Two Players Game Theory with Information: Introducing the Witsenhausen Intrinsic Model

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## Outline of the presentation

Introduction to games in intrinsic form

Witsenhausen intrinsic model and game theory with information

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Nash equilibrium with information

Witsenhausen intrinsic model and principal-agent models

Games solvable by dynamic programming

Open questions

# Outline of the presentation

#### Introduction to games in intrinsic form

Witsenhausen intrinsic model and game theory with information

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Open questions

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Information plays a crucial role in competition

- Information who knows what and when plays a crucial role in competitive contexts
- Concealing, dissimulation, cheating, lying, deception are effective strategies

Our goals are to

- $1_{\cdot}$  introduce the notion of game in intrinsic form
- 2. contribute to the analysis of decentralized, non-cooperative decision settings
- 3. provide a (very) general mathematical language for mechanism design

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We will distinguish an individual from an agent

- An individual who makes a first, followed by a second decision, is represented by two agents (two decision makers)
- An individual who makes a sequence of decisions — one for each period t = 0, 1, 2, ..., T − 1 is represented by T agents, labelled t = 0, 1, 2, ..., T − 1
- N individuals each *i* of whom makes a sequence of decisions, one for each period t = 0, 1, 2, ..., T<sub>i</sub> − 1 is represented by ∏<sup>N</sup><sub>i=1</sub> T<sub>i</sub> agents, labelled by

$$(i,t)\in igcup_{j=1}^{N}\{j\} imes\{0,1,2,\ldots,T_{j}-1\}$$

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## What is a game in intrinsic form?

Nature, the source of all randomness, or states of Nature

- Agents, who
  - hold information
  - make *decisions*, by means of *admissible strategies*, those fueled by information
- Players, who
  - hold beliefs about states of Nature
  - hold a subset of agents under their exclusive control (executives)

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hold objectives, that they achieve by selecting proper admissible strategies for the agents under their control

### What is a game in intrinsic form?

- Nature, the source of all randomness a set Ω equipped with a σ-field F
- Agents, who hold information and make decisions a set A
  - ▶ for each agent  $a \in A$ , an action set  $\mathbb{U}_a$  equipped with a  $\sigma$ -field  $\mathcal{U}_a$
  - for each agent  $a \in A$ , an information field

$$\mathfrak{I}_{a} \subset \mathfrak{H} = \mathfrak{U}_{A} \otimes \mathfrak{F} = \bigotimes_{b \in A} \mathfrak{U}_{b} \otimes \mathfrak{F}$$

- ► Players, who hold objectives and beliefs a partition (A<sub>p</sub>)<sub>p∈P</sub> of the set A of agents
  - for each player  $p \in P$ , a criterion

$$j_p: \mathbb{H} = \mathbb{U}_A \times \Omega = \prod_{b \in A} \mathbb{U}_b \times \Omega \to \mathbb{R}$$

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• for each player  $p \in P$ , a probability  $\mathcal{P}_p$  over  $(\Omega, \mathcal{F})$ 

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Witsenhausen intrinsic model with Nature and two players, each made of a single agent

We lay out

- basic sets
  - decision sets
  - states of Nature
  - history set

and their  $\sigma$ -fields

- objective functions
- beliefs
- information  $\sigma$ -fields, admissible strategies and predecessors

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# Nature's moves and agents decisions

- Let Ω be a measurable set equipped with a σ-field F which represents all uncertainties: any ω ∈ Ω is called a state of Nature
- ► The agent a makes one decision u<sub>a</sub> ∈ U<sub>a</sub> where the decision set U<sub>a</sub> is equipped with a σ-field U<sub>a</sub>
- ► The agent b makes one decision u<sub>b</sub> ∈ U<sub>b</sub> where the decision set U<sub>b</sub> is equipped with a σ-field U<sub>b</sub>

#### History space

The history space is the product space

 $\mathbb{H} = \mathbb{U}_{a} \times \mathbb{U}_{b} \times \Omega$ 

equipped with the product history field

 $\mathcal{H} = \mathcal{U}_{a} \otimes \mathcal{U}_{b} \otimes \mathcal{F}$ 

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## Players, criteria and beliefs

From now on, we consider the partition {a}, {b} of players, and we identify player {a} with agent a, and player {b} with agent b
The two players a, b have a criterion,

#### $j_{a}: \mathbb{U}_{a} \times \mathbb{U}_{b} \times \Omega \to \mathbb{R} , \ j_{b}: \mathbb{U}_{a} \times \mathbb{U}_{b} \times \Omega \to \mathbb{R}$

that are measurable functions over history  $\mathbb H$ 

► The two players *a*, *b* have a belief,

 $\mathcal{P}_{a}: \mathfrak{F} \rightarrow [0,1], \ \mathcal{P}_{b}: \mathfrak{F} \rightarrow [0,1]$ 

that are probability distributions over  $(\Omega, \mathcal{F})$ 

Information and predecessors

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## Information

 When making a decision, agent a and agent b can make use of information, materialized under the form of σ-fields

