

# COURS RTE 2022

## OPTIMISATION STOCHASTIQUE



### Méthodes de décomposition spatiale I

Mélange des techniques de décomposition  
et de programmation dynamique

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# Ultimate goal of the lecture

How to obtain “good” **strategies** for a **large scale** stochastic optimal control problem, for example a problem corresponding to the optimal management over a given time horizon of a system involving a large amount of dynamical production units.

- In order to obtain **decision strategies** (closed-loop controls), we have to use **Dynamic Programming** or related methods.
  - **Assumption**: Markovian case,
  - **Difficulty**: **curse of dimensionality**.
- In order to take into account the size of the system, we have to use **decomposition/coordination** techniques.
  - **Assumption**: convexity,
  - **Difficulty**: **information pattern** of the problem.

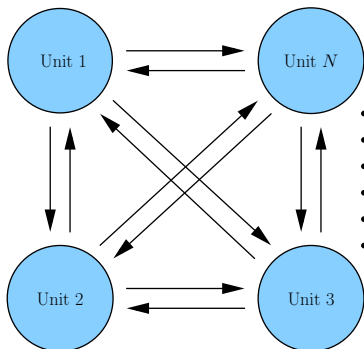
**Mixture of spatial and temporal decompositions**

# Lecture outline

- 1 Examples and background
  - Examples of interconnected systems
  - Convex optimization background
- 2 Decomposition in the deterministic case
  - Additive model: 3 decomposition methods
  - General model: Auxiliary Problem Principle
- 3 About decomposition in the stochastic case
  - Dynamic Programming and decomposition
  - Couplings in stochastic optimization

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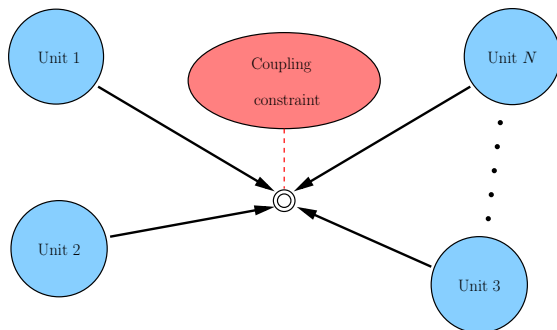
# Decomposition and coordination



Interconnected units

- The (**large**) system to be optimized consists of **interconnected** subsystems: we want to use this structure in order to formulate optimization **subproblems** of **reasonable** complexity.
- But the presence of **interactions** requires a level of **coordination**.
- Coordination must provide a **local model** of the interactions to each subproblem: it is an **iterative** process.
- The ultimate goal is to obtain the solution of the **overall problem** by concatenation of the solutions of the **subproblems**.

# Example: the “flower model”

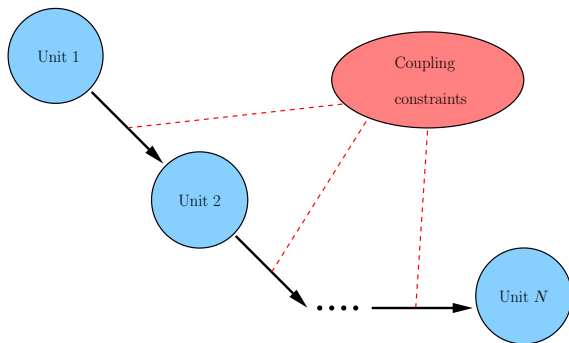


$$\min_u \sum_{i=1}^N J_i(u_i),$$

$$\text{s.t.} \quad \sum_{i=1}^N \Theta_i(u_i) = \theta.$$

Unit Commitment Problem

# Example: the “cascade model”



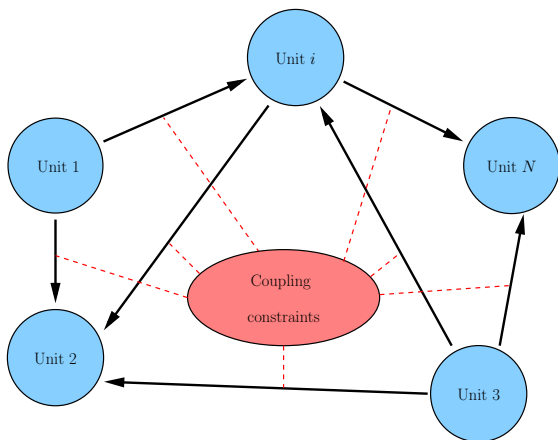
$$\min_{u,v} \sum_{i=1}^N J_i(u_i, v_i),$$

$$\text{s.t. } H_i(u_i, v_i) = v_{i+1} \quad \forall i.$$

Dams Management Problem

Link with the flower model:  $\Theta_i \rightsquigarrow (0, \dots, -v_i, H_i(u_i, v_i), \dots, 0)^\top$ .

# A general model



$$\min_{u,v} \sum_{i=1}^N J_i \left( u_i, \sum_{j \neq i} v_{j,i} \right),$$

$$\text{s.t. } H_i \left( u_i, \sum_{j \neq i} v_{j,i} \right) = v_i.$$

Microgrid Management Problem



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# Optimization without explicit constraint

$$\min_{u \in \mathcal{U}^{\text{ad}}} J(u) . \quad (\mathcal{P}_S)$$

- $\mathcal{U}$ : Hilbert space with scalar product  $\langle \cdot, \cdot \rangle$ .  
Examples:  $\mathcal{U} = \mathbb{R}^n$  (vectors) or  $\mathcal{U} = L^2(\Omega, \mathcal{A}, \mathbb{P}; \mathbb{R}^n)$  (random variables).
- $\mathcal{U}^{\text{ad}}$ : closed convex subset of  $\mathcal{U}$ .
- $J : \mathcal{U} \rightarrow \mathbb{R}$ : function satisfying some properties (convexity, continuity, differentiability, coercivity).

**Characterization** of a solution  $u^\sharp$  (**optimality conditions**):

$$\langle \nabla J(u^\sharp), u - u^\sharp \rangle \geq 0 \quad \forall u \in \mathcal{U}^{\text{ad}} .$$

**Computation** of the solution  $u^\sharp$  (**projected gradient algorithm**):

$$u^{(k+1)} = \text{proj}_{\mathcal{U}^{\text{ad}}} \left( u^{(k)} - \rho \nabla J(u^{(k)}) \right) .$$

# Optimization with explicit constraints

$$\min_{u \in \mathcal{U}^{\text{ad}}} J(u) \quad \text{subject to} \quad \Theta(u) \in -C. \quad (\mathcal{P}_C)$$

- $\mathcal{U}$ : Hilbert space.
- $\mathcal{U}^{\text{ad}}$ : closed convex subset of  $\mathcal{U}$ .
- $\mathcal{V}$ : another Hilbert space.
- $C$ : **cone** of  $\mathcal{V}$  (examples:  $C = \{0\}$ ,  $C = \{v \geq 0\}$ ).
- $J : \mathcal{U} \rightarrow \mathbb{R}$ : cost function.
- $\Theta : \mathcal{U} \rightarrow \mathcal{V}$ : constraint function satisfying some properties (convexity w.r.t.  $C$ , continuity, differentiability).

**Constraint Qualification Condition**, e.g.  $0 \in \text{int}(\Theta(\mathcal{U}^{\text{ad}}) + C)$ .

The **dual cone** of  $C$  is defined by:  $C^* = \{\lambda \in \mathcal{V}, \langle \lambda, v \rangle \geq 0 \ \forall v \in C\}$ .

# Optimization with explicit constraints

II

## Karush-Kuhn-Tucker Conditions

In addition to standard conditions on  $J$  and  $\Theta$ , we assume that the constraints are **qualified**.

Then a **necessary and sufficient** condition for  $u^\# \in \mathcal{U}^{\text{ad}}$  to be a solution of Problem  $(\mathcal{P}_C)$  is that there exists  $\lambda^\# \in \mathcal{V}$  such that:

- ①  $\langle \nabla J(u^\#) + [\Theta'(u^\#)]^* \lambda^\#, u - u^\# \rangle \geq 0 \quad \forall u \in \mathcal{U}^{\text{ad}},$
- ②  $\Theta(u^\#) \in -C,$
- ③  $\lambda^\# \in C^*,$
- ④  $\langle \lambda^\#, \Theta(u^\#) \rangle = 0$  (*Complementary Slackness*).

# Optimization with explicit constraints



Let  $L : \mathcal{U}^{\text{ad}} \times C^* \rightarrow \mathbb{R}$  be the **Lagrangian** associated to  $(\mathcal{P}_C)$ :

$$L(u, \lambda) = J(u) + \langle \lambda, \Theta(u) \rangle .$$

A point  $(u^\#, \lambda^\#) \in \mathcal{U}^{\text{ad}} \times C^*$  is a **saddle point** of  $L$  if

$$L(u^\#, \lambda) \leq L(u^\#, \lambda^\#) \leq L(u, \lambda^\#), \quad \forall (u, \lambda) \in \mathcal{U}^{\text{ad}} \times C^* .$$

- If  $(u^\#, \lambda^\#)$  is a **saddle point** of  $L$ , then  $u^\#$  is a **solution** of  $(\mathcal{P}_C)$ .
- If  $u^\#$  is a **solution** of  $(\mathcal{P}_C)$  and if the **KKT** conditions are met for some  $\lambda^\#$ , then  $(u^\#, \lambda^\#)$  is a **saddle point** of  $L$ .

