

Dynamic Programming

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CIRM, Luminy
November 2019

Outline of the presentation

Introduction to Dynamic Programming
The All-pairs Shortest Path Problem

The One Step Newsvendor Problem

The T -Step Newsvendor Problem

The Dynamic Programming Equation

Exercise

Extensions

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The All-Pairs Shortest Path Problem

- ▶ (V, E) a given oriented graph, nodes V , edges $E \subset V \times V$.
- ▶ P : sequences (v_1, \dots, v_n) , s.t $v_i \in V$, $(v_i, v_{i+1}) \in E$. given $\ell : E \rightarrow \mathbb{R} \cup \{+\infty\}$ we extend $\ell : P \rightarrow \mathbb{R} \cup \{+\infty\}$ by

$$\ell((v_1, \dots, v_n)) := \sum_{i=1}^{n-1} \ell((v_i, v_{i+1})) \quad (\text{path length})$$

- ▶ Shortest path from v_i to v_j :

$$sp(v_i, v_j) \in \arg \min \{ \ell(p) \mid p = (w_1, \dots, w_n) \in P, w_1 = v_i, w_n = v_j \}$$

Find $sp(v, w)$ for all $(v, w) \in V^2$ is called the all-pairs shortest path problem.

- ▶ Assumption : there does not exist cycles with negative length

$$p = (w_1, \dots, w_n) \in P, w_1 = w_n \text{ then } \ell(p) \geq 0$$

Matrix Notations

- ▶ To simplify the notation we consider that $V = \{1, \dots, n\}$
- ▶ $\ell : E \rightarrow \mathbb{R} \cup \{+\infty\}$ is represented by a matrix L

$$L_{i,j} = \begin{cases} 0 & \text{when } i = j, \\ +\infty & \text{when } (i,j) \notin E \\ \ell(i,j) & \text{when } (i,j) \in E \end{cases}$$

- ▶ We define $(D^{(k)})_{k \in \mathbb{N}}$, a sequences of matrices given by

$$D_{i,j}^{(k)} = \begin{cases} 0 & \text{when } i = j, \\ lsp_k(i,j) \end{cases}$$

$$lsp_k(i,j) = \min \{ \ell(p) \mid p = (i, \dots, j) \in P, |p| \leq k + 1 \}$$

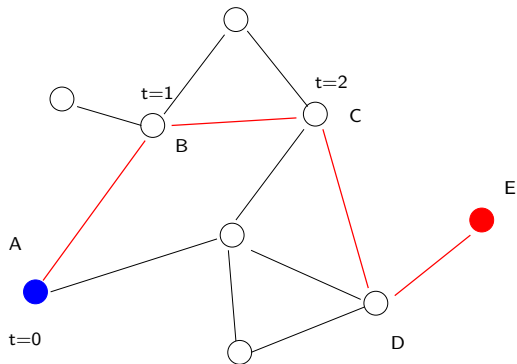
Dynamic Programming Approach

- ▶ We have that $D_{i,j}^{(1)} = L_{i,j}$
- ▶ The no-negative-cycle assumption implies that $D_{i,j}^{(k)} = \ell(sp(i,j))$ when $k \geq n - 1$ ($n = |V|$).
- ▶ We thus have to compute $D^{(1)}, \dots, D^{(n-1)}$
- ▶ We have introduced a family of problems, the original problem being one of them
- ▶ We compute **value functions** (value of optimal path lengths in set of constrained paths) instead of optimal paths.

A recursion formula giving the value function of problem k given the value of problem $k - 1$ will be obtained through an optimality principle.

Optimality Principle

A sub-path of an optimal path is an optimal path
If the path $A \rightarrow B \rightarrow C \rightarrow D \rightarrow E$ is a shortest path from A to E
Then, the path $C \rightarrow D \rightarrow E$ is a shortest path from C to E .



Optimality Principle at Work

- ▶ Consider $D_{i,j}^{(m)}$ the value of the shortest path going from i to j with at most $m + 1$ nodes.
- ▶ If the shortest path have at most m nodes then
 $D_{i,j}^{(m)} = D_{i,j}^{(m-1)}$
- ▶ If the shortest path have exactly $m + 1$ nodes : then it is composed of a path from i to k with m nodes and an edge (k, j) whose length is $L_{k,j}$.
- ▶ **Optimality principle** : The path from i to k must be optimal in the set of path with at most m nodes. Thus,

$$D_{i,j}^{(m)} = D_{i,k}^{(m-1)} + L_{k,j}$$

Gathering all the cases

$$D_{i,j}^{(m)} = \min \left(D_{i,j}^{(m-1)}, \min_{k \neq j} \left(D_{i,k}^{(m-1)} + L_{k,j} \right) \right) = \min_k \left(D_{i,k}^{(m-1)} + L_{k,j} \right)$$

Find the optimal path using $D^{(1)}, \dots, D^{(n-1)}$

- ▶ Compute the shortest path from i to j in p^\sharp :
 1. if $D_{i,j}^{(n-1)} = \infty$ then stop (there's no path from i to j)
 2. Fix $last = j$, set $q = 1$, and $p^\sharp = ()$
 3. If $n - q = 1$ then $p^\sharp = (i, p^\sharp)$ and go to end 7
 4. let $k^\sharp \in \arg \min_k \left(D_{s,k}^{(n-q)} + L_{k,last} \right)$.
 5. If $k^\sharp \neq last$ then $p^\sharp = (k^\sharp, p^\sharp)$.
 6. $q = q + 1$, $last = k^\sharp$ and iterate at 3.
 7. p^\sharp gives the optimal path

If $k^\sharp \in \arg \min_k \left(D_{i,k}^{(p-1)} + L_{k,j} \right)$

$$\underbrace{i \rightsquigarrow j}_{p\text{-shortest path}} = \underbrace{i \rightsquigarrow k^\sharp}_{p-1 \text{ shortest path}} \mapsto j$$

Recursive functions

- ▶ $Fib(n) = (n \leq 1)?1 : Fib(n-1) + Fib(n-2)$.

