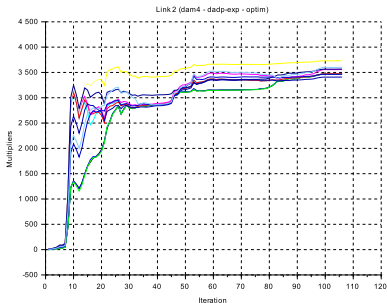
Table: Results obtained by DADP "Expectation"

<sup>4</sup>The SDDP method will be explained in detail next Monday.

## CPU time summary

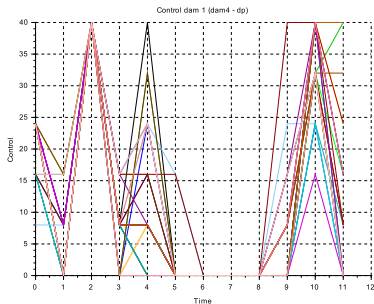


## 4-Dams in detail: DADP convergence

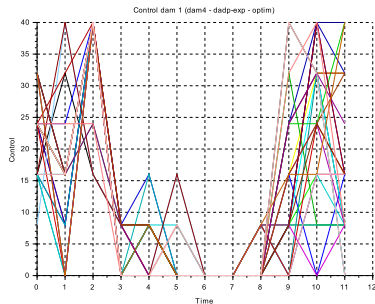


Multipliers convergence (dam1 $\leftrightarrow$ dam2 and dam2 $\leftrightarrow$ dam3)

## 4-Dams in detail: control trajectories

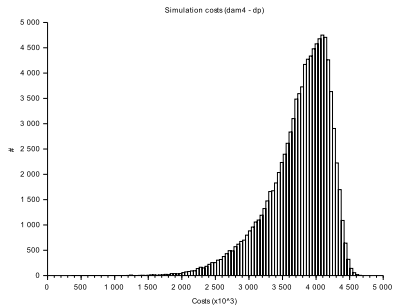


DP: dam 1 trajectories

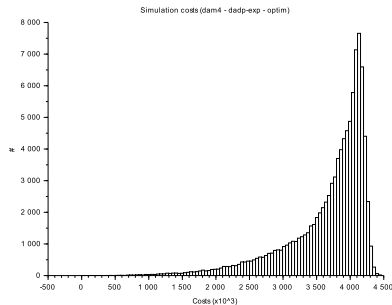


DADP: dam 1 trajectories

## 4-Dams in detail: payoff distributions

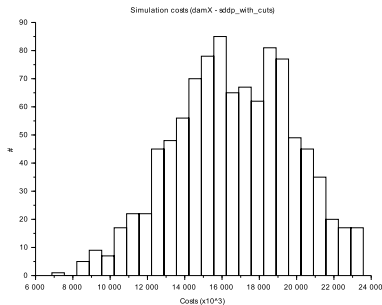


DP payoff distribution

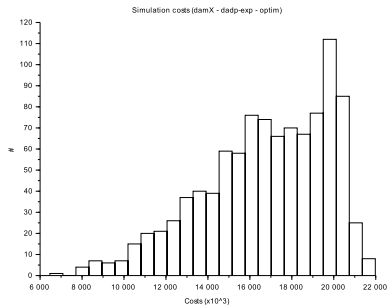


DADP payoff distribution

## 10-Dams in detail: payoff distribution



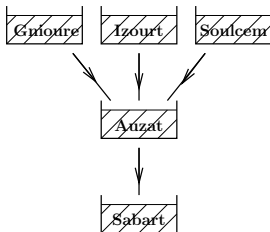
SDDP payoff distribution



DADP payoff distribution



## Two “true” valleys

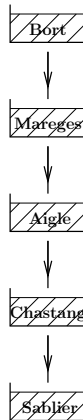


*Discretization*

$T \rightsquigarrow 12$ ,  $W \rightsquigarrow 10$

realistic grids for  $U$  and  $X$

Vicdessos



Dordogne

## Results

Valley	Vicdessos	Dordogne
SDDP <sub>c</sub> CPU time	9'	17'
SDDP <sub>c</sub> value	2244	22145

Table: Results obtained by SDDP<sub>c</sub>

Valley	Vicdessos	Dordogne
DADP CPU time	9'	210'
DADP value	2237	21652
Gap with SDDP <sub>c</sub>	-0.3%	-2.2%

Table: Results obtained by DADP "Expectation"

## Conclusions and perspectives

### Conclusions for DADP

- Fast numerical convergence of the method.
- Near-optimal results even when using a “crude” relaxation.
- Method that can be used for very large valleys

### General perspectives

- Apply to more complex topologies (smart grids).
- Connection with other decomposition methods.
- Theoretical study.



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