Moreover we have that

$$J(u^\#) = \min_{u \in \mathcal{U}^{\text{ad}}} \max_{\lambda \in C^*} L(u, \lambda) = \max_{\lambda \in C^*} \min_{u \in \mathcal{U}^{\text{ad}}} L(u, \lambda) = L(u^\#, \lambda^\#) .$$

# Optimization with explicit constraints

IV

Define the **dual function** associated to the Lagrangian  $L$  as

$$\Phi(\lambda) = \min_{u \in \mathcal{U}^{\text{ad}}} L(u, \lambda),$$

and assume that  $\arg \min L(\cdot, \lambda) = \{\hat{u}_\lambda\}$ , so that  $\nabla \Phi(\lambda) = \Theta(\hat{u}_\lambda)$ .

To compute the solution  $u^\sharp$ , use a **gradient algorithm** for Problem:

$$\max_{\lambda \in C^*} \Phi(\lambda) \quad \left( \Leftrightarrow \max_{\lambda \in C^*} \min_{u \in \mathcal{U}^{\text{ad}}} L(u, \lambda) \right).$$

## Uzawa's Algorithm

Choose  $\lambda^{(0)} \in C^*$ . At each itération  $k$ ,

- obtain the solution  $u^{(k+1)} = \arg \min_{u \in \mathcal{U}^{\text{ad}}} J(u) + \langle \lambda^{(k)}, \Theta(u) \rangle$ ,
- update the multiplier  $\lambda^{(k+1)} = \text{proj}_{C^*} (\lambda^{(k)} + \rho \Theta(u^{(k+1)}))$ .

# Optimization with explicit constraints

## Uzawa's algorithm convergence theorem

- H1**  $\mathcal{U}^{\text{ad}}$  is a closed convex subset of the Hilbert space  $\mathcal{U}$ ,  
 $\mathcal{C}$  is a closed convex cone of the Hilbert space  $\mathcal{V}$ .
- H2**  $J$  is a proper l.s.c. **strongly convex** function with modulus  $a$ ,  
 Gâteaux différentiable.
- H3**  $\Theta$  is a  $\mathcal{C}$ -convex, Lipschitz with constant  $\tau$ .
- H4**  $L$  admits a saddle point  $(u^\#, \lambda^\#) \in \mathcal{U}^{\text{ad}} \times \mathcal{C}^*$ .
- H5**  $\rho$  is such that  $0 < \rho < 2a/\tau^2$ .
- R1** The sequence  $\{u^{(k)}\}_{k \in \mathbb{N}}$  converges toward  $u^\#$ .
- R2** The sequence  $\{\lambda^{(k)}\}_{k \in \mathbb{N}}$  is bounded, and any of its cluster  
 points  $\bar{\lambda}$  is such that  $(u^\#, \bar{\lambda})$  is a saddle point of  $L$ .

# Uzawa's geometric interpretation

For the sake of simplicity, we consider here **equality** constraints:

$$u^{(k+1)} \in \arg \min_{u \in \mathcal{U}^{\text{ad}}} J(u) + \langle \lambda^{(k)}, \Theta(u) \rangle ,$$

$$\lambda^{(k+1)} = \lambda^{(k)} + \rho \Theta(u^{(k+1)}) .$$

The minimization step is equivalent to:

$$\min_{v \in \mathcal{V}} \min_{u \in \mathcal{U}^{\text{ad}}} J(u) + \langle \lambda^{(k)}, v \rangle \quad \text{s.t.} \quad \Theta(u) - v = 0 .$$

Introducing the **perturbation function**  $G$ :

$$G(v) = \min_{u \in \mathcal{U}^{\text{ad}}} J(u) \quad \text{s.t.} \quad \Theta(u) - v = 0 ,$$

this minimization step also writes:

$$\min_{v \in \mathcal{V}} G(v) + \langle \lambda^{(k)}, v \rangle .$$



# Uzawa's geometric interpretation



With the help of  $G$ , Uzawa's algorithm writes:

$$v^{(k+1)} \in \arg \min_{v \in \mathcal{V}} G(v) + \langle \lambda^{(k)}, v \rangle,$$

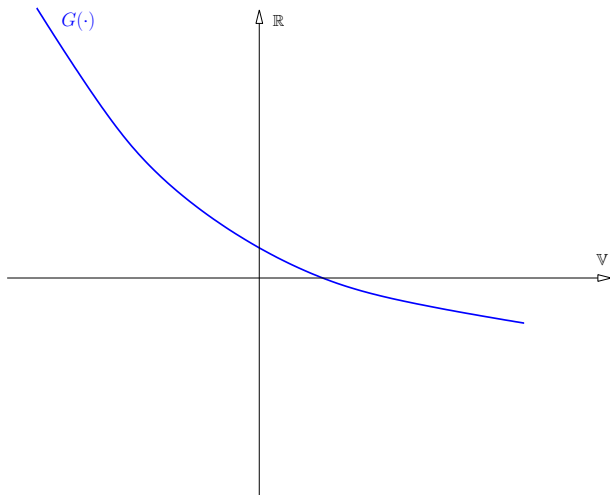
$$\lambda^{(k+1)} = \lambda^{(k)} + \rho v^{(k+1)}.$$

From a (conceptual) geometric point of view, it amounts to:

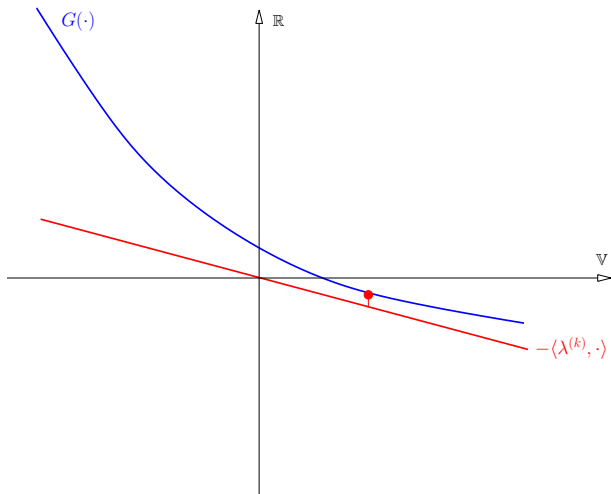
- **Step 1:** **minimize** the gap between  $G(\cdot)$  et  $\langle -\lambda^{(k)}, \cdot \rangle$ .
- **Step 2:** **adjust** the slope  $-\lambda^{(k)}$  if  $v^{(k+1)} \neq 0$ .

*Recall that the initial problem consists in obtaining  $G(0)$ ...*

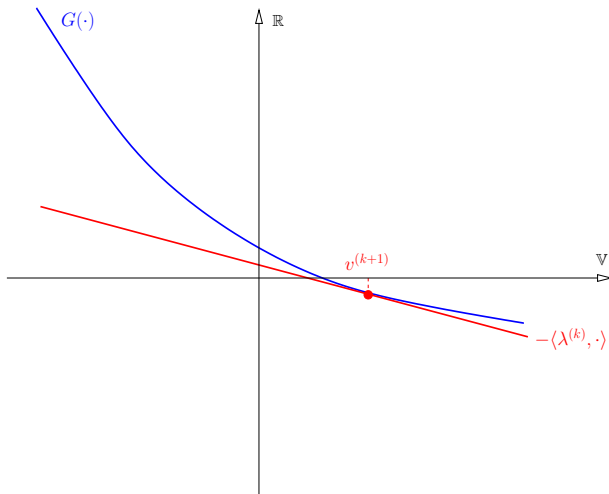
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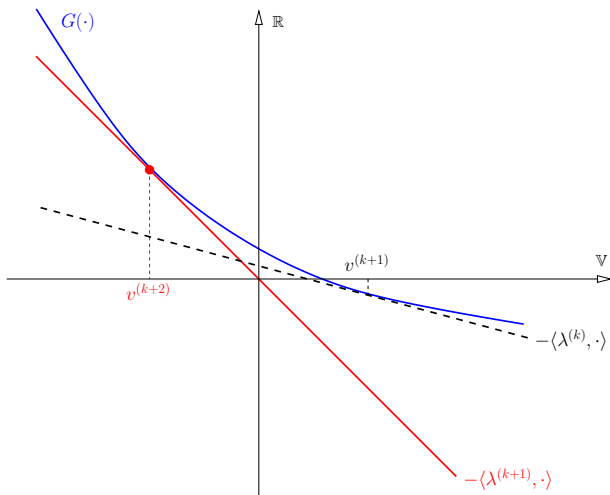
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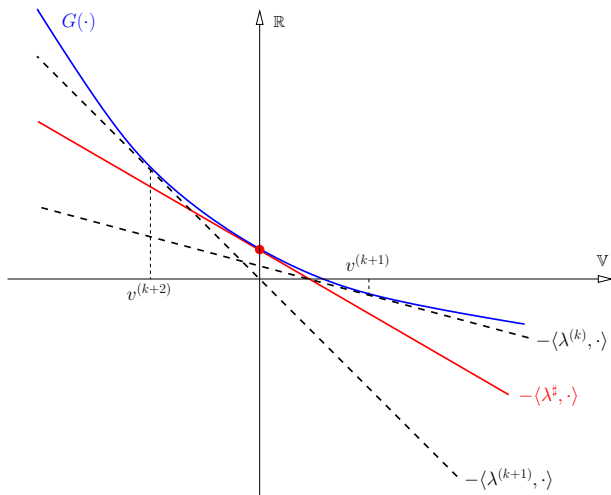
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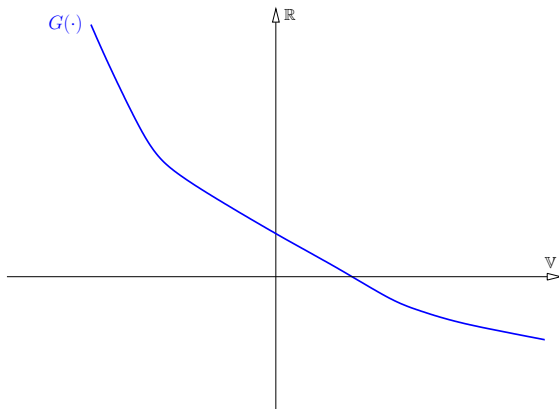