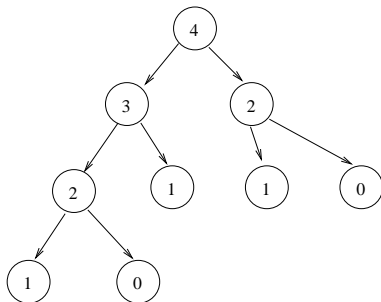


FIGURE – Fibonacci

Recursion

- ▶ Exponential complexity of the naive recursive algorithm.
- ▶ Solution 1 : iterative computation from $(Fib(0), Fib(1))$
- ▶ Solution 2 : use recursion but keep track of already computed values (memoization).

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The One Day Newsvendor Problem



- ▶ Each morning, The newsvendor orders a number of newspaper $u \in \mathbb{U} = \{0, 1, \dots\}$ at unit price $c > 0$.
- ▶ The demand is **incertain** $w \in \mathbb{W} = \{0, 1, \dots\}$
 - ▶ If at the end of the day, if it remains unsold newspapers, the incurred cost is

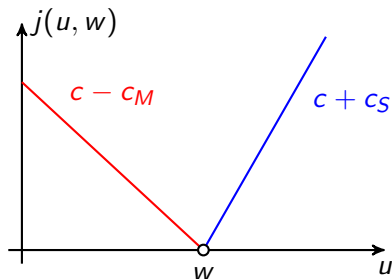
$$c_S(u - w)_+ = c_S \underbrace{\max(u - w, 0)}_{\text{unsold}} \quad \text{with } c > -c_S$$

- ▶ If during the day facing the demand was not possible, the incurred cost for unsatisfied demand is

$$c_M(w - u)_+ = c_M \underbrace{\max(w - u, 0)}_{\text{missing}}$$

Uncertain Demand Induces **Uncertain Cost**

$$j(u, w) = \underbrace{cu}_{\text{purchases}} + \underbrace{c_S(u - w)_+}_{\text{unsold}} + \underbrace{c_M(w - u)_+}_{\text{missing}}$$



$\text{Argmin}_{u \in \mathbb{U}} j(u, w) = \{w\}$: **unknown quantity w !**

Attitude Toward Risk

The choice of the cost to minimize depends on the risk attitude of the newsvendor. We will use to aggregate random values through expectation

$$\min_{u \in \mathbb{U}} J(u) \quad \text{with } J(u) := \underbrace{\mathbb{E}_w[j(u, w)]}_{\text{Expectation}}$$

A pessimistic newsvendor could decide to use the worse case among all the possible values of the demand

$$\min_{u \in \mathbb{U}} J(u) \quad \text{with } J(u) := \underbrace{\max_{w \in \mathbb{W}} j(u, w)}_{\text{Worse case}}$$

Expected cost

- ▶ The demand, \mathbf{W} , is a random variable. The newsvendor knows the **distribution** $\mathbb{P}_{\mathbf{W}}$, that is the law of \mathbf{W} .
- ▶ The cost to be minimized is

$$J(u) = \mathbb{E}_{\mathbf{W}} [cu + c_S(u - \mathbf{W})_+ + c_M(\mathbf{W} - u)_+]$$

- ▶ The newsvendor problem is to find

$$u^* \stackrel{?}{\in} \underset{u \in \mathbb{U}}{\text{Argmin}} J(u).$$

The newsvendor takes a deterministic decision facing a future random variable. He knows the law of the random variable but not the realization (Decision-Hazard framework).

Optimal Control

The optimal control u^* when $u \in \mathbb{R}_+$:

- ▶ If $c_M \leq c$ then J is a non-decreasing function and the optimal value is not to order $u^* = 0$
- ▶ If $c_M > c$ then the minimum of J is reached at u^* :

$$u^* = \inf\{z \in \mathbb{R} \mid F(z) \geq (c_M - c)/(c_M + c_S)\}$$

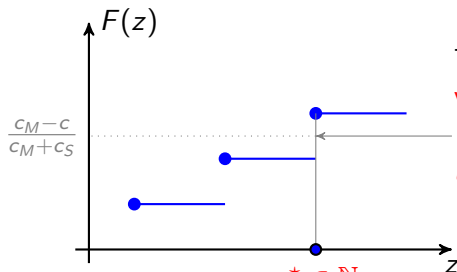
where F is the cumulative distribution function of \mathbf{W} ,
 $F(z) = \mathbb{P}(\mathbf{W} \leq z)$.

- ▶ If \mathbf{W} is a discrete random variable taking values in \mathbb{N} then the optimal value of the relaxed problem u^* is in \mathbb{N}

$$\begin{aligned} J(u) &= c_M \mathbb{E}_{\mathbf{W}}[\mathbf{W}] + (c - c_M)u + (c_M + c_S) \mathbb{E}_{\mathbf{W}}[(u - \mathbf{W})_+] \\ &= c_M \mathbb{E}_{\mathbf{W}}[\mathbf{W}] + (c - c_M)u + (c_M + c_S) \int_0^u F(z) dz \end{aligned}$$

J is continuous and coercive. When $c_M > c$ it is non-increasing then non-decreasing with a minimum at u^* .

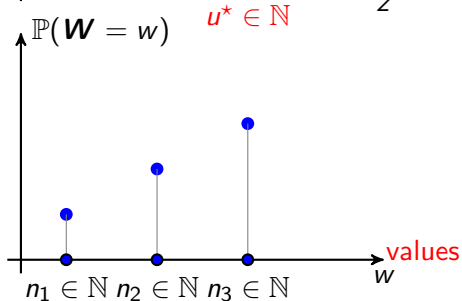
Optimal Control (II)



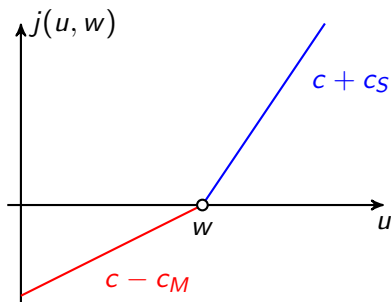
The optimal control takes **integer values**

$$u^* = \inf \left\{ z \in \mathbb{R} \mid F(z) \geq \frac{c_M - c}{c_M + c_S} \right\},$$