The information field J<sub>a</sub> of the agent a is a subfield of the history field H

 $\mathfrak{I}_{a}\subset\mathfrak{U}_{a}\otimes\mathfrak{U}_{b}\otimes\mathfrak{F}$ 

The information field J<sub>b</sub> of the agent b is a subfield of the history field H

 $\mathfrak{I}_b\subset\mathfrak{U}_a\otimes\mathfrak{U}_b\otimes\mathfrak{F}$ 

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Absence of "self-information"

The information fields J<sub>a</sub> and J<sub>b</sub> display the absence of "self-information" when

 $\mathbb{J}_{a} \subset \{\emptyset, \mathbb{U}_{a}\} \otimes \mathbb{U}_{b} \otimes \mathfrak{F}$ 

 $\mathbb{J}_b \subset \mathbb{U}_a \otimes \{ \emptyset, \mathbb{U}_b \} \otimes \mathfrak{F}$ 

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 In what follows, we always assume absence of "self-information" (otherwise, we would be led to paradoxes)

## Classical information patterns in game theory

Two agents: the principal Pr (leader) and the agent Ag (follower)

 Moral hazard (the insurance company cannot observe if the insured plays with matches at home)

 $\mathfrak{I}_{\mathtt{Pr}} \subset \{\emptyset, \mathbb{U}_{\mathtt{Ag}}\} \otimes \{\emptyset, \mathbb{U}_{\mathtt{Pr}}\} \otimes \mathfrak{F}$ 

Stackelberg leadership model

 $\mathbb{J}_{\mathtt{Ag}} \subset \{\emptyset, \mathbb{U}_{\mathtt{Ag}}\} \otimes \mathbb{U}_{\mathtt{Pr}} \otimes \mathfrak{F} \,, \ \mathbb{J}_{\mathtt{Pr}} \subset \{\emptyset, \mathbb{U}_{\mathtt{Ag}}\} \otimes \{\emptyset, \mathbb{U}_{\mathtt{Pr}}\} \otimes \mathfrak{F}$ 

 Adverse selection (the insurance company cannot observe if the insured has good health)

 $\{\emptyset,\mathbb{U}_{\mathtt{Ag}}\}\otimes\{\emptyset,\mathbb{U}_{\mathtt{Pr}}\}\otimes \mathfrak{F}\subset \mathtt{J}_{\mathtt{Ag}}\;,\;\;\mathtt{J}_{\mathtt{Pr}}\subset \mathtt{\mathcal{U}}_{\mathtt{Ag}}\otimes\{\emptyset,\mathbb{U}_{\mathtt{Pr}}\}\otimes\{\emptyset,\Omega\}$ 

Signaling

 $\{\emptyset,\mathbb{U}_{\mathtt{Ag}}\}\otimes\{\emptyset,\mathbb{U}_{\mathtt{Pr}}\}\otimes \mathfrak{F}\subset \mathtt{J}_{\mathtt{Ag}}\;,\;\;\mathtt{J}_{\mathtt{Pr}}=\mathtt{U}_{\mathtt{Ag}}\otimes\{\emptyset,\mathbb{U}_{\mathtt{Pr}}\}\otimes\{\emptyset,\Omega\}$ 

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## Cylindric subfields

Information only carried by the moves of Nature

 $\mathcal{H}_{\emptyset} = \{\emptyset, \mathbb{U}_{a}\} \otimes \{\emptyset, \mathbb{U}_{b}\} \otimes \mathcal{F}$ 

 Information only carried by the moves of Nature and by the decisions of agent a

 $\mathcal{H}_{\{a\}} = \mathcal{U}_a \otimes \{\emptyset, \mathbb{U}_b\} \otimes \mathcal{F}$ 

 Information only carried by the moves of Nature and by the decisions of agent b

 $\mathcal{H}_{\{b\}} = \{\emptyset, \mathbb{U}_a\} \otimes \mathcal{U}_b \otimes \mathcal{F}$ 

 Information carried by the moves of Nature and by the decisions of agents a and b

 $\mathcal{H}_{\{a,b\}} = \mathcal{U}_a \otimes \mathcal{U}_b \otimes \mathcal{F} = \mathcal{H}$ 

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## Definition of predecessor, excluding Nature

Consider a subset B of  $\{a, b\} \longrightarrow B \in \{\emptyset, \{a\}, \{b\}, \{a, b\}\}$  — and define

$$\mathcal{H}_{B} = \prod_{c \in B} \mathcal{U}_{c} \otimes \prod_{c \notin B} \{\emptyset, \mathbb{U}_{c}\} \otimes \mathcal{F}$$

#### Predecessor

For any agent  $c \in \{a, b\}$ , we define  $\langle c \rangle_{\mathfrak{P}}$ as the intersection of all subsets *B* of  $\{a, b\}$  such that  $\mathfrak{I}_c \subset \mathfrak{H}_B$ 

$$\langle c \rangle_{\mathfrak{P}} = \bigcap_{B, \ \mathfrak{I}_{c} \subset \mathfrak{H}_{B}} B$$