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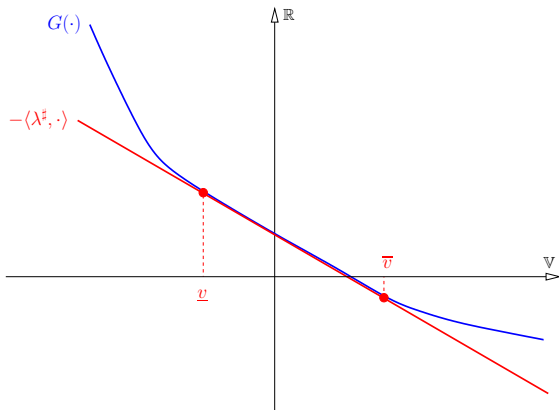
# Uzawa's geometric interpretation

IV



## Uzawa's geometric interpretation

## IV

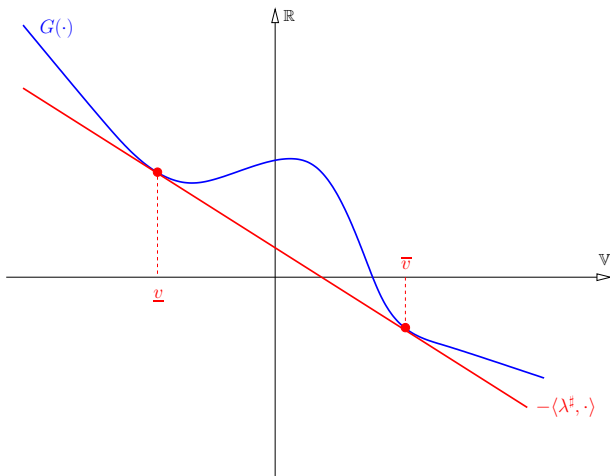


Even if  $\{\lambda^{(k)}\}_{k \in \mathbb{N}}$  converges towards  $\lambda^\sharp$ , the constraint level  $v^{(k)}$  oscillates between  $\underline{v}$  and  $\bar{v}$ , but the value  $v^\sharp = 0$  is **never reached**.



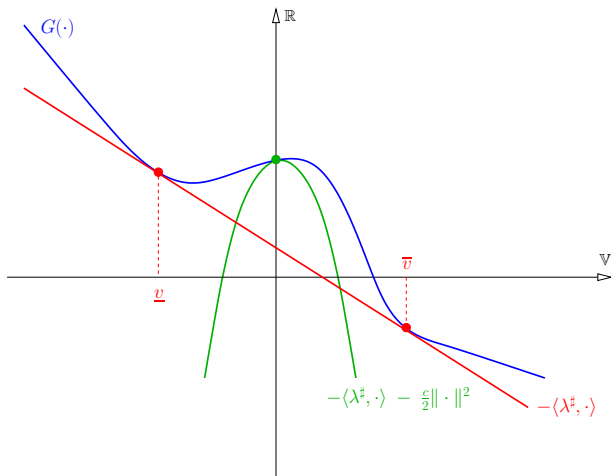
## Uzawa's geometric interpretation

V



## Uzawa's geometric interpretation

V



In the non convex case, use an **augmented Lagrangian**...

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# Additive model

Consider the following problem:

$$\min_{u \in \mathcal{U}^{\text{ad}} \subset \mathcal{U}} J(u) \quad \text{subject to} \quad \Theta(u) - \theta = 0 \in \mathcal{V},$$

and consider a **decomposition** of the space  $\mathcal{U} = \mathcal{U}_1 \times \dots \times \mathcal{U}_N$ , so that  $u \in \mathcal{U}$  writes  $u = (u_1, \dots, u_N)$  with  $u_i \in \mathcal{U}_i$ . Assume that

- $\mathcal{U}^{\text{ad}} = \mathcal{U}_1^{\text{ad}} \times \dots \times \mathcal{U}_N^{\text{ad}} \quad \mathcal{U}_i^{\text{ad}} \subset \mathcal{U}_i,$
- $J(u) = J_1(u_1) + \dots + J_N(u_N) \quad u_i \in \mathcal{U}_i,$
- $\Theta(u) = \Theta_1(u_1) + \dots + \Theta_N(u_N) \quad u_i \in \mathcal{U}_i.$

Then the problem displays the following **additive structure**:

$$\min_{\substack{u_1 \in \mathcal{U}_1^{\text{ad}} \\ \vdots \\ u_N \in \mathcal{U}_N^{\text{ad}}}} \sum_{i=1}^N J_i(u_i) \quad \text{subject to} \quad \sum_{i=1}^N \Theta_i(u_i) - \theta = 0.$$

Note that the **coupling** between the  $i$ 's **only** arises from the constraint  $\Theta$ .

# Additive model — Price decomposition

$$\min_{u \in \mathcal{U}^{\text{ad}}} \sum_{i=1}^N J_i(u_i) \quad \text{subject to} \quad \sum_{i=1}^N \Theta_i(u_i) - \theta = 0 .$$

- Form the **Lagrangian** of the problem. We assume that a saddle point exists, so that solving the initial problem is equivalent to:

$$\max_{\lambda \in \mathcal{V}} \min_{u \in \mathcal{U}^{\text{ad}}} \sum_{i=1}^N \left( J_i(u_i) + \langle \lambda, \Theta_i(u_i) \rangle \right) - \langle \lambda, \theta \rangle ,$$

- Solve this problem by the Uzawa algorithm:

$$u_i^{(k+1)} \in \arg \min_{u_i \in \mathcal{U}_i^{\text{ad}}} J_i(u_i) + \langle \lambda^{(k)}, \Theta_i(u_i) \rangle, \quad i = 1, \dots, N,$$

$$\lambda^{(k+1)} = \lambda^{(k)} + \rho \left( \sum_{i=1}^N \Theta_i(u_i^{(k+1)}) - \theta \right).$$

## Additive model — Price decomposition

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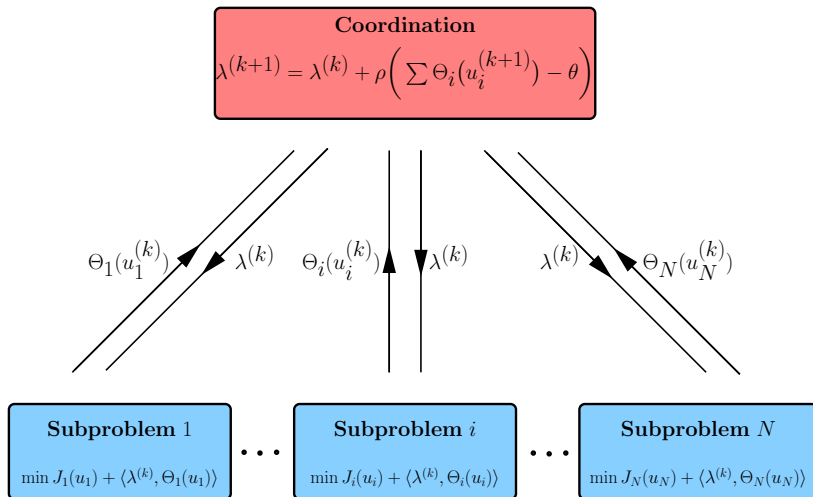
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$$\lambda^{(k+1)} = \lambda^{(k)} + \rho \left( \sum_{i=1}^N \Theta_i(u_i^{(k+1)}) - \theta \right).$$

## Additive model — Price decomposition

II





## Additive model — Resource allocation

$$\min_{u \in \mathcal{U}^{\text{ad}}} \sum_{i=1}^N J_i(u_i) \quad \text{subject to} \quad \sum_{i=1}^N \Theta_i(u_i) - \theta = 0 .$$

- 1 Write the constraint in an equivalent manner by introducing **new variables**  $v = (v_1, \dots, v_N)$  (the so-called “**allocation**”):

$$\sum_{i=1}^N \Theta_i(u_i) - \theta = 0 \quad \Leftrightarrow \quad \Theta_i(u_i) - v_i = 0 \quad \text{and} \quad \sum_{i=1}^N v_i = \theta ,$$

and minimize the criterion w.r.t.  $u$  and  $v$ :

$$\min_{v \in \mathcal{V}^N} \sum_{i=1}^N \left( \min_{u_i \in \mathcal{U}_i^{\text{ad}}} J_i(u_i) \text{ s.t. } \Theta_i(u_i) - v_i = 0 \right) \text{ s.t. } \sum_{i=1}^N v_i = \theta ,$$

## Additive model — Resource allocation

II

$$\min_{v \in \mathcal{V}^N} \sum_{i=1}^N \underbrace{\left( \min_{u_i \in \mathcal{U}_i^{\text{ad}}} J_i(u_i) \text{ s.t. } \Theta_i(u_i) - v_i = 0 \right)}_{G_i(v_i)} \text{ s.t. } \sum_{i=1}^N v_i = \theta ,$$