If the demand takes **integer**



The case $c_M \leq c$

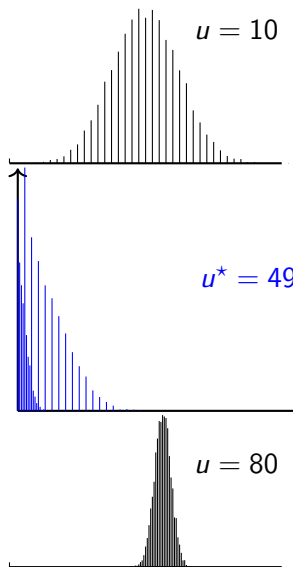


We note that

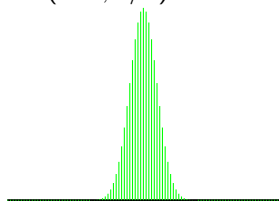
$$\min_{u \in \mathbb{U}} J(u) \geq \min_{u = \phi(\mathbf{W})} J(u) = \mathbb{E}_{\mathbf{W}} [\min_{u \in \mathbb{U}} j(u, \mathbf{W})]$$

$\min_{u \in \mathbb{U}} j(u, w)$ is attained by $u = 0$ for all w and gives an admissible command for the problem $\min_{u \in \mathbb{U}} j(u, w)$.

Costs Distribution $j(u, W)$



Distribution of the demand
 $\text{bin}(100, 1/2)$.



Initial Stock and Fixed Cost for Buying Newspapers

We change the model :

- ▶ We start with an initial stock $x \in \mathbb{Z}$
- ▶ When increasing stock a fixed cost c_F occurs

$$\begin{aligned}\tilde{J}(u) &= \mathbb{E}_{\mathbf{W}}[c_F \mathbb{I}_{\{u > 0\}} + cu + c_s(x + u - \mathbf{W})_+ + c_M(\mathbf{W} - x - u)_+] \\ &= c_F \mathbb{I}_{\{u > 0\}} - cx + \mathbb{E}_{\mathbf{W}}[j(u + x, \mathbf{W})] \\ &= c_F \mathbb{I}_{\{u > 0\}} + J(u + x) - cx\end{aligned}$$

- ▶ The optimal control $u^*(x)$ (for $c_M > c$) is a function of x
- ▶ The optimal control depends on two bounds (s, S) :

$$u^*(x) = (S - x) \mathbb{I}_{\{x \leq s\}}$$

- ▶ If the stock is smaller than s , buy newspapers to bring the stock to S .
- ▶ If the stock is bigger than s , do not buy newspapers
- ▶ The value of S is given by $\text{Argmin}_{u \in \mathbb{U}} J(u)$.

Initial Stock and Fixed Cost for Buying Newspapers

Let S given by $\{S\} = \text{Argmin}_{u \in \mathbb{U}} J(u)$.

- ▶ If $x \geq S$, not buying is optimal since $J(\cdot) \nearrow$ and $c_F \geq 0$:

$$\tilde{J}(0) = J(x) - cx \leq J(x + u) - cx + c_F = \tilde{J}(u)$$

- ▶ If $x \leq S$, since $J(\cdot)$ is minimal for S :

- ▶ If buying, we need to order $u = S - x$ whose cost is

$$\tilde{J}(u) = J(S) - cx + c_F$$

- ▶ If not buying the cost is $J(x) - cx$

The solution for minimizing the costs is to fill the stock up to S if $J(x) \geq c_F + J(S)$ and do nothing otherwise. Noting that

$$\{x \mid J(x) \geq c_F + J(S)\} = \{x \mid x \leq s\}$$

where s is given by

$$s := \sup \{z \in (-\infty, S) \mid J(z) \geq c_F + J(S)\}$$

Initial Stock and Fixed Cost for Buying Newspapers

Let $\mathbf{X}_0 = x$ and $\mathbf{X}_1 = f(\mathbf{X}_0, u)$ with $f(x, u) := x + u - w$ et

$$\tilde{j}(u, x_1) := c_F \mathbb{I}_{\{u > 0\}} + cu + c_S(x_1)_+ + c_M(-x_1)_+$$

The newsvendor problem is

$$\begin{aligned} \min_{u \in \mathcal{U}, \mathbf{X}_1, \mathbf{X}_0} \mathbb{E}_{\mathbf{W}}[\tilde{j}(u, \mathbf{X}_1)] \\ \mathbf{X}_0 = x \quad \mathbf{X}_1 = f(\mathbf{X}_0, u, \mathbf{W}) \end{aligned}$$

With a non-anticipative constraint

$$\mathcal{U} = \{\mathbf{U} : \Omega \rightarrow \mathbb{N} \mid \mathbf{U}(\omega) = \phi(\mathbf{X}_0(\omega))\}$$

Note that $u^\#(x)$ is obtained by

$$u^\#(x) \in \underset{u \in \mathbb{N}}{\text{Argmin}} h(u, x) \text{ with } h(u, x) = \mathbb{E}_{\mathbf{W}}[\bar{j}(u, x, \mathbf{W})]$$

with $\bar{j}(u, x, \mathbf{W}) = \tilde{j}(u, f(u, x, w))$.

Initial Stock and Fixed Cost for Buying Newspapers

- ▶ Suppose now that \mathbf{X}_0 is a given random variable
- ▶ The newsvendor problem becomes

$$\min_{\mathbf{U} \in \mathcal{U}} \mathbb{E} \left[\bar{j}(\mathbf{U}, \mathbf{X}_0, \mathbf{W}) \right]$$

Where $\mathcal{U} = \{ \mathbf{U} : \Omega \rightarrow \mathbb{N} \mid \mathbf{U}(\omega) = \phi(\mathbf{X}_0(\omega)) \}$