When non empty, an element of  $\langle c \rangle_{\mathfrak{P}}$  is called a predecessor of c

- Nature has no predecessor: Nature plays before the agents (but is not necessarily revealed to the agents)
- As an illustration, absence of "self-information" is equivalent to c ∉ ⟨c⟩<sub>𝔅</sub>, for any c ∈ {a, b}

## Sequential and non-sequential information patterns

#### Sequential patterns

- When  $\langle a \rangle_{\mathfrak{P}} = \emptyset$  and  $\langle b \rangle_{\mathfrak{P}} = \emptyset$ , agent *a* and agent *b* both play first (static team)
- ▶ When  $\langle a \rangle_{\mathfrak{P}} = \emptyset$  and  $\langle b \rangle_{\mathfrak{P}} = \{a\}$ , agent *a* plays first, agent *b* plays second
- When ⟨a⟩<sub>𝔅</sub> = {b} and ⟨b⟩<sub>𝔅</sub> = ∅, agent b plays first, agent a plays second
- Non-sequential pattern
  - When ⟨a⟩<sub>𝔅</sub> = {b} and ⟨b⟩<sub>𝔅</sub> = {a}, agent a and agent b
    - can be in a deadlock (non causal system)
    - or can be first and second agents depending on Nature's move (causal system)

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Strategies and admissible strategies

### Pure strategies

A (pure) strategy the agent a is a measurable mapping  $\lambda_a: \mathbb{U}_a \times \mathbb{U}_b \times \Omega \to \mathbb{U}_a, \ \lambda_a^{-1}(\mathcal{U}_a) \subset \mathcal{H}$ and the set of strategies of agent a is  $\Lambda_a = \{\lambda_a : (\mathbb{H}, \mathcal{H}) \to (\mathbb{U}_a, \mathcal{U}_a) \mid \lambda_a^{-1}(\mathcal{U}_a) \subset \mathcal{H} \}$ A (pure) strategy of agent b is a measurable mapping  $\lambda_b: \mathbb{U}_a \times \mathbb{U}_b \times \Omega \to \mathbb{U}_b, \ \lambda_b^{-1}(\mathfrak{U}_b) \subset \mathcal{H}$ and the set of strategies of agent b is  $\Lambda_a = \{\lambda_a : (\mathbb{H}, \mathcal{H}) \to (\mathbb{U}_a, \mathcal{U}_a) \mid \lambda_{\scriptscriptstyle h}^{-1}(\mathcal{U}_a) \subset \mathcal{H} \}$ 

We denote the set of strategies of all agents in A by

$$\Lambda_A = \Lambda_a \times \Lambda_b$$

### Mixed strategies

- A mixed strategy (or randomized strategy) for agent a is an element of Δ(Λ<sub>a</sub>), the set of probability distributions over the set of strategies of agent a
- A mixed strategy (or randomized strategy) for agent b is an element of Δ(Λ<sub>b</sub>), the set of probability distributions over the set of strategies of agent b
- We denote the set of mixed strategies of players by

 $\Delta(\Lambda_a) imes \Delta(\Lambda_b) \subset \Delta(\Lambda_a imes \Lambda_b)$ 

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We introduce admissible strategies to account for the interplay between decision and information

Information is the fuel of strategies

Admissible strategy

An admissible strategy of the agent  $c \in \{a, b\}$  is a mapping

 $\lambda_c: \mathbb{U}_a \times \mathbb{U}_b \times \Omega \to \mathbb{U}_c$  such that  $\lambda_c^{-1}(\mathbb{U}_c) \subset \mathfrak{I}_c$ 

▶ The set of admissible strategies of the agent  $c \in \{a, b\}$  is

 $\Lambda_c^{ad} = \{\lambda_c \mid \mathbb{U}_a \times \mathbb{U}_b \times \Omega \to \mathbb{U}_c \ , \ \lambda_c^{-1}(\mathfrak{U}_c) \subset \mathfrak{I}_c\}$ 

The set of admissible strategies is

 $\Lambda^{ad} = \Lambda^{ad}_a imes \Lambda^{ad}_b$ 

The set of mixed admissible strategies is

 $\Delta\big(\Lambda^{\mathit{ad}}_{\mathit{a}}\big)\times \Delta\big(\Lambda^{\mathit{ad}}_{\mathit{b}}\big)\subset \Delta\big(\Lambda^{\mathit{ad}}_{\mathit{a}}\times \Lambda^{\mathit{ad}}_{\mathit{b}}\big)$ 

Absence of "self-information" and structure of admissible strategies

> The information fields J<sub>a</sub> and J<sub>b</sub> display the absence of "self-information" when

$$\mathbb{J}_{a} \subset \{\emptyset, \mathbb{U}_{a}\} \otimes \mathbb{U}_{b} \otimes \mathfrak{F} \iff a 
ot \in \langle a 
angle_{\mathfrak{P}}$$