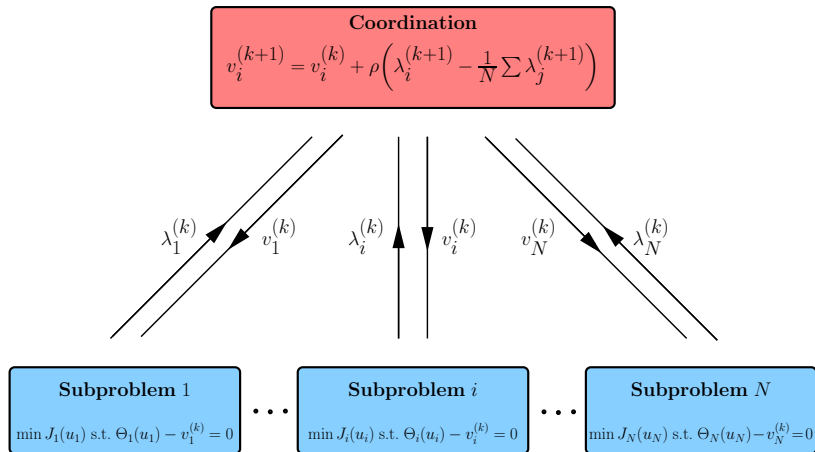
$$\iff \min_{v \in \mathcal{V}^N} \sum_{i=1}^N G_i(v_i) \text{ s.t. } \sum_{i=1}^N v_i = \theta .$$

- 2 Solve the last problem using a **projected gradient method**:

$$G_i(v_i^{(k)}) = \min_{u_i \in \mathcal{U}_i^{\text{ad}}} J_i(u_i) \text{ s.t. } \Theta_i(u_i) - v_i^{(k)} = 0 \rightsquigarrow \lambda_i^{(k+1)} ,$$

$$v_i^{(k+1)} = v_i^{(k)} + \rho \left( \lambda_i^{(k+1)} - \frac{1}{N} \sum_{j=1}^N \lambda_j^{(k+1)} \right) .$$

## Additive model — Resource allocation



## Additive model — Prediction

$$\min_{u \in \mathcal{U}^{\text{ad}}} \sum_{i=1}^N J_i(u_i) \quad \text{subject to} \quad \sum_{i=1}^N \Theta_i(u_i) - \theta = 0 .$$

We assume for the moment that the constraint is *scalar*...

- 1 Choose the unit that will drive the constraint (e.g. unit 1) and **split the constraint** according to that choice:

$$\Theta_1(u_1) - v = 0 \quad , \quad \sum_{i \neq 1} \Theta_i(u_i) - \theta + v = 0 .$$

- 2 Formulate the problem obtained by dualizing **only** the second part of the constraint:

$$\max_{\lambda \in \mathbb{R}} \min_{v \in \mathcal{V}} \left( \min_{u \in \mathcal{U}^{\text{ad}}} \sum_{i=1}^N J_i(u_i) + \left\langle \lambda, \sum_{i \neq 1} \Theta_i(u_i) - \theta + v \right\rangle \right)$$

subject to  $\Theta_1(u_1) - v = 0 .$

## Additive model — Prediction



- ③ With  $v = v^{(k)}$  and  $\lambda = \lambda^{(k)}$  fixed, the problem decomposes:

$$\min_{u_1 \in \mathcal{U}_1^{\text{ad}}} J_1(u_1) \text{ s.t. } \Theta_1(u_1) - v^{(k)} = 0 \rightsquigarrow \lambda_1^{(k+1)},$$

$$\min_{u_i \in \mathcal{U}_i^{\text{ad}}} J_i(u_i) + \langle \lambda^{(k)}, \Theta_i(u_i) \rangle \quad \forall i \neq 1 \rightsquigarrow \Theta_i(u_i^{(k+1)}).$$

- ④ Update  $v$  and  $\lambda$  by solving the **optimality conditions** in  $\lambda$  and  $v$  of the global problem:

$$v^{(k+1)} = \theta - \sum_{i \neq 1} \Theta_i(u_i^{(k+1)}),$$

$$\lambda^{(k+1)} = \lambda_1^{(k+1)}.$$

*In case of multiple constraints, incorporate them one by one.  
A choice has to be done for each constraint. The constraints  
are thus distributed among the units.*

## Additive model — Prediction



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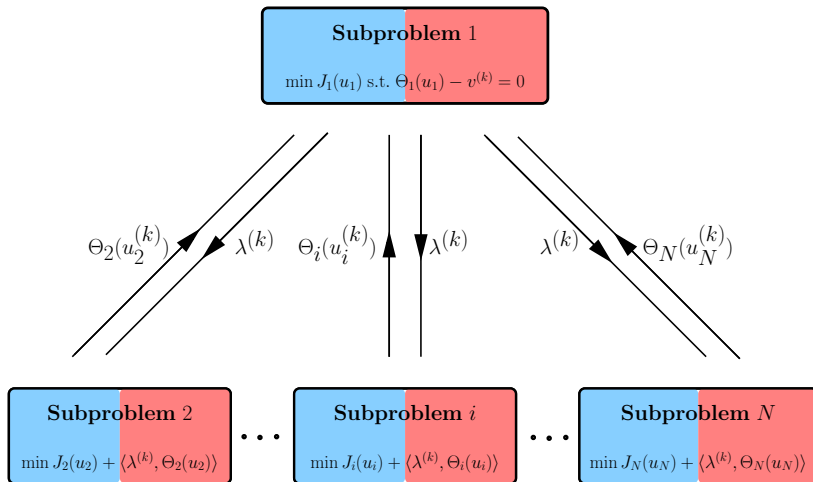
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*In case of multiple constraints, incorporate them one by one. A choice has to be done for each constraint. The constraints are thus **distributed among the units**.*

## Additive model — Prediction



# Additive model: conclusions

## 1 Price decomposition

- **Pros:** “non-destructive” method.
- **Cons:** admissible solution once convergence achieved.

Gives at each iteration a **lower bound** of the optimal cost.

## 2 Resource allocation

- **Pros:** admissible solution at each iteration.
- **Cons:** potential existence of unfeasible subproblems.

Gives at each iteration an **upper bound** of the optimal cost.

## 3 Prediction

- **Pros and Cons:** depending on the constraints distribution. . .

*Straightforward extension to inequality constraints. . .*



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- 3 About decomposition in the stochastic case
  - Dynamic Programming and decomposition
  - Couplings in stochastic optimization

## General model — Auxiliary Problem Principle

The 3 decomposition schemes we have presented seem to depend crucially on the **additive structure** of the underlying problems. . .  
In fact they can be **extended** to general problems:

$$\min_{u \in \mathcal{U}^{\text{ad}}} J(u_1, \dots, u_N) \quad \text{s.t.} \quad \Theta(u_1, \dots, u_N) - \theta = 0 .$$

This generalization is achieved by the **Auxiliary Problem Principle (APP)**, whose aim is to recover **additivity** by replacing the two functions  $J$  and  $\Theta$  by their **first-order approximation** around the current point  $u^{(k)}$ :

$$J(u) \rightsquigarrow \sum_{i=1}^N \langle \nabla_{u_i} J(u^{(k)}), u_i \rangle, \quad \Theta(u) \rightsquigarrow \sum_{i=1}^N \Theta'_{u_i}(u^{(k)}) \cdot u_i .$$

The solution  $u^{(k+1)}$  of the **auxiliary problem** built around  $u^{(k)}$  is used to formulate the next auxiliary problem (**iterative process**).

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The solution  $u^{(k+1)}$  of the **auxiliary problem** built around  $u^{(k)}$  is used to formulate the next auxiliary problem (**iterative process**).

## APP without explicit constraint

$$\min_{u \in \mathcal{U}^{\text{ad}}} J(u).$$

- ① Replace  $J(u)$  by its **first order approximation** around  $u^{(k)}$ :

$$J(u^{(k)}) + \langle \nabla J(u^{(k)}), u - u^{(k)} \rangle.$$

- ② **Choose** a strongly convex function  $K$ , some  $\epsilon > 0$  and form:

$$\frac{1}{\epsilon} \left( K(u) - K(u^{(k)}) - \langle \nabla K(u^{(k)}), u - u^{(k)} \rangle \right).$$

Add these two terms to obtain the **auxiliary problem** at iteration  $k$ :

$$\min_{u \in \mathcal{U}^{\text{ad}}} K(u) + \langle \epsilon \nabla J(u^{(k)}) - \nabla K(u^{(k)}), u \rangle,$$

whose unique solution is denoted by  $u^{(k+1)}$ .

## APP without explicit constraint



## Convergence theorem

- H1**  $\mathcal{U}^{\text{ad}}$  is a closed convex subset of the Hilbert space  $\mathcal{U}$ .
- H2**  $J$  is a proper l.s.c. convex function, coercive over  $\mathcal{U}^{\text{ad}}$ , and its derivative  $J'$  is Lipschitz with constant  $A$ .
- H3**  $K$  is a proper l.s.c. strongly convex function with modulus  $b$ , and its derivative  $K'$  is Lipschitz with constant  $B$ .
- H4**  $\epsilon$  is a coefficient such that  $0 < \epsilon < 2b/A$ .
- R1**  $\{J(u^{(k)})\}_{k \in \mathbb{N}}$  is a strictly decreasing real sequence which converges towards  $J(u^\#)$ .
- R2**  $\{u^{(k)}\}_{k \in \mathbb{N}}$  is a bounded sequence, and each of its cluster points is a solution of the initial problem.

Moreover, if  $J$  is strongly convex, then  $\{u^{(k)}\}_{k \in \mathbb{N}}$  converges to  $u^\#$ .