- ▶ Suppose that \mathbf{X}_0 and \mathbf{W} are independent r.v.
 - ▶ $h(\gamma(\mathbf{X}_0), \mathbf{X}_0) = \mathbb{E} \left[\bar{j}(\gamma(\mathbf{X}_0), \mathbf{X}_0, \mathbf{W}) \mid \mathbf{X}_0 \right]$
 - ▶ If $h(u^\sharp(\mathbf{X}_0), \mathbf{X}_0) \leq h(\gamma(\mathbf{X}_0), \mathbf{X}_0)$ for all γ then $\mathbb{E} \left[\bar{j}(u^\sharp(\mathbf{X}_0), \mathbf{X}_0, \mathbf{W}) \mid \mathbf{X}_0 \right] \leq \mathbb{E} \left[\bar{j}(\gamma(\mathbf{X}_0), \mathbf{X}_0, \mathbf{W}) \mid \mathbf{X}_0 \right]$
 - ▶ If $h(u^\sharp(\mathbf{X}_0), \mathbf{X}_0) \leq h(\gamma(\mathbf{X}_0), \mathbf{X}_0)$ for all γ then $\mathbb{E} \left[\bar{j}(u^\sharp(\mathbf{X}_0), \mathbf{X}_0, \mathbf{W}) \right] \leq \mathbb{E} \left[\bar{j}(\gamma(\mathbf{X}_0), \mathbf{X}_0, \mathbf{W}) \right]$
- ▶ The optimal control is $\mathbf{U}^\sharp = u^\sharp(\mathbf{X}_0)$
- ▶ We have

$$\min_{\mathbf{U} \in \mathcal{U}} \mathbb{E} \left[\bar{j}(\mathbf{U}, \mathbf{X}_0, \mathbf{W}) \right] = \mathbb{E} \left[\min_{u \in \mathbb{N}} \mathbb{E} \left[\bar{j}(u, \mathbf{X}_0, \mathbf{W}) \mid \mathbf{X}_0 \right] \right]$$

Dynamics as a Markov Chain

The two stocks \mathbf{X}_0 and \mathbf{X}_1 can be seen as two consecutive states of a controlled Markov chain.

- ▶ Assume that $u \in \mathbb{N}$ is fixed, the transition matrix is

$$P_{x_0, x_1}^u = \mathbb{P}(\mathbf{X}_1 = x_1 | \mathbf{X}_0 = x_0)$$

$$P_{x_0, x_1}^u = \begin{cases} \mathbb{P}(\mathbf{W} = w_0) & \text{if } x_1 = x_0 + u - w_0 \\ 0 & \text{if not} \end{cases}$$

- ▶ Assume that \mathbf{U} is chosen as a function of X_0 , $U = \phi(\mathbf{X}_0)$ then

$$P_{x_0, x_1}^\phi = \begin{cases} \mathbb{P}(\mathbf{W} = w_0) & \text{if } x_1 = x_0 + \phi(x_0) - w_0 \\ 0 & \text{if not} \end{cases}$$

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From One Stage To Finite Horizon Problem

The one stage problem with initial stock

$$\begin{aligned} \min_{u \in \mathcal{U}, \mathbf{X}_1, \mathbf{X}_0} \mathbb{E}_{\mathbf{W}}[\tilde{j}(u, \mathbf{X}_1)] \\ \text{s.t. } \mathbf{X}_0 = x \quad \mathbf{X}_1 = f(\mathbf{X}_0, u, \mathbf{W}_1) \end{aligned}$$

with

$$\begin{aligned} f(x, u, w) &:= x + u - w \\ \tilde{j}(u, x) &:= c_F \mathbb{I}_{\{u > 0\}} + cu + \alpha (c_S(x)_+ + c_M(-x)_+) \end{aligned}$$

The stock \mathbf{X}_t can be positive (physical stock) or negative (the opposite of missing newspapers)

The law of the demand, \mathbf{W} , is known (finite expectation)

Finite Horizon problem

The newsvendor minimizes the costs over a period T

$$\min_{\mathbf{u} \in \mathcal{U}, \mathbf{x}} \mathbb{E}_{\mathbf{w}} \left[\sum_{t=0}^{T-1} \alpha^t \tilde{J}(\mathbf{u}_t, \mathbf{x}_{t+1}) \right]$$
$$\mathbf{x}_0 = x \quad \mathbf{x}_{t+1} = f(\mathbf{x}_t, \mathbf{u}_t, \mathbf{w}_{t+1})$$

- ▶ $(\mathbf{x}_0, \mathbf{w}_1, \mathbf{w}_2, \dots, \mathbf{w}_T)$ are independent
- ▶ $\alpha \in (0, 1]$ is an actualization rate.

Canonical Form

The newsvendor minimizes the costs over a T period of time

$$\min_{\mathbf{U} \in \mathcal{U}, \mathbf{X}} \mathbb{E}_{\mathbf{W}} \left[\sum_{t=0}^{T-1} \alpha^t c_t(\mathbf{U}_t, \mathbf{X}_t) + \alpha^T K(\mathbf{X}_T) \right]$$
$$\mathbf{X}_0 = x \quad \mathbf{X}_{t+1} = f(\mathbf{X}_t, \mathbf{U}_t, \mathbf{W}_{t+1})$$

with

$$c_t(u, x) := c_F \mathbb{I}_{\{u > 0\}} + cu + c_S(x)_+ + c_M(-x)_+$$

$$c_0(u, x) := c_F \mathbb{I}_{\{u > 0\}} + cu$$

$$K(x) := c_S(x)_+ + c_M(-x)_+$$

Canonical Form (II)

$$\begin{aligned}\tilde{J}(u_0, x_1) + \alpha \tilde{J}(u_1, x_2) &= \overbrace{c_F \mathbb{I}_{\{u_0 > 0\}} + cU_0 + \alpha (c_S(x_1)_+ + c_M(-x_1)_+)}^{\tilde{J}(u_0, x_1)} \\ &+ \alpha \overbrace{(c_F \mathbb{I}_{\{u_1 > 0\}} + cU_1 + \alpha (c_S(x_2)_+ + c_M(-x_2)_+))}^{\tilde{J}(u_1, x_2)} \\ &= \overbrace{c_F \mathbb{I}_{\{u_0 > 0\}} + cU_0}^{c_0(u_0, x_0)} \\ &+ \alpha \overbrace{(c_F \mathbb{I}_{\{u_1 > 0\}} + cU_1 + c_S(x_1)_+ + c_M(-x_1)_+)}^{c_1(u_1, x_1)} \\ &+ \alpha^2 \overbrace{(c_S(x_2)_+ + c_M(-x_2)_+)}^{K(x_2)} \\ &= c_0(u_0, x_0) + \alpha c_1(u_1, x_1) + \alpha^2 K(x_2)\end{aligned}$$

Non Anticipativity

- ▶ The newsvendor collects over time the demand of each day.
 - ▶ At time t , he knows $(\mathbf{W}_1, \dots, \mathbf{W}_t)$ and \mathbf{X}_0 and can use this information to compute \mathbf{U}_t . He could also collect the past controls.
- ▶ Under the independence assumption of the r.v $(\mathbf{X}_0, \mathbf{W}_1, \dots, \mathbf{W}_t)$ the optimal control at time t only depends of the stock \mathbf{X}_t .