 $\mathbb{J}_b \subset \mathbb{U}_a \otimes \{ \emptyset, \mathbb{U}_b \} \otimes \mathfrak{F} \iff b \not\in \langle b \rangle_{\mathfrak{P}}$ 

When σ-fields include singletons and we exclude "self-information", then, for any admissible strategy λ<sub>c</sub> of the agent c ∈ {a, b}, we have that the expression λ<sub>c</sub>(u<sub>a</sub>, u<sub>b</sub>, ω) does not depend on u<sub>c</sub>:

$$\lambda_{a}(\mathscr{Y}_{a}, u_{b}, \omega) = \widetilde{\lambda_{a}}(u_{b}, \omega), \quad \lambda_{b}(u_{a}, \mathscr{Y}_{b}, \omega) = \widetilde{\lambda_{b}}(u_{a}, \omega)$$

# Sequential patterns and structure of admissible strategies

• When  $\langle a 
angle_{\mathfrak{P}} = \emptyset$  and  $\langle b 
angle_{\mathfrak{P}} = \emptyset$ 

$$\lambda_{a}(\mathscr{y}_{a}, \mathscr{y}_{b}, \omega) = \widetilde{\lambda_{a}}(\omega) , \ \lambda_{b}(\mathscr{y}_{a}, \mathscr{y}_{b}, \omega) = \widetilde{\lambda_{b}}(\omega)$$

► When 
$$\langle a \rangle_{\mathfrak{P}} = \emptyset$$
 and  $\langle b \rangle_{\mathfrak{P}} = \{a\}$   
 $\lambda_a(y_a, y_b, \omega) = \widetilde{\lambda_a}(\omega), \ \lambda_b(u_a, y_b, \omega) = \widetilde{\lambda_b}(u_a, \omega)$ 

► When  $\langle a \rangle_{\mathfrak{P}} = \{b\}$  and  $\langle b \rangle_{\mathfrak{P}} = \emptyset$  $\lambda_a(\mathscr{Y}_a, u_b, \omega) = \widetilde{\lambda_a}(u_b, \omega), \ \lambda_b(\mathscr{Y}_a, \mathscr{Y}_b, \omega) = \widetilde{\lambda_b}(\omega)$ 

# Non-sequential information patterns and structure of admissible strategies

When  $\langle a \rangle_{\mathfrak{P}} = \{b\}$  and  $\langle b \rangle_{\mathfrak{P}} = \{a\}$ , agent *a* and agent *b*  $\triangleright$  can be in a deadlock

$$\lambda_{a}(\mathscr{Y}_{a}, u_{b}, \omega) = \widetilde{\lambda_{a}}(u_{b}, \omega) , \ \lambda_{b}(u_{a}, \mathscr{Y}_{b}, \omega) = \widetilde{\lambda_{b}}(u_{a}, \omega)$$

or can be first and second agents depending on Nature's move
 when Nature's move is ω<sup>+</sup>, agent a plays first, agent b plays second
 λ<sub>a</sub>(𝒫<sub>a</sub>, 𝒫<sub>b</sub>, ω<sup>+</sup>) = λ<sub>a</sub>(ω<sup>+</sup>), λ<sub>b</sub>(𝑢<sub>a</sub>, 𝒫<sub>b</sub>(𝑢<sub>a</sub>) = λ<sub>b</sub>(𝑢<sub>a</sub>, ω<sup>+</sup>)

▶ when Nature's move is  $\omega^-$ , agent *b* plays first, agent *a* plays second  $\lambda_a(\mathscr{Y}_a, u_b, \omega^-) = \widetilde{\lambda_a}(u_b, \omega^-), \ \lambda_b(\mathscr{Y}_a, \mathscr{Y}_b, \omega^-) = \widetilde{\lambda_b}(\omega^-)$ 

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# Solvability property

The information fields  $\mathcal{I}_a$  and  $\mathcal{I}_b$  display the solvability property when,

- for any couple (λ<sub>a</sub>, λ<sub>b</sub>) ∈ Λ<sup>ad</sup><sub>a</sub> × Λ<sup>ad</sup><sub>b</sub> of admissible strategies and any state of Nature ω ∈ Ω,
- ► there exists one, and only one, couple (u<sub>a</sub>, u<sub>b</sub>) ∈ U<sub>a</sub> × U<sub>b</sub> of decisions such that

$$u_{a} = \lambda_{a}(u_{a}, u_{b}, \omega)$$

 $u_b = \lambda_b(u_a, u_b, \omega)$ 

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Solvability property and solution map

#### Solution map

In case of solvability, we can define  $S_{(\lambda_{a},\lambda_{b})}(\omega)$ , for any  $\omega \in \Omega$ , by

$$S_{(\lambda_{a},\lambda_{b})}(\omega) = (u_{a}, u_{b}, \omega) \iff \begin{cases} u_{a} = \lambda_{a}(u_{a}, u_{b}, \omega) \\ u_{b} = \lambda_{b}(u_{a}, u_{b}, \omega) \end{cases}$$

Hence, we obtain a mapping called the solution map

 $S_{(\lambda_a,\lambda_b)}:\Omega \to \mathbb{U}_a \times \mathbb{U}_b \times \Omega$ 

- The solvability property holds true in the sequential cases
- The graph of  $S_{(\lambda_a,\lambda_b)}$  belongs to  $\mathfrak{I}_a \vee \mathfrak{U}_a \vee \mathfrak{I}_b \vee \mathfrak{U}_b$ .