## APP without explicit constraint

Consider the auxiliary problem obtained at iteration  $k$ :

$$\min_{u \in \mathcal{U}^{\text{ad}}} K(u) + \langle \epsilon \nabla J(u^{(k)}) - \nabla K(u^{(k)}), u \rangle .$$

Assume that there exists a **decomposition**  $\mathcal{U}_1 \times \dots \times \mathcal{U}_N$  of  $\mathcal{U}$ , that is,  $u \in \mathcal{U}$  writes  $u = (u_1, \dots, u_N)$  with  $u_i \in \mathcal{U}_i$ , such that:

$$\mathcal{U}^{\text{ad}} = \mathcal{U}_1^{\text{ad}} \times \dots \times \mathcal{U}_N^{\text{ad}} \quad \text{with} \quad \mathcal{U}_i^{\text{ad}} \subset \mathcal{U}_i .$$

A **additive** choice of  $K$  leads to decomposition. Indeed, using

$$K(u) = \sum_{i=1}^N K_i(u_i) ,$$

the  $k$ -th auxiliary problem can be decomposed in  $N$  subproblems:

$$\min_{u_i \in \mathcal{U}_i^{\text{ad}}} K_i(u_i) + \langle \epsilon \nabla_{u_i} J(u^{(k)}) - \nabla_{u_i} K(u^{(k)}), u_i \rangle , \quad i = 1, \dots, N .$$

## APP without explicit constraint

## IV

## Variants of the algorithm

- Take into account an additional cost function  $J^\Sigma$ :

$$\min_{u \in \mathcal{U}^{\text{ad}}} K(u) + \langle \epsilon \nabla J(u^{(k)}) - \nabla K(u^{(k)}), u \rangle + \epsilon J^\Sigma(u).$$

- $K$  and  $\epsilon$  may depend on the iteration index  $k$ :

$$\min_{u \in \mathcal{U}^{\text{ad}}} K^{(k)}(u) + \langle \epsilon^{(k)} \nabla J(u^{(k)}) - \nabla K^{(k)}(u^{(k)}), u \rangle.$$

- Use  $\epsilon \equiv 1$  by adding an under-relaxation step in the algorithm:

$$u^{(k+\frac{1}{2})} = \arg \min_{u \in \mathcal{U}^{\text{ad}}} K(u) + \langle \nabla J(u^{(k)}) - \nabla K(u^{(k)}), u \rangle,$$

$$u^{(k+1)} = \rho u^{(k+\frac{1}{2})} + (1 - \rho) u^{(k)}, \quad 0 < \rho < 1.$$

## APP with explicit constraints

$$\min_{u \in \mathcal{U}^{\text{ad}}} J(u) \quad \text{s.t.} \quad \Theta(u) \in -C,$$

Denote by  $L(u, \lambda) = J(u) + \langle \lambda, \Theta(u) \rangle$  the associated Lagrangian.

- ① Replace  $L$  by its **first order approximation** around  $(u^{(k)}, \lambda^{(k)})$ :

$$L(u^{(k)}, \lambda^{(k)}) + \langle \nabla_u L(u^{(k)}, \lambda^{(k)}), u - u^{(k)} \rangle + \langle \nabla_\lambda L(u^{(k)}, \lambda^{(k)}), \lambda - \lambda^{(k)} \rangle.$$

- ② **Choose** a convex-concave operator  $M(u, \lambda)$  and some  $\epsilon > 0$ .

Use these elements to form the **auxiliary Lagrangian** at iteration  $k$ :

$$M(u, \lambda) + \langle (\epsilon \nabla_u L - \nabla_u M)(u^{(k)}, \lambda^{(k)}), u \rangle + \langle (\epsilon \nabla_\lambda L - \nabla_\lambda M)(u^{(k)}, \lambda^{(k)}), \lambda \rangle,$$

and obtain a point  $(u^{(k+1)}, \lambda^{(k+1)})$  satisfying optimality conditions.



## APP with explicit constraints



We denote by  $\mathcal{L}^{(k)}$  the **auxiliary Lagrangian** at iteration  $k$ :

$$\mathcal{L}^{(k)}(u, \lambda) = M(u, \lambda) + \langle \epsilon \nabla_u L(u^{(k)}, \lambda^{(k)}) - \nabla_u M(u^{(k)}, \lambda^{(k)}), u \rangle + \langle \epsilon \nabla_\lambda L(u^{(k)}, \lambda^{(k)}) - \nabla_\lambda M(u^{(k)}, \lambda^{(k)}), \lambda \rangle .$$

We have two possible algorithms to solve the auxiliary problem.

- ① SIM: solve **simultaneously** the optimality conditions:

$$u^{(k+1)} = \arg \min_{u \in \mathcal{U}^{\text{ad}}} \mathcal{L}^{(k)}(u, \lambda^{(k+1)}) ,$$

$$\lambda^{(k+1)} = \arg \max_{\lambda \in C^*} \mathcal{L}^{(k)}(u^{(k+1)}, \lambda) .$$

- ② SEQ: solve **sequentially** the optimality conditions:

$$u^{(k+1)} = \arg \min_{u \in \mathcal{U}^{\text{ad}}} \mathcal{L}^{(k)}(u, \lambda^{(k)}) ,$$

$$\lambda^{(k+1)} = \arg \max_{\lambda \in C^*} \mathcal{L}^{(k)}(u^{(k+1)}, \lambda) .$$

## APP with explicit constraints

We denote by  $\mathfrak{L}^{(k)}$  the **auxiliary Lagrangian** at iteration  $k$ :

$$\mathfrak{L}^{(k)}(u, \lambda) = M(u, \lambda) + \langle \epsilon \nabla_u L(u^{(k)}, \lambda^{(k)}) - \nabla_u M(u^{(k)}, \lambda^{(k)}), u \rangle + \langle \epsilon \nabla_\lambda L(u^{(k)}, \lambda^{(k)}) - \nabla_\lambda M(u^{(k)}, \lambda^{(k)}), \lambda \rangle .$$

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## APP with explicit constraints: one-level algorithm

I

Possible **choice**:  $M(u, \lambda) = K(u) + \langle \lambda, \Omega(u) \rangle$  and Algorithm SIM.

The expression of the auxiliary Lagrangian is as follows:

$$\begin{aligned} \mathcal{L}^{(k)}(u, \lambda) &= M(u, \lambda) + \langle \epsilon \nabla_u L(u^{(k)}, \lambda^{(k)}) - \nabla_u M(u^{(k)}, \lambda^{(k)}), u \rangle \\ &\quad + \langle \epsilon \nabla_\lambda L(u^{(k)}, \lambda^{(k)}) - \nabla_\lambda M(u^{(k)}, \lambda^{(k)}), \lambda \rangle \\ &= K(u) + \langle \epsilon \nabla J(u^{(k)}) - \nabla K(u^{(k)}), u \rangle \\ &\quad + \langle \lambda^{(k)}, (\epsilon \Theta'(u^{(k)}) - \Omega'(u^{(k)})), u \rangle \\ &\quad + \langle \lambda, \Omega(u) + \epsilon \Theta(u^{(k)}) - \Omega(u^{(k)}) \rangle. \end{aligned}$$

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## APP with explicit constraints: one-level algorithm

II

The **saddle point**  $(u^{(k+1)}, \lambda^{(k+1)})$  of  $\mathcal{L}^{(k)}$  is obtained by solving the associated constrained optimization problem:

$$\min_{u \in \mathcal{U}^{\text{ad}}} K(u) + \langle \epsilon \nabla J(u^{(k)}) - \nabla K(u^{(k)}), u \rangle + \langle \lambda^{(k)}, (\epsilon \Theta'(u^{(k)}) - \Omega'(u^{(k)})) \cdot u \rangle,$$

$$\text{subject to } \Omega(u) - \Omega(u^{(k)}) + \epsilon \Theta(u^{(k)}) \in -C.$$

The convergence proof of this algorithm is available for problems involving a **quadratic** cost function and **linear equality** constraints. Moreover, a geometric condition, namely  $\Theta J^{-1} \Omega^* + \Omega J^{-1} \Theta^* > 0$  (**weak coupling** through the constraints) has to be met.

## APP with explicit constraints: one-level algorithm

## III

With regard to decomposition, consider the following choices:

$$K(u) = \sum_{i=1}^N K_i(u_i) \quad , \quad \Omega(u) = \begin{pmatrix} \Omega_1(u_1) & \cdots & 0 \\ \vdots & \ddots & \vdots \\ 0 & \cdots & \Omega_N(u_N) \end{pmatrix} ,$$

that is,

- an **additive** auxiliary cost function  $K$ ,
- a **block diagonal** auxiliary constraint  $\Omega$ ,

and assume that  $U^{\text{ad}} = U_1^{\text{ad}} \times \dots \times U_N^{\text{ad}}$ .

Then the auxiliary problem can be decomposed in  $N$  subproblems.

This algorithm is in fact a generalization of the **decomposition by prediction** that has been studied for additive models. The choice of  $\Omega$  as a block-diagonal operator corresponds to the **distribution of the constraints** among the units.

## APP with explicit constraints: two-level algorithm

I

Alternative **choice**:  $M(u, \lambda) = K(u) - \frac{\|\lambda\|^2}{2\alpha}$  and Algorithm SEQ.

The expression of the auxiliary Lagrangian is as follows:

$$\begin{aligned} \mathcal{L}^{(k)}(u, \lambda) = & M(u, \lambda) + \langle \epsilon \nabla_u L(u^{(k)}, \lambda^{(k)}) - \nabla_u M(u^{(k)}, \lambda^{(k)}), u \rangle \\ & + \langle \epsilon \nabla_\lambda L(u^{(k)}, \lambda^{(k)}) - \nabla_\lambda M(u^{(k)}, \lambda^{(k)}), \lambda \rangle, \end{aligned}$$

so that

$$\begin{aligned} \mathcal{L}^{(k)}(u, \lambda^{(k)}) \leftrightarrow & K(u) + \langle \epsilon \nabla J(u^{(k)}) - \nabla K(u^{(k)}), u \rangle \\ & + \epsilon \langle \lambda^{(k)}, \Theta'(u^{(k)}), u \rangle, \end{aligned}$$

$$\mathcal{L}^{(k)}(u^{(k+1)}, \lambda) \leftrightarrow -\frac{1}{2} \|\lambda\|^2 + \langle \alpha \epsilon \Theta(u^{(k+1)}) + \lambda^{(k)}, \lambda \rangle.$$

## APP with explicit constraints: two-level algorithm

Alternative **choice**:  $M(u, \lambda) = K(u) - \frac{\|\lambda\|^2}{2\alpha}$  and Algorithm SEQ.