Policies

$$\min_{\mathbf{U} \in \mathcal{U}} J(\mathbf{U}) \text{ with } J(\mathbf{U}) = \mathbb{E}_{\mathbf{W}} \left[\sum_{t=0}^{T-1} \alpha^t c_t(\mathbf{U}_t, \mathbf{X}_t) + \alpha^T K(\mathbf{X}_T) \right]$$
$$\mathbf{X}_0 = \mathbf{x} \quad \mathbf{X}_{t+1} = f(\mathbf{X}_t, \mathbf{U}_t, \mathbf{W}_{t+1})$$

1. $\min_{\mathbf{U} \in \mathcal{U}_H} J(\mathbf{U})$: \mathcal{U}_H space of history dependent controls depending on past states, past controls, past noises
2. $\leq \min_{\mathbf{U} \in \mathcal{U}_M} J(\mathbf{U})$: \mathcal{U}_M space of markovian controls depending on current state
3. $\leq \min_{\mathbf{U} \in \mathcal{U}_{OL}} J(\mathbf{U})$: \mathcal{U}_{OL} space of open-loop controls deterministic (constant random variables)

We have that $\min_{\mathbf{U} \in \mathcal{U}_H} J(\mathbf{U}) = \min_{\mathbf{U} \in \mathcal{U}_M} J(\mathbf{U})$

State Feedback Versus Open-Loop Feedback

$$\min_{\mathbf{u}_1, \mathbf{u}_2} \mathbb{E} \left[\mathbf{x}_1^2 + \mathbf{x}_2^2 \right]$$

$$\text{s.t.} \quad \mathbf{x}_2 = \mathbf{x}_1 - \mathbf{u}_1 + \mathbf{w}_2, \quad \mathbf{x}_1 = \mathbf{x}_0 - \mathbf{u}_0 + \mathbf{w}_1, \quad \mathbf{x}_0 = 0$$
$$\mathbf{w}_0, \mathbf{w}_1 \text{ i.i.d (Bernoulli with } p = 1/2)$$

- ▶ State Feedback $\mathbf{u}_0 = \gamma_0(\mathbf{x}_0)$ $\mathbf{u}_1 = \gamma_1(\mathbf{x}_1)$

- ▶ $V_2(x) = x^2$

- ▶ $V_1(x) = \min_u \mathbb{E} [x^2 + V_2(x - u + W_2)]$
 $= \min_u x^2 + (x - u + 1)^2/2 + (x - u)^2/2 = x^2 + 1/4$

- ▶ $V_0(x) = \min_u \mathbb{E} [V_1(x - u + W_1)]$
 $= \min_u 1/4 + (x - u + 1)^2/2 + (x - u)^2/2 = 1/2$

- ▶ Open Loop Controls $\mathbf{u}_0 = u_0$, $\mathbf{u}_1 = u_1$

- ▶ $\min_{u_0, u_1} \mathbb{E} \left[(-u_0 + \mathbf{w}_1)^2 + (-u_0 + \mathbf{w}_1 - u_1 + \mathbf{w}_2)^2 \right]$
 $= \min_{u_0, u_1} 2u_0^2 + 2u_0u_1 + u_1^2 - 3u_0 - 2u_1 + 2 = 3/4$

$$1/2 = \min_{\mathbf{u}_0, \mathbf{u}_1} \mathbb{E} \left[J(\mathbf{u}_0, \mathbf{u}_1) \right] < \min_{u_0, u_1} \mathbb{E} \left[J(u_0, u_1) \right] = 3/4$$

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Start with a Simplified Problem \rightarrow just a Final Cost

Problem (\mathcal{P}_0) starting at position x at initial time $t = 0$:

$$V_0(x) = \min_{\mathbf{x}, \mathbf{u} \in \mathcal{U}} \mathbb{E} [K(\mathbf{x}_T)],$$

s.c. $\mathbf{x}_0 = x$,

$$\mathbf{x}_{t+1} = f_t(\mathbf{x}_t, \mathbf{u}_t, \mathbf{w}_{t+1}), \quad \forall t = 0, \dots, T-1,$$

- ▶ Noises $\mathbf{W} = (\mathbf{w}_t)_{t=1, \dots, T}$ (The demand)
- ▶ Controls $\mathbf{U} = (\mathbf{u}_t)_{t=0, \dots, T-1}$ (Newspaper to order)
- ▶ States $(\mathbf{x}_t)_{t=0, \dots, T-1}$ (Stock of Newspaper)

Markovian Dynamics

The noises and initial state $\mathbf{X}_0, \mathbf{W}_1, \dots, \mathbf{W}_T$ are independent r.v

- ▶ Transition Matrix : uncontrolled case $f : \mathbb{X} \times \mathbb{W} \rightarrow \mathbb{X}$

$$\mathbf{X}_{t+1} = f(\mathbf{X}_t, \mathbf{W}_{t+1}), \quad P(x, y) = \mathbb{P}(f(x, \mathbf{W}_1) = y)$$

- ▶ Transition Matrix : controlled case $f : \mathbb{X} \times \mathbb{U} \times \mathbb{W} \rightarrow \mathbb{X}$ with markovian policy $(\phi_s)_{s \in [0, T-1]}$, $\mathbf{U}_t = \phi_t(\mathbf{X}_t)$

$$\mathbf{X}_{t+1} = f(\mathbf{X}_t, \phi_t(\mathbf{X}_t), \mathbf{W}_{t+1}), \quad P_t(x, y) = \mathbb{P}(f(x, \phi_t(x), \mathbf{W}_1) = y)$$