Co-cycle property of the solution map (I)

- We suppose that ⟨a⟩<sub>𝔅</sub> = {b} and ⟨b⟩<sub>𝔅</sub> = ∅, that is, agent b plays first, agent a plays second
- ▶ We consider a couple  $(\lambda_a, \lambda_b) \in \Lambda_a^{ad} \times \Lambda_b^{ad}$  of admissible strategies

### Co-cycle property of the solution map

We have that

- ▶ the strategy  $\lambda_b$  can be identified with  $\lambda_b : \Omega \to \mathbb{U}_b$  and the partial solution map  $S_{\lambda_b} : \Omega \to \mathbb{U}_b \times \Omega$  is such that  $S_{\lambda_b}(\omega) = (\lambda_b(\omega), \omega)$
- ▶ the strategy  $\lambda_a$  can be identified with  $\lambda_a : \mathbb{U}_b \times \Omega \to \mathbb{U}_a$
- the solution map has the following co-cycle property

$$\begin{split} \mathbf{S}_{(\lambda_{a},\lambda_{b})} &= (\lambda_{a} \circ \mathbf{S}_{\lambda_{b}}, \mathbf{S}_{\lambda_{b}}) : \Omega \to \mathbb{U}_{a} \times (\mathbb{U}_{b} \times \Omega) \\ \mathbf{S}_{(\lambda_{a},\lambda_{b})}(\omega) &= (\lambda_{a} (\lambda_{b}(\omega), \omega), \lambda_{b}(\omega), \omega) , \ \forall \omega \in \Omega \end{split}$$

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Co-cycle property of the solution map (II)

The co-cycle property

$$S_{(\lambda_{a},\lambda_{b})} = (\lambda_{a} \circ S_{\lambda_{b}}, S_{\lambda_{b}})$$

is equivalent to

$$S_{(\lambda_{a},\lambda_{b})}(\omega) = (u_{a}, u_{b}, \omega) \iff \begin{cases} (u_{b}, \omega) &= S_{\lambda_{b}}(\omega) \\ u_{a} &= \lambda_{a}(u_{b}, \omega) \end{cases}$$

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Criteria composed with solution map

Costs or payoffs are

 $j_a: \mathbb{U}_a \times \mathbb{U}_b \times \Omega \to \mathbb{R}$ 

 $j_b: \mathbb{U}_a \times \mathbb{U}_b \times \Omega \to \mathbb{R}$ 

Solution map is

$$S_{(\lambda_{a},\lambda_{b})}:\Omega \to \mathbb{U}_{a} \times \mathbb{U}_{b} imes \Omega$$

The composition of criteria with the solution map provides random variables

 $j_a \circ S_{(\lambda_a,\lambda_b)} : \Omega \to \mathbb{R}$ 

 $j_b \circ S_{(\lambda_a,\lambda_b)} : \Omega \to \mathbb{R}$ 

### Pure Bayesian Nash equilibrium

We recall that player a has belief  $\mathcal{P}_a$  and player b has belief  $\mathcal{P}_b$ 

#### Bayesian Nash equilibrium

We say that the couple  $(\overline{\lambda}_a, \overline{\lambda}_b) \in \Lambda_a^{ad} \times \Lambda_b^{ad}$  of admissible strategies is a Bayesian Nash equilibrium if (in case of payoffs)

$$\begin{split} \mathcal{E}_{\mathcal{P}_{a}}\Big[j_{a}\circ S_{(\overline{\lambda}_{a},\overline{\lambda}_{b})}\Big] \geq & \mathcal{E}_{\mathcal{P}_{a}}\Big[j_{a}\circ S_{(\lambda_{a},\overline{\lambda}_{b})}\Big] \ , \ \forall\lambda_{a}\in\Lambda_{a}^{ad} \\ \\ \mathcal{E}_{\mathcal{P}_{b}}\Big[j_{b}\circ S_{(\overline{\lambda}_{a},\overline{\lambda}_{b})}\Big] \geq & \mathcal{E}_{\mathcal{P}_{b}}\Big[j_{b}\circ S_{(\overline{\lambda}_{a},\lambda_{b})}\Big] \ , \ \forall\lambda_{b}\in\Lambda_{b}^{ad} \end{split}$$