The expression of the auxiliary Lagrangian is as follows:

$$\begin{aligned} \mathcal{L}^{(k)}(u, \lambda) = & M(u, \lambda) + \langle \epsilon \nabla_u L(u^{(k)}, \lambda^{(k)}) - \nabla_u M(u^{(k)}, \lambda^{(k)}), u \rangle \\ & + \langle \epsilon \nabla_\lambda L(u^{(k)}, \lambda^{(k)}) - \nabla_\lambda M(u^{(k)}, \lambda^{(k)}), \lambda \rangle, \end{aligned}$$

so that

$$\begin{aligned} \mathcal{L}^{(k)}(u, \lambda^{(k)}) \quad \leftrightarrow \quad & K(u) + \langle \epsilon \nabla J(u^{(k)}) - \nabla K(u^{(k)}), u \rangle \\ & + \epsilon \langle \lambda^{(k)}, \Theta'(u^{(k)}) \cdot u \rangle, \end{aligned}$$

$$\mathcal{L}^{(k)}(u^{(k+1)}, \lambda) \leftrightarrow -\frac{1}{2} \|\lambda\|^2 + \langle \alpha \epsilon \Theta(u^{(k+1)}) + \lambda^{(k)}, \lambda \rangle.$$



## APP with explicit constraints: two-level algorithm

II

The optimization problems are solved **sequentially**. Solving the first problem  $\min_{u \in \mathcal{U}^{\text{ad}}} \mathfrak{L}^{(k)}(u, \lambda^{(k)})$  leads to

$$\min_{u \in \mathcal{U}^{\text{ad}}} K(u) + \langle \epsilon \nabla J(u^{(k)}) - \nabla K(u^{(k)}), u \rangle + \epsilon \langle \lambda^{(k)}, \Theta'(u^{(k)}) \cdot u \rangle,$$

whose solution is denoted  $u^{(k+1)}$ , and solving the second problem  $\max_{\lambda \in C^*} \mathfrak{L}^{(k)}(u^{(k+1)}, \lambda)$  is equivalent to

$$\lambda^{(k+1)} = \text{proj}_{C^*}(\lambda^{(k)} + \underbrace{\alpha \epsilon}_{\rho} \Theta(u^{(k+1)})),$$

that is, an **update** of the multiplier  $\lambda$ .

The convergence proof of this algorithm can be established under standard assumptions in the **convex** (sub)differentiable framework.

## APP with explicit constraints: two-level algorithm

III

## Convergence theorem

- H1**  $\mathcal{U}^{\text{ad}}$  is a closed convex subset of the Hilbert space  $\mathcal{U}$ , and  $\mathcal{C}$  is a closed convex cone of the Hilbert space  $\mathcal{V}$ .
- H2**  $J$  is a proper l.s.c. **strongly convex** function with modulus  $a$ , and its derivative  $J'$  is Lipschitz with constant  $A$ .
- H3**  $\Theta$  is a  $\mathcal{C}$ -convex, Lipschitz with constant  $\tau$ , differentiable.
- H4** A saddle point  $(u^\#, \lambda^\#)$  of  $L$  exists.
- H5**  $K$  is a proper l.s.c. strongly convex function with modulus  $b$ , and its derivative  $K'$  is Lipschitz with constant  $B$ .
- H6**  $\epsilon$  and  $\rho$  are such that  $0 < \epsilon < b/A$  and  $0 < \rho < a/\tau^2$ .
- R1** The sequence  $\{u^{(k)}\}_{k \in \mathbb{N}}$  converges toward  $u^\#$ .
- R2** The sequence  $\{\lambda^{(k)}\}_{k \in \mathbb{N}}$  is bounded, and any of its cluster points  $\bar{\lambda}$  is such that  $(u^\#, \bar{\lambda})$  is a saddle point of  $L$ .

## APP with explicit constraints: two-level algorithm

## IV

This algorithm corresponds to a **generalization** of both Uzawa and Arrow-Hurwicz algorithms. Roughly speaking,

- $K(u) = J(u)$  and  $\epsilon = 1 \rightsquigarrow$  **Uzawa**.
- $K(u) = \frac{1}{2} \|u\|^2 \rightsquigarrow$  **Arrow-Hurwicz**.

Choosing an **additive** auxiliary function  $K$ :

$$K(u) = \sum_{i=1}^N K_i(u_i),$$

and assuming that  $\mathcal{U}^{\text{ad}} = \mathcal{U}_1^{\text{ad}} \times \dots \times \mathcal{U}_N^{\text{ad}}$ , the minimization step in the previous algorithm **splits** into  $N$  independent subproblems:

$$\min_{u_i \in \mathcal{U}_i^{\text{ad}}} K_i(u_i) + \langle \epsilon \nabla_{u_i} J(u^{(k)}) - \nabla_{u_i} K(u^{(k)}), u_i \rangle + \epsilon \langle \lambda^{(k)}, \Theta'_{u_i}(u^{(k)}) \cdot u_i \rangle.$$

## APP with explicit constraints: augmented Lagrangian

For the sake of simplicity, consider an optimization problem under **equality** constraints:

$$\min_{u \in \mathcal{U}^{\text{ad}}} J(u) \quad \text{s.t.} \quad \Theta(u) = 0,$$

The **two-level APP algorithm** writes in the following equivalent form:

$$u^{(k+1)} \in \arg \min_{u \in \mathcal{U}^{\text{ad}}} K(u) + \langle \epsilon \nabla_u L(u^{(k)}, \lambda^{(k)}) - \nabla K(u^{(k)}), u \rangle,$$

$$\lambda^{(k+1)} = \lambda^{(k)} + \rho \nabla_\lambda L(u^{(k+1)}, \lambda^{(k)}),$$

$L$  being the **standard Lagrangian**:  $L(u, \lambda) = J(u) + \langle \lambda, \Theta(u) \rangle$ .

## APP with explicit constraints: augmented Lagrangian II

Introduce now the **augmented Lagrangian**  $L_c$ , whose expression in the case of equality constraints is given by

$$L_c(u, \lambda) = L(u, \lambda) + \frac{c}{2} \|\Theta(u)\|^2 .$$

Applying the APP methodology to this new Lagrangian leads to the following two-level algorithm:

$$u^{(k+1)} \in \arg \min_{u \in \mathcal{U}^{\text{ad}}} K(u) + \langle \epsilon \nabla_u L_c(u^{(k)}, \lambda^{(k)}) - \nabla K(u^{(k)}), u \rangle ,$$

$$\lambda^{(k+1)} = \lambda^{(k)} + \rho \nabla_\lambda L_c(u^{(k+1)}, \lambda^{(k)}) ,$$

that is, APP allows to decompose **augmented Lagrangians**!

## APP with explicit constraints: augmented Lagrangian III

## Convergence theorem

- H1**  $\mathcal{U}^{\text{ad}}$  is a closed convex subset of the Hilbert space  $\mathcal{U}$ , and  $\mathcal{C}$  is a closed convex cone of the Hilbert space  $\mathcal{V}$ .
- H2**  $J$  is a proper l.s.c **convex** function, and its derivative  $J'$  is Lipschitz with constant  $A$ .
- H3**  $\Theta$  is a  $C$ -convex, Lipschitz with constant  $\tau$ , differentiable.
- H4** A saddle point  $(u^\#, \lambda^\#)$  exists.
- H5**  $K$  is a proper l.s.c strongly convex function with modulus  $b$ , and its derivative  $K'$  is Lipschitz with constant  $B$ .
- H6**  $\epsilon$  and  $\rho$  are such that  $0 < \epsilon < b/(A + c\tau^2)$  and  $0 < \rho < 2c$ .
- R1** The sequence  $\{(u^{(k)}, \lambda^{(k)})\}_{k \in \mathbb{N}}$  is bounded, and any of its cluster points is a saddle point.

## References on decomposition/coordination methods

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G. COHEN & B. MIARA, "Optimization with an Auxiliary Constraint and Decomposition". *SIAM Journal on Control and Optimization*, **28**, 137-157, 1990.

# Final remarks on decomposition methods

The theory is available for general (infinite dimensional) Hilbert spaces, and thus applies in the **stochastic framework**, that is, the case where  $\mathcal{U}$  is a space of **random variables**.

The **minimization algorithm** used for solving the subproblems is not specified in the decomposition process and is left to the user! It is however assumed that the user is able to solve the subproblem, for example in the price decomposition case:

$$\min_{u_i \in \mathcal{U}_i^{\text{ad}}} J_i(u_i) + \langle \lambda^{(k)}, \Theta_i(u_i) \rangle ,$$

and to send the requested information, namely  $\Theta_i(u_i^{(k+1)})$ , to the coordination level.

**Question:** *what methods are suitable in the stochastic case?*



# Final remarks on decomposition methods

II

Whatever the **decomposition/coordination scheme** used (price, allocation, prediction, APP), **new variables** (depending on  $u^{(k)}$  and/or  $\lambda^{(k)}$ ) appear in the subproblems arising at iteration  $k$  of the optimization process.