For $u \in \mathbb{U}$, let $P^u := \mathbb{P}(f(x, u, \mathbf{W}_1) = y)$. For a Markovian policy $(\phi_s)_{s \in [0, T-1]}$, P_t^ϕ is defined by $P_t^\phi(x, y) := P^{\phi_t(x)}(x, y)$

$$\mathbf{X}_{t+1} = f(\mathbf{X}_t, \phi_t(\mathbf{X}_t), \mathbf{W}_{t+1}), \quad P_t^\phi(x, y) = P^{\phi_t(x)}(x, y)$$

A Family of problems

- ▶ Problem (\mathcal{P}_{t_0}) starting with stock x at time t_0 :

$$V_{t_0}(x) = \min_{\mathbf{x}, \phi(\cdot)} V_{t_0}^{\phi}(x)$$

$$V_{t_0}^{\phi}(x) = \mathbb{E} [K(\mathbf{X}_T)],$$

$$\text{s.t } \mathbf{X}_{t_0} = x, \quad \mathbf{X}_{t+1} = f_t(\mathbf{X}_t, \phi_t(\mathbf{X}_t), \mathbf{W}_{t+1})$$

- ▶ Problem (\mathcal{P}'_{t_0}) démarrant en μ à t_0 :

$$\mathcal{V}_{t_0}(\mu) = \min_{\mu, \phi(\cdot)} \mathcal{V}_{t_0}^{\phi}(\mu)$$

$$\mathcal{V}_{t_0}^{\phi}(\mu) = \sum_x \mu_T(x) K(x)$$

$$\text{with } \mu_{t_0} = \mu, \quad \mu_{t+1} = \mu_t P_t^{\phi}$$

- ▶ Dynamic $\mu_{t+1}(y) = \sum_x \mu_t(x) P_{x,y}^{\phi(x)}$

Links between $\mathcal{V}_{t_0}^\phi(\cdot)$ and $V_{t_0}^\phi(\cdot)$

We have that

$$\mathcal{V}_{t_0}^\phi(\mu) = \langle \mu, V_{t_0}^\phi \rangle := \sum_x \mu(x) V_{t_0}^\phi(x), \text{ and } V_{t_0}^\phi(x) = \mathcal{V}_{t_0}^\phi(\delta_x(\cdot))$$

Indeed :

- ▶ Problem (\mathcal{P}_{t_0}) starting with stock x at t_0 :

$$V_{t_0}^\phi(x) = (P_{t_0}^\phi \cdots P_{T-1}^\phi K)(x)$$

$$\text{(Ex. for } t_0 = T - 1 \quad V_{T-1}^\phi(x) = \sum_{y \in \mathbb{X}} P_{T-1}^\phi(x, y) K(y))$$

- ▶ Problem (\mathcal{P}'_{t_0}) starting with μ at t :

$$\mathcal{V}_{t_0}^\phi(\mu) = \mu P_{t_0}^\phi \cdots P_{T-1}^\phi K$$

$$\text{(Ex. for } t_0 = T - 1 \quad \mathcal{V}_{T-1}^\phi(\mu) = \sum_{x, y \in \mathbb{X}} \mu(x) P_{T-1}^\phi(x, y) K(y))$$

The case $T = 1$

$$X_0 = x$$

$$V(x) = \mathbb{E} [K(\mathbf{X}_1)] = \sum_y P_{x,y} K(y)$$

Assume that the law of \mathbf{X}_0 is μ

$$\mathcal{V}(\mu) = \mathbb{E} [K(\mathbf{X}_1)] = \sum_x \mu(x) \sum_y P_{x,y} K(y)$$

We obtain

$$\mathcal{V}(\mu) = \sum_y \mu(x) V(x) = \langle \mu, V \rangle ,$$

and

$$\mathcal{V}(\delta_{x'}) = \sum_x \delta_{x'}(x) \sum_y P_{x,y} K(y) = \sum_y \mu(x) P_{x',y} K(y) = V(x').$$

Recursive Computation of \mathcal{V}_t

We have that :

$$\mathcal{V}_t(\mu) = \min_{\phi_t} \mathcal{V}_{t+1}(\mu P_t^\phi)$$

Proof : The problem (\mathcal{P}'_t) starting with μ at time t :

$$\mathcal{V}_t(\mu) = \min_{\phi(\cdot)} \mu P_t^\phi \cdots P_{T-1}^\phi K$$

- ▶ At time t , P_t^ϕ only depends on ϕ_t .
- ▶ The raw vector $\mu P_{t_0}^\phi$ is non negative (its a probability law)

$$\mathcal{V}_t(\mu) = \min_{\phi_t} \left\langle \mu P_t^\phi, \min_{(\phi_s)_{s>t}} P_{t+1}^\phi \cdots P_{T-1}^\phi K \right\rangle$$

$$\mathcal{V}_t(\mu) = \min_{\phi_t} \left\langle \mu P_t^\phi, \mathcal{V}_{t+1}(\cdot) \right\rangle = \min_{\phi_t} \mathcal{V}_{t+1}(\mu P_t^\phi)$$

Recursive Equation for V_t with $t \in \{0, \dots, T\}$

Bellman Equation :

$$V_t(x) = \min_{u \in \mathcal{U}} \mathbb{E} [V_{t+1}(f_t(x, u, \mathbf{W}_{t+1}))]$$

$$u^\sharp(x) \in \underset{u \in \mathcal{U}}{\text{Argmin}} \mathbb{E} [V_{t+1}(f_t(x, u, \mathbf{W}_{t+1}))]$$

$$V_T(x) = K(x)$$

Proof : We already have for \mathcal{V}_t that

$$\mathcal{V}_t(\mu) = \min_{\phi_t} \mathcal{V}_{t+1}(\mu P_t^\phi)$$

- ▶ $V_t(x) = \mathcal{V}_t(\delta_x(\cdot))$
- ▶ $V_{t+1}(\delta_x(\cdot) P_t^\phi) = \mathbb{E} [V_{t+1}(f_t(x, \phi_t, \mathbf{W}_{t+1}))]$

Recursive Equation for V_t with $t \in \{0, \dots, T\}$

Bellman Equation

$$V_t(x) = \min_{u \in \mathcal{U}} \sum_y P_{x,y}^u V_{t+1}(y)$$

$$V_T(x) = K(x)$$

Optimal control

$$u^\#(x) \in \operatorname{Argmin}_{u \in \mathcal{U}} \sum_y P_{x,y}^u V_{t+1}(y)$$