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### Mixed Bayesian Nash equilibrium

We say that the couple of mixed admissible strategies

 $\left(\overline{\mu}_{a},\overline{\mu}_{b}\right)\in\Delta\left(\Lambda_{a}^{ad}\right)\times\Delta\left(\Lambda_{b}^{ad}\right)$ 

is a Bayesian Nash equilibrium if (in case of payoffs)

$$\begin{split} & \int_{\Lambda_{a}^{ad} \times \Lambda_{b}^{ad}} \overline{\mu}_{a}(d\lambda_{a}) \otimes \overline{\mu}_{b}(d\lambda_{b}) \mathcal{E}_{\mathcal{P}_{a}} \Big[ j_{a} \circ S_{(\lambda_{a},\lambda_{b})} \Big] \geq \\ & \int_{\Lambda_{a}^{ad} \times \Lambda_{b}^{ad}} \mu_{a}(d\lambda_{a}) \otimes \overline{\mu}_{b}(d\lambda_{b}) \mathcal{E}_{\mathcal{P}_{a}} \Big[ j_{a} \circ S_{(\lambda_{a},\lambda_{b})} \Big] , \ \forall \mu_{a} \in \Delta(\Lambda_{a}^{ad}) \end{split}$$

$$\int_{\Lambda_{a}^{ad} \times \Lambda_{b}^{ad}} \overline{\mu}_{a}(d\lambda_{a}) \otimes \overline{\mu}_{b}(d\lambda_{b}) \mathcal{E}_{\mathcal{P}_{b}}\left[j_{b} \circ S_{(\lambda_{a},\lambda_{b})}\right] \geq \int_{\Lambda_{a}^{ad} \times \Lambda_{b}^{ad}} \overline{\mu}_{a}(d\lambda_{a}) \otimes \mu_{b}(d\lambda_{b}) \mathcal{E}_{\mathcal{P}_{b}}\left[j_{b} \circ S_{(\overline{\lambda}_{a},\lambda_{b})}\right], \ \forall \mu_{a} \in \Delta(\Lambda_{b}^{ad})$$

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## Principal-agent models with two players

- ► A branch of Economics studies so-called principal-agent models
- Principal-agent models display a general information structure, which can be transparently expressed thanks to Witsenhausen intrinsic model
- The model exhibits two players
  - the principal Pr (leader), makes decisions u<sub>Pr</sub> ∈ U<sub>Pr</sub>, where the set of decisions is equipped with a σ-field U<sub>Pr</sub>
  - ► the agent Ag (follower) makes decisions  $u_{Ag} \in U_{Ag}$ , where the set of decisions is equipped with a  $\sigma$ -field  $U_{Ag}$

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- and Nature, corresponding to private information (or type) of the agent Ag
  - Nature selects  $\omega \in \Omega$ ,

where  $\Omega$  is equipped with a  $\sigma$ -field  $\mathfrak F$ 

Here is the most general information structure of principal-agent models

 $\mathbb{J}_{\mathtt{Pr}} \subset \mathfrak{U}_{\mathtt{Ag}} \otimes \{ \emptyset, \mathbb{U}_{\mathtt{Pr}} \} \otimes \mathfrak{F}$ 

 $\mathbb{J}_{\mathtt{Ag}} \subset \{ \emptyset, \mathbb{U}_{\mathtt{Ag}} \} \otimes \mathbb{U}_{\mathtt{Pr}} \otimes \mathfrak{F}$ 

- By these expressions of the information fields
  - J<sub>Pr</sub> of the principal Pr (leader)
  - J<sub>Ag</sub> of the agent Ag (follower)
- we have excluded self-information, that is, we suppose that the information of a player cannot be influenced by his actions

## Classical information patterns in game theory

Now, we will make the information structure more specific

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- Stackelberg leadership model
- Moral hazard
- Adverse selection
- Signaling

## Stackelberg leadership model

- In the Stackelberg leadership model of game theory,
- the follower Ag may partly observe the action of the leader Pr

 $\mathbb{J}_{\mathtt{Ag}} \subset \{\emptyset, \mathbb{U}_{\mathtt{Ag}}\} \otimes \mathfrak{U}_{\mathtt{Pr}} \otimes \mathfrak{F}$ 

whereas the leader Pr observes at most the state of Nature

 $\mathfrak{I}_{\mathtt{Pr}} \subset \{\emptyset, \mathbb{U}_{\mathtt{Ag}}\} \otimes \{\emptyset, \mathbb{U}_{\mathtt{Pr}}\} \otimes \mathfrak{F}$ 

As a consequence, the system is sequential

- with the principal Pr as first player (leader)
- and the agent Ag as second player (follower)
- Stackelberg games can be solved by bi-level optimization, for some information structures, like when

 $\mathbb{J}_{\mathtt{Pr}} \vee \{ \emptyset, \mathbb{U}_{\mathtt{Ag}} \} \otimes \mathbb{U}_{\mathtt{Pr}} \otimes \{ \emptyset, \Omega \} \subset \mathbb{J}_{\mathtt{Ag}}$ 