**Example:** subproblem  $i$  in price decomposition:

$$\min_{u_i \in \mathcal{U}_i^{\text{ad}}} J_i(u_i) + \langle \lambda^{(k)}, \Theta_i(u_i) \rangle .$$

All these new variables are considered as **fixed** when solving the subproblems (they only depend on the iteration index  $k$ ). They are nothing but constants, and therefore do not cause any trouble in the **deterministic** case.

**Question:** *what happens in the stochastic case?*

- 1 Examples and background
  - Examples of interconnected systems
  - Convex optimization background
- 2 Decomposition in the deterministic case
  - Additive model: 3 decomposition methods
  - General model: Auxiliary Problem Principle
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## Reminder of our ultimate goal

How to obtain “good” **strategies** for a **large scale** stochastic optimal control problem, for example a problem corresponding to the optimal management over a given time horizon of a system involving a large amount of dynamical production units.

- In order to obtain **decision strategies** (closed-loop controls), we have to use **Dynamic Programming** or related methods.
  - **Assumption**: Markovian case,
  - **Difficulty**: **curse of dimensionality**.
- In order to take into account the size of the system, we have to use **decomposition/coordination** techniques.
  - **Assumption**: convexity,
  - **Difficulty**: **information pattern** of the problem.

# Stochastic optimal control problems

We consider a **SOC problem** (in the *Decision-Hazard* setting):

$$\min_{\mathbf{U}, \mathbf{X}} \mathbb{E} \left( \sum_{i=1}^N \left( \sum_{t=0}^{T-1} L_t^i(\mathbf{x}_t^i, \mathbf{u}_t^i, \mathbf{w}_{t+1}) + K^i(\mathbf{x}_T^i) \right) \right),$$

subject to the constraints:

$$\mathbf{x}_0^i = f_{-1}^i(\mathbf{w}_0), \quad i = 1 \dots N,$$

$$\mathbf{x}_{t+1}^i = f_t^i(\mathbf{x}_t^i, \mathbf{u}_t^i, \mathbf{w}_{t+1}), \quad t = 0 \dots T-1, \quad i = 1 \dots N,$$

$$\sigma(\mathbf{u}_t^i) \subset \sigma(\mathbf{w}_0, \dots, \mathbf{w}_t), \quad t = 0 \dots T-1, \quad i = 1 \dots N,$$

$$\sum_{i=1}^N \Theta_t^i(\mathbf{x}_t^i, \mathbf{u}_t^i) = 0, \quad t = 0 \dots T-1.$$

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$$\min_{\mathbf{U}, \mathbf{X}} \sum_{i=1}^N \left( \mathbb{E} \left( \sum_{t=0}^{T-1} L_t^i(\mathbf{x}_t^i, \mathbf{u}_t^i, \mathbf{w}_{t+1}) + K^i(\mathbf{x}_T^i) \right) \right),$$

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$$\sum_{i=1}^N \Theta_t^i(\mathbf{x}_t^i, \mathbf{u}_t^i) = 0, \quad t = 0 \dots T-1.$$

# Dynamic Programming yields centralized controls

Remember that we want to solve this SOC problem using **Dynamic Programming** (DP) or related methods (such as SDDP).

The system is made of  $N$  **interconnected** subsystems, and we have denoted the **control** and the **state** of subsystem  $i$  at time  $t$  by  $U_t^i$  and  $X_t^i$ . Recall that the **optimal** control of subsystem  $i$  when using DP is a function of the **whole** system state:

$$U_t^i = \gamma_t^i(X_t^1, \dots, X_t^N),$$

but a **straightforward use** of DP is prohibited for  $N$  large...

*Decomposition would seem to lead to decentralized feedbacks:*

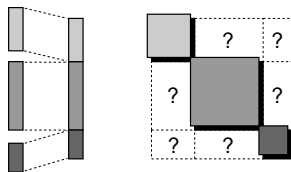
$$U_t^i = \hat{\gamma}_t^i(X_t^i),$$

*which are, in most cases, far from being optimal!*

# Straightforward decomposition of Dynamic Programming?

The crucial point is that the **optimal feedback** of a subsystem a priori depends on the state of all other subsystems, so that using a decomposition scheme by subsystems is far from being obvious. . .

As far as we have to deal with **Dynamic Programming**, the central concern for decomposition/coordination purpose is resumed as:



- how to decompose a feedback  $\gamma_t$  w.r.t. its **domain**  $\mathbb{X}_t$  rather than its **range**  $\mathbb{U}_t$ ?

**And the answer is:**

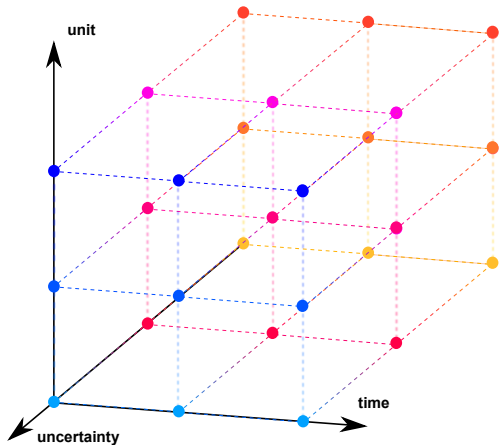
- **impossible** in the general case!

- 1 Examples and background
  - Examples of interconnected systems
  - Convex optimization background
- 2 Decomposition in the deterministic case
  - Additive model: 3 decomposition methods
  - General model: Auxiliary Problem Principle
- 3 About decomposition in the stochastic case
  - Dynamic Programming and decomposition
  - Couplings in stochastic optimization



## Couplings and decompositions for SOC problems

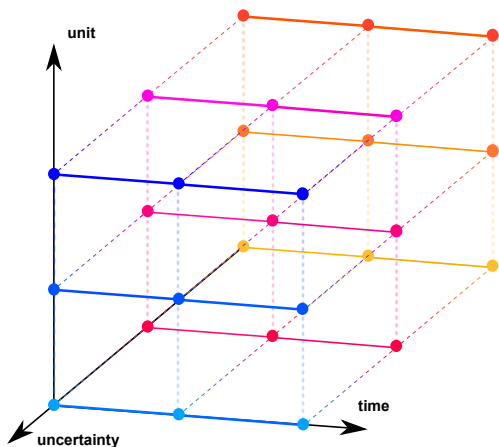
I



$$\min \sum_{\omega} \sum_i \sum_t \pi_{\omega} L_t^i(x_t^i, u_t^i, w_{t+1})$$

## Couplings and decompositions for SOC problems

II

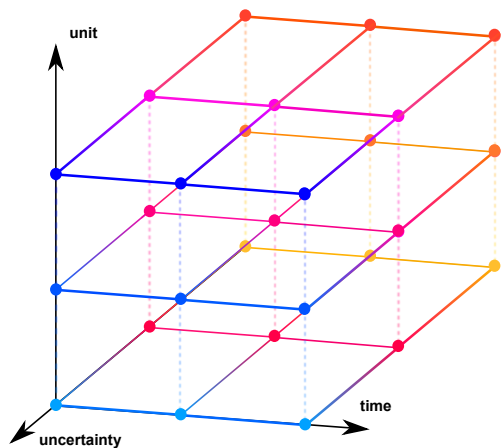


$$\min \sum_{\omega} \sum_i \sum_t \pi_{\omega} L_t^i(\mathbf{x}_t^i, \mathbf{u}_t^i, \mathbf{w}_{t+1})$$

$$\text{s.t. } \mathbf{x}_{t+1}^i = f_t^i(\mathbf{x}_t^i, \mathbf{u}_t^i, \mathbf{w}_{t+1})$$

## Couplings and decompositions for SOC problems

III



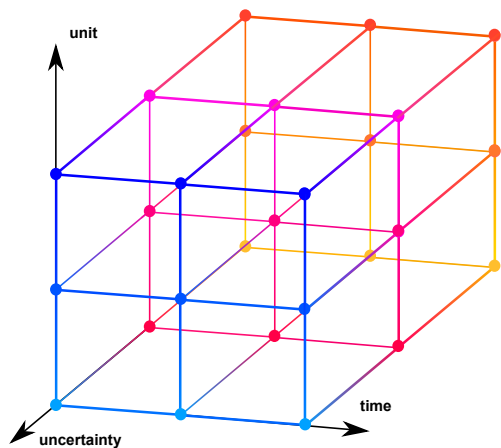
$$\min \sum_{\omega} \sum_i \sum_t \pi_{\omega} L_t^i(\mathbf{x}_t^i, \mathbf{u}_t^i, \mathbf{w}_{t+1})$$

$$\text{s.t. } \mathbf{x}_{t+1}^i = f_t^i(\mathbf{x}_t^i, \mathbf{u}_t^i, \mathbf{w}_{t+1})$$

$$\sigma(\mathbf{u}_t^i) \subset \sigma(\mathbf{w}_0, \dots, \mathbf{w}_t)$$

## Couplings and decompositions for SOC problems

IV



$$\min \sum_{\omega} \sum_i \sum_t \pi_{\omega} L_t^i(\mathbf{x}_t^i, \mathbf{u}_t^i, \mathbf{w}_{t+1})$$