Proof :

$$\begin{aligned} \sum_y P_{x,y}^u V_{t+1}(y) &= \sum_y \mathbb{P}(f(x, u, \mathbf{w}_{t+1}) = y) V_{t+1}(y) \\ &= \mathbb{E}[V_{t+1}(f_t(x, u, \mathbf{w}_{t+1}))] \end{aligned}$$

Finite Horizon with instantaneous costs

The problem (\mathcal{P}_0) starting with x at initial time $t = 0$:

$$V_0(x) = \min_{\mathbf{x}, \mathbf{u} \in \mathcal{U}} \mathbb{E} \left[\sum_{t=0}^{T-1} L_t(\mathbf{x}_t, \mathbf{u}_t, \mathbf{w}_{t+1}) + K(\mathbf{x}_T) \right],$$

s.t. $\mathbf{x}_0 = x$,

$$\mathbf{x}_{t+1} = f_t(\mathbf{x}_t, \mathbf{u}_t, \mathbf{w}_{t+1}), \quad \forall t = 0, \dots, T-1,$$

Dynamic Programming Equation (Bellman Equation)

$$V_t(x) = \min_{u \in \mathcal{U}} \mathbb{E} [L_t(x, u, \mathbf{w}_{t+1}) + V_{t+1}(f_t(x, u, \mathbf{w}_{t+1}))]$$

$$V_T(x) = K(x)$$

Finite Horizon with instantaneous costs

The cost of problem (\mathcal{P}_0) starting at x at $t = 0$ equals $\tilde{V}(0, x)$

$$\begin{aligned}\tilde{V}_0(z, x) &= \min_{\mathbf{x}, \mathbf{u} \in \mathcal{U}} \mathbb{E} [Z_T + K(\mathbf{X}_T)], \\ \text{s.t. } \mathbf{X}_0 &= x, \mathbf{X}_{t+1} = f_t(\mathbf{X}_t, \mathbf{U}_t, \mathbf{W}_{t+1}), \\ \mathbf{Z}_0 &= z, \mathbf{Z}_{t+1} = \mathbf{Z}_t + L_t(\mathbf{X}_t, \mathbf{U}_t, \mathbf{W}_{t+1}),\end{aligned}$$

- ▶ (\mathbf{Z}, \mathbf{X}) is Markovian for feedback controls $\phi_t(z, x)$.
- ▶ The Bellman Equation is $\tilde{V}_T(z, x) = z + K(x)$ and

$$\tilde{V}_t(z, x) = \min_{u \in \mathcal{U}} \mathbb{E} [\tilde{V}_{t+1}(z + L_t(x, u, \mathbf{W}_{t+1}), f_t(x, u, \mathbf{W}_{t+1}))]$$

Finite Horizon with instantaneous costs

- ▶ We recursively show that $\tilde{V}_t(z, x) = z + V_t(x)$
- ▶ true for $t = T$ since $\tilde{V}_T(z, x) = z + K(x)$
- ▶ at time t

$$\tilde{V}_t(z, x) = \min_{u \in \mathbb{U}} \mathbb{E} [z + L_t(x, u, \mathbf{W}_{t+1}) + V_{t+1}(f_t(x, u, \mathbf{W}_{t+1}))]$$

$$\tilde{V}_t(z, x) = z + \underbrace{\min_{u \in \mathbb{U}} \mathbb{E} [L_t(x, u, \mathbf{W}_{t+1}) + V_{t+1}(f_t(x, u, \mathbf{W}_{t+1}))]}_{V_t(x)}$$

- ▶ the minimization for $u \in \mathbb{U}$ only depends on x . Thus, the optimal control is a feedback on the state x .
- ▶ we note that $\tilde{V}_t(0, x) = V_t(x)$, giving the Bellman Equation for the problem with instantaneous cost.

Final Result : The Bellman Equation

$$V_0(x) = \min_{\mathbf{x}, \mathbf{u} \in \mathcal{U}} \mathbb{E} \left[\sum_{t=0}^{T-1} L_t(\mathbf{x}_t, \mathbf{u}_t, \mathbf{w}_{t+1}) + K(\mathbf{x}_T) \right],$$

s.t. $\mathbf{x}_0 = x$,

$$\mathbf{x}_{t+1} = f_t(\mathbf{x}_t, \mathbf{u}_t, \mathbf{w}_{t+1}), \quad \forall t = 0, \dots, T-1,$$

Dynamic Programming Equation (Bellman Equation)

$$V_t(x) = \min_{u \in \mathcal{U}} \mathbb{E} \left[L_t(x, u, \mathbf{w}_{t+1}) + V_{t+1}(f_t(x, u, \mathbf{w}_{t+1})) \right]$$

$$V_T(x) = K(x)$$

$$u^*(x) \in \operatorname{Argmin}_{u \in \mathcal{U}} \mathbb{E} \left[L_t(x, u, \mathbf{w}_{t+1}) + V_{t+1}(f_t(x, u, \mathbf{w}_{t+1})) \right]$$

Computing value functions

1. Time loop backward to compute V_t for all t
2. State loop for all x
3. Find the optimal control (control loop, L.P., Q.P)
4. Loop on random values to compute the expected cost

State Feedback Versus Open-Loop Feedback (II)

$$\min_{\mathbf{u}_1, \mathbf{u}_2} \mathbb{E} \left[\mathbf{x}_1^2 + \mathbf{x}_2^2 \right]$$

$$\text{s.t.} \quad \mathbf{x}_2 = \mathbf{x}_1 - \mathbf{u}_1 + \mathbf{w}_2, \quad \mathbf{x}_1 = \mathbf{x}_0 - \mathbf{u}_0 + \mathbf{w}_1, \quad \mathbf{x}_0 = 0$$
$$\mathbf{w}_0, \mathbf{w}_1 \text{ i.i.d (Bernoulli with } p = 1/2)$$