### Moral hazard

- An insurance company (the principal Pr) cannot observe the efforts of the insured (the agent Ag) to avoid risky behavior
- The firm faces the hazard that insured persons behave "immorally" (playing with matches at home)
- Moral hazard (hidden action) occurs when the decisions of the agent Ag are hidden to the principal Pr

 $\mathfrak{I}_{\mathtt{Pr}} \subset \{\emptyset, \mathbb{U}_{\mathtt{Ag}}\} \otimes \{\emptyset, \mathbb{U}_{\mathtt{Pr}}\} \otimes \mathfrak{F}$ 

- In case of moral hazard, the system is sequential with the principal as first player, (which does not preclude to choose the agent as first player in some special cases, as in a static team situation)
- Moral hazard games can be solved by bi-level optimization, for some information structures

### Adverse selection

- In the absence of observable information on potential customers (the agent Ag), an insurance company (the principal Pr) offers a unique price for a contract hence screens and selects the "bad" ones
- Adverse selection occurs when
  - the agent Ag knows the state of nature (his type, or private information)

 $\{\emptyset, \mathbb{U}_{\mathtt{Ag}}\} \otimes \{\emptyset, \mathbb{U}_{\mathtt{Pr}}\} \otimes \mathfrak{F} \subset \mathfrak{I}_{\mathtt{Ag}}$ 

(the agent Ag can possibly observe the principal Pr action)
but the principal Pr does not know the state of nature

 $\mathbb{J}_{\mathtt{Pr}} \subset \mathbb{U}_{\mathtt{Ag}} \otimes \{ \emptyset, \mathbb{U}_{\mathtt{Pr}} \} \otimes \{ \emptyset, \Omega \}$ 

(the principal Pr can possibly observe the agent Ag action)

In case of adverse selection, the system may or may not be sequential

# Signaling

- In biology, a peacock signals its "good genes" (genotype) by its lavish tail (phenotype)
- In economics, a worker signals his working ability (productivity) by his educational level (diplomas)
- There is room for signaling
  - when the agent Ag knows the state of nature (his type)

 $\{\emptyset, \mathbb{U}_{\mathtt{Ag}}\} \otimes \{\emptyset, \mathbb{U}_{\mathtt{Pr}}\} \otimes \mathfrak{F} \subset \mathfrak{I}_{\mathtt{Ag}}$ 

(the agent Ag can possibly observe the principal Pr action)

 whereas the principal Pr does not know the state of nature, but the principal Pr observes the agent Ag action

 ${\tt J}_{\tt Pr}={\tt U}_{\tt Ag}\otimes\{\emptyset,{\tt U}_{\tt Pr}\}\otimes\{\emptyset,\Omega\}$ 

as the agent Ag may reveal his type by his decision which is observable by the principal Pr > The system is sequential (with the agent as first player) when

 $\mathtt{J}_{\mathtt{Ag}} = \{ \emptyset, \mathbb{U}_{\mathtt{Ag}} \} \otimes \{ \emptyset, \mathbb{U}_{\mathtt{Pr}} \} \otimes \mathtt{F}$ 

The system is non causal when

 $\{\emptyset,\mathbb{U}_{\mathtt{Ag}}\}\otimes\{\emptyset,\mathbb{U}_{\mathtt{Pr}}\}\otimes\mathfrak{F}\subsetneq\mathfrak{I}_{\mathtt{Ag}}\subset\{\emptyset,\mathbb{U}_{\mathtt{Ag}}\}\otimes\mathfrak{U}_{\mathtt{Pr}}\otimes\mathfrak{F}$ 

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# Outline of the presentation

Introduction to games in intrinsic form

Witsenhausen intrinsic model and game theory with information

Nash equilibrium with information

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Open questions

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## Stackelberg leadership model

- In the Stackelberg leadership model of game theory, we consider a leader Pr (principal) and a follower Ag (agent)
- We suppose that  $\langle \Pr \rangle_{\mathfrak{P}} = \emptyset$ , that is, leader  $\Pr$  plays first,

 $\mathfrak{I}_{\mathtt{Pr}} \subset \{\emptyset, \mathbb{U}_{\mathtt{Ag}}\} \otimes \{\emptyset, \mathbb{U}_{\mathtt{Pr}}\} \otimes \mathfrak{F}$ 

 $\mathbb{J}_{\mathtt{Ag}} \subset \{ \emptyset, \mathbb{U}_{\mathtt{Ag}} \} \otimes \mathbb{U}_{\mathtt{Pr}} \otimes \mathbb{F}$ 

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### We work on a reduced history space