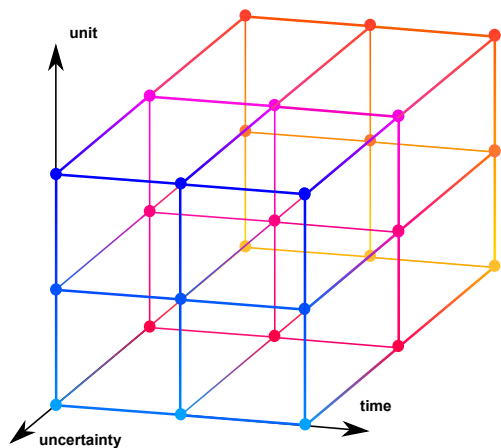
$$\text{s.t. } \mathbf{x}_{t+1}^i = f_t^i(\mathbf{x}_t^i, \mathbf{u}_t^i, \mathbf{w}_{t+1})$$

$$\sigma(\mathbf{u}_t^i) \subset \sigma(\mathbf{w}_0, \dots, \mathbf{w}_t)$$

$$\sum_i \Theta_t^i(\mathbf{x}_t^i, \mathbf{u}_t^i) = 0$$

## Couplings and decompositions for SOC problems

V



$$\min \sum_{\omega} \sum_i \sum_t \pi_{\omega} L_t^i(\mathbf{x}_t^i, \mathbf{u}_t^i, \mathbf{w}_{t+1})$$

$$\text{s.t. } \mathbf{x}_{t+1}^i = f_t^i(\mathbf{x}_t^i, \mathbf{u}_t^i, \mathbf{w}_{t+1})$$

$$\sigma(\mathbf{u}_t^i) \subset \sigma(\mathbf{w}_0, \dots, \mathbf{w}_t)$$

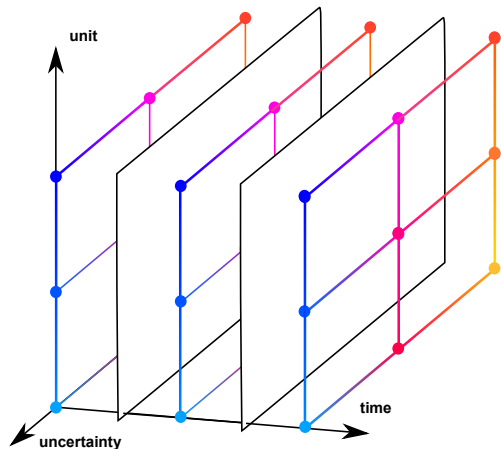
$$\sum_i \Theta_t^i(\mathbf{x}_t^i, \mathbf{u}_t^i) = 0$$

3 additive structures!

Multiple decompositions...

## Couplings and decompositions for SOC problems

VI



$$\min \sum_{\omega} \sum_i \sum_t \pi_{\omega} L_t^i(\mathbf{x}_t^i, \mathbf{u}_t^i, \mathbf{w}_{t+1})$$

$$\text{s.t. } \mathbf{x}_{t+1}^i = f_t^i(\mathbf{x}_t^i, \mathbf{u}_t^i, \mathbf{w}_{t+1})$$

$$\sigma(\mathbf{u}_t^i) \subset \sigma(\mathbf{w}_0, \dots, \mathbf{w}_t)$$

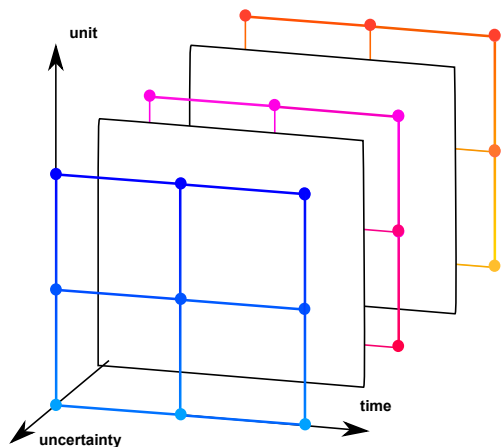
$$\sum_i \Theta_t^i(\mathbf{x}_t^i, \mathbf{u}_t^i) = 0$$

Time decomposition

Dynamic Programming

## Couplings and decompositions for SOC problems

## VII



$$\min \sum_{\omega} \sum_i \sum_t \pi_{\omega} L_t^i(\mathbf{x}_t^i, \mathbf{u}_t^i, \mathbf{w}_{t+1})$$

$$\text{s.t. } \mathbf{x}_{t+1}^i = f_t^i(\mathbf{x}_t^i, \mathbf{u}_t^i, \mathbf{w}_{t+1})$$

$$\sigma(\mathbf{u}_t^i) \subset \sigma(\mathbf{w}_0, \dots, \mathbf{w}_t)$$

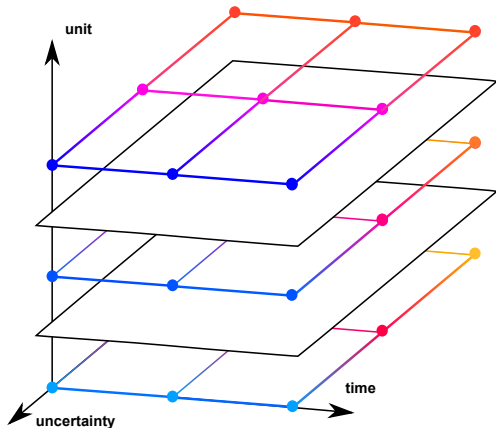
$$\sum_i \Theta_t^i(\mathbf{x}_t^i, \mathbf{u}_t^i) = 0$$

Scenario decomposition

Progressive Hedging

## Couplings and decompositions for SOC problems

## VIII



$$\min \sum_{\omega} \sum_i \sum_t \pi_{\omega} L_t^i(\mathbf{x}_t^i, \mathbf{u}_t^i, \mathbf{w}_{t+1})$$

$$\text{s.t. } \mathbf{x}_{t+1}^i = f_t^i(\mathbf{x}_t^i, \mathbf{u}_t^i, \mathbf{w}_{t+1})$$

$$\sigma(\mathbf{u}_t^i) \subset \sigma(\mathbf{w}_0, \dots, \mathbf{w}_t)$$

$$\sum_i \Theta_t^i(\mathbf{x}_t^i, \mathbf{u}_t^i) = 0$$

Spatial decomposition

Purpose of the lecture



# Price decomposition in the stochastic case

Dualize the **spatial coupling constraints** in the SOC problem:

$$\min_{\mathbf{U}, \mathbf{X}} \sum_{i=1}^N \left( \mathbb{E} \left( \sum_{t=0}^{T-1} L_t^i(\mathbf{x}_t^i, \mathbf{u}_t^i, \mathbf{w}_{t+1}) + K^i(\mathbf{x}_T^i) \right) \right),$$

subject to the constraints:

$$\mathbf{x}_0^i = f_{-1}^i(\mathbf{w}_0), \quad i = 1 \dots N,$$

$$\mathbf{x}_{t+1}^i = f_t^i(\mathbf{x}_t^i, \mathbf{u}_t^i, \mathbf{w}_{t+1}), \quad t = 0 \dots T-1, \quad i = 1 \dots N,$$

$$\sigma(\mathbf{u}_t^i) \subset \sigma(\mathbf{w}_0, \dots, \mathbf{w}_t), \quad t = 0 \dots T-1, \quad i = 1 \dots N,$$

$$\sum_{i=1}^N \Theta_t^i(\mathbf{x}_t^i, \mathbf{u}_t^i) = 0, \quad t = 0 \dots T-1 \quad \rightsquigarrow \quad \boldsymbol{\Lambda}_t.$$

# Price decomposition in the stochastic case

II

Applying **price decomposition** to the previous **SOC** problem leads to a collection of **local stochastic optimal control subproblems** indexed by  $i \in \llbracket 1, N \rrbracket$ :

$$\min_{\mathbf{u}^i, \mathbf{x}^i} \mathbb{E} \left( \sum_{t=0}^{T-1} \left( L_t^i(\mathbf{x}_t^i, \mathbf{u}_t^i, \mathbf{w}_{t+1}) + \boldsymbol{\lambda}_t^{(k)} \cdot \Theta_t^i(\mathbf{x}_t^i, \mathbf{u}_t^i) \right) + K^i(\mathbf{x}_T^i) \right),$$

subject to the constraints:

$$\mathbf{x}_0^i = f_{-1}^i(\mathbf{w}_0),$$

$$\mathbf{x}_{t+1}^i = f_t^i(\mathbf{x}_t^i, \mathbf{u}_t^i, \mathbf{w}_{t+1}), \quad t = 0 \dots T-1,$$

$$\sigma(\mathbf{u}_t^i) \subset \sigma(\mathbf{w}_0, \dots, \mathbf{w}_t), \quad t = 0 \dots T-1.$$

## Price decomposition in the stochastic case



- As pointed out in the deterministic case, **new variables**, that is, dual multipliers  $\Lambda_t^{(k)}$ , appear in the subproblems arising at iteration  $k$ : these variables, **fixed** at this stage of calculation, corresponds to **random variables**.

$$\min_{u^i, x^i} \mathbb{E} \left( \sum_t L_t^i(x_t^i, u_t^i, w_{t+1}) + \Lambda_t^{(k)} \cdot \Theta_t^i(x_t^i, u_t^i) \right).$$

- The process  $\Lambda^{(k)}$  acts as an **additional** input (data) in the subproblems, but the structure of this process is a priori unknown: it may be **correlated** in time, so that the **white noise** assumption, crucial for the **optimality** of Dynamic Programming, has no reason to be fulfilled in that context!

# Summary

- On the one hand, it seems that Dynamic Programming **cannot** be decomposed in a straightforward manner.
- On the other hand, applying a decomposition scheme to a SOC problem introduces **coordination instruments** in the subproblems, e.g. the multipliers  $\Lambda_t^{(k)}$  in the case of price decomposition, which correspond to additional fixed random variables whose **time structure** is **unknown**.

**Question:** how to **handle** the coordination instruments (random variables  $\Lambda_t^{(k)}$  in the case of price decomposition) in order to obtain **an approximation of the overall optimum** of the SOC problem?

BREAK