- ▶ State Feedback $\mathbf{u}_0 = \gamma_0(\mathbf{x}_0)$ $\mathbf{u}_1 = \gamma_1(\mathbf{x}_1)$
 - ▶ $V_2(x) = x^2$
 - ▶ $V_1(x) = \min_u \mathbb{E} [x^2 + V_2(x - u + W_2)]$
 $= \min_u x^2 + (x - u + 1)^2/2 + (x - u)^2/2 = x^2 + 1/4$
 - ▶ $V_0(x) = \min_u \mathbb{E} [V_1(x - u + W_1)]$
 $= \min_u 1/4 + (x - u + 1)^2/2 + (x - u)^2/2 = 1/2$
 - ▶ Open Loop Controls $\mathbf{u}_0 = u_0$, $\mathbf{u}_1 = u_1$
 - ▶ $\min_{u_0, u_1} \mathbb{E} \left[(-u_0 + \mathbf{w}_1)^2 + (-u_0 + \mathbf{w}_1 - u_1 + \mathbf{w}_2)^2 \right]$
 $= \min_{u_0, u_1} 2u_0^2 + 2u_0u_1 + u_1^2 - 3u_0 - 2u_1 + 2 = 3/4$
- $$1/2 = \min_{\mathbf{u}_0, \mathbf{u}_1} \mathbb{E} \left[J(\mathbf{u}_0, \mathbf{u}_1) \right] < \min_{u_0, u_1} \mathbb{E} \left[J(u_0, u_1) \right] = 3/4$$

A Farmer problem

- ▶ when annual production is x units of a certain crop
 - ▶ he stores $(1 - u)x$ units,
 - ▶ he uses the remaining ux units for next year production, where $u \in (0, 1)$
- ▶ then, the level of next year production will be $\mathbf{W}ux$, where \mathbf{W} is a positive random variable not depending on x or u with known expectation $\mathbb{E}[\mathbf{W}] = \bar{w}$.
- ▶ Optimization problem : find the optimal investment policy that maximizes the total expected product stored over N years

$$\mathbb{E} \left[\sum_{k=0}^{N-1} (1 - \mathbf{U}_k) \mathbf{X}_k + \mathbf{X}_N \right],$$

assuming that $\mathbf{X}_{k+1} = \mathbf{W}_{k+1} \mathbf{U}_k \mathbf{X}_k$.

Bellman Equation

$$V_n(x) = \max_{\mathbf{u}_n, \mathbf{u}_{n+1}, \dots, \mathbf{u}_{N-1}} \mathbb{E} \left[\sum_{k=n}^{N-1} (1 - \mathbf{u}_k) \mathbf{x}_k + \mathbf{x}_N \right]$$
$$\mathbf{x}_{k+1} = \mathbf{w}_{k+1} \mathbf{u}_k \mathbf{x}_k \text{ and } \mathbf{x}_n = x.$$

We obtain that $V_N(x) = x$ and

$$V_n(x) = \max_{u \in [0,1]} (1 - u)x + \mathbb{E} [V_{n+1}(\mathbf{w}ux)]$$

Assume that $V_{n+1}(x) = a_{n+1}x$ then we have that

$$V_n(x) = \begin{cases} x & \text{when } a_{n+1} \bar{w} \leq 1 \\ a_{n+1} \bar{w} x & \text{when } a_{n+1} \bar{w} \geq 1 \end{cases}$$

That is $V_n(x) = a_n x$ with $a_n = \max(1, a_{n+1} \bar{w})$ (with $a_N = 1$).

Increasing the State Space

How to solve

$$V_0(x) = \min_{\mathbf{x}, \mathbf{u} \in \mathcal{U}} \mathbb{E} \left[K \left(\max_{s \in \{0, \dots, T\}} \mathbf{X}_s \right) \right],$$

s.t. $\mathbf{X}_0 = x$,

$$\mathbf{X}_{t+1} = f_t(\mathbf{X}_t, \mathbf{U}_t, \mathbf{W}_{t+1}), \quad \forall t = 0, \dots, T-1,$$

The noises and initial state $\mathbf{X}_0, \mathbf{W}_1, \dots, \mathbf{W}_T$ are independent r.v

- ▶ $\mathbf{Y}_t = \max_{s \in \{0, \dots, t\}} \mathbf{X}_s$, is not a Markov chain.
- ▶ $(\mathbf{X}_t, \mathbf{Y}_t)$ is a Markov chain. $(\mathbf{X}_0, \mathbf{Y}_0) = (x, x)$ and

$$\begin{aligned} \mathbf{X}_{t+1} &= f_t(\mathbf{X}_t, \mathbf{U}_t, \mathbf{W}_{t+1}) \\ \mathbf{Y}_{t+1} &= \max_{s \in \{0, \dots, t+1\}} \mathbf{X}_s = \max(\mathbf{Y}_t, \mathbf{X}_{t+1}) \\ &= \max(\mathbf{Y}_t, f_t(\mathbf{X}_t, \mathbf{U}_t, \mathbf{W}_{t+1})) \end{aligned}$$

Increasing the State Space

How to solve

$$V_0(x) = \min_{\mathbf{x}, \mathbf{u} \in \mathcal{U}} \mathbb{E} \left[K(\mathbf{X}_T) \right],$$

s.t. $\mathbf{X}_0 = x$ fixed,

$$\mathbf{X}_{t+1} = f_t(\mathbf{X}_t, \mathbf{U}_t, \mathbf{W}_{t+1}), \quad \forall t = 0, \dots, T-1,$$

The noises and initial state $\mathbf{X}_0, \mathbf{W}_1, \dots, \mathbf{W}_T$ are **not** independent

- ▶ $\mathbf{W}_{t+1} = g_t(\mathbf{W}_t, \overline{\mathbf{W}}_{t+1})$
- ▶ $\mathbf{X}_0, \overline{\mathbf{W}}_1, \dots, \overline{\mathbf{W}}_T$ are independent
- ▶ $(\mathbf{X}_t, \mathbf{W}_t)$ is a Markov chain

$$\mathbf{X}_{t+1} = f_t(\mathbf{X}_t, \mathbf{U}_t, \mathbf{W}_{t+1}) = f_t(\mathbf{X}_t, \mathbf{U}_t, g_t(\mathbf{W}_t, \overline{\mathbf{W}}_{t+1}))$$
$$\mathbf{W}_{t+1} = g_t(\mathbf{W}_t, \overline{\mathbf{W}}_{t+1})$$