- ▶ As both information fields  $J_{Pr} \subset \{\emptyset, \mathbb{U}_{Ag}\} \otimes \{\emptyset, \mathbb{U}_{Pr}\} \otimes \mathcal{F}$ and  $J_{Ag} \subset \{\emptyset, \mathbb{U}_{Ag}\} \otimes \mathcal{U}_{Pr} \otimes \mathcal{F}$ — do not depend on  $\mathcal{U}_{Ag}$ , the actions of the follower Ag (agent) do not fuel strategies (via information), so that we introduce
- ▶ the reduced history space  $\widetilde{\mathbb{H}}$  (without the actions of the follower Ag) equipped with the reduced history field  $\widetilde{\mathcal{H}}$

$$\widetilde{\mathbb{H}} = \mathbb{U}_{\mathtt{Pr}} \times \Omega \;, \;\; \widetilde{\mathcal{H}} = \mathcal{U}_{\mathtt{Pr}} \otimes \mathfrak{F}$$

▶ and the reduced information fields  $\widetilde{J}_{Pr}$  and  $\widetilde{J}_{Ag}$  defined by

$$\begin{array}{lll} \mathbb{J}_{\mathtt{Pr}} = & \{ \emptyset, \mathbb{U}_{\mathtt{Ag}} \} \otimes \widetilde{\mathbb{J}}_{\mathtt{Pr}} & \mbox{ with } \widetilde{\mathbb{J}}_{\mathtt{Pr}} \subset \{ \emptyset, \mathbb{U}_{\mathtt{Pr}} \} \otimes \mathfrak{F} \subset \widetilde{\mathcal{H}} \\ \mathbb{J}_{\mathtt{Ag}} = & \{ \emptyset, \mathbb{U}_{\mathtt{Ag}} \} \otimes \widetilde{\mathbb{J}}_{\mathtt{Ag}} & \mbox{ with } \widetilde{\mathbb{J}}_{\mathtt{Ag}} \subset \mathcal{U}_{\mathtt{Pr}} \otimes \mathfrak{F} = \widetilde{\mathcal{H}} \end{array}$$

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# Here is what become the admissible strategies on the reduced history space

We consider a couple  $(\lambda_{Ag}, \lambda_{Pr}) \in \Lambda_{Ag}^{ad} \times \Lambda_{Pr}^{ad}$  of admissible strategies

► As  $J_{Pr} = \{\emptyset, \mathbb{U}_{Ag}\} \otimes \widetilde{J}_{Pr}$  with  $\widetilde{J}_{Pr} \subset \{\emptyset, \mathbb{U}_{Pr}\} \otimes \mathcal{F}$ , the strategy  $\lambda_{Pr}$  of the leader Pr can be identified with

 $\widetilde{\lambda}_{\mathtt{Pr}}:\Omega o \mathbb{U}_{\mathtt{Pr}}$ 

(indeed, the strategies of the leader Pr depend at most upon Nature)
As J<sub>Ag</sub> = {Ø, U<sub>Ag</sub>} ⊗ J̃<sub>Ag</sub>, with J̃<sub>Ag</sub> ⊂ U<sub>Pr</sub> ⊗ F, the strategy λ<sub>Ag</sub> of the follower Ag can be identified with

 $\widetilde{\lambda}_{\mathtt{Ag}}: \mathbb{U}_{\mathtt{Pr}} imes \Omega o \mathbb{U}_{\mathtt{Ag}}$ 

Therefore, we can work with reduced admissible strategies

 $(\widetilde{\lambda}_{\mathtt{Ag}},\widetilde{\lambda}_{\mathtt{Pr}})\in\widetilde{\Lambda}_{\mathtt{Ag}}^{ad} imes\widetilde{\Lambda}_{\mathtt{Pr}}^{ad}$ 

Strategy independence of conditional expectation (SICE)

#### Assumption SICE

There exists a probability  $\mathbb{Q}$  on  $\widetilde{\mathbb{H}} = \mathbb{U}_{Pr} \times \Omega$  such that

 $\mathcal{P}_{\mathtt{Ag}} \circ S_{\widetilde{\lambda}_{\mathtt{Pr}}}^{-1} = \mathcal{T}_{\widetilde{\lambda}_{\mathtt{Pr}}} \mathbb{Q} \text{ with } \mathcal{E}_{\mathbb{Q}} \big[ \mathcal{T}_{\widetilde{\lambda}_{\mathtt{Pr}}} \mid \widetilde{J}_{\mathtt{Ag}} \big] > 0 \;, \; \forall \widetilde{\lambda}_{\mathtt{Pr}} \in \widetilde{\Lambda}_{\mathtt{Pr}}^{ad}$ 

and that, under  $\mathbb{Q}$ , the conditional expected gain of the follower Ag does not change when one adds to his information both the actions and the information available of the leader Pr, namely

 $\mathcal{E}_{\mathbb{Q}}\big[j_{\mathtt{Ag}}(u_{\mathtt{Ag}},\cdot)\mid\widetilde{\mathbb{J}}_{\mathtt{Ag}}\big]=\mathcal{E}_{\mathbb{Q}}\big[j_{\mathtt{Ag}}(u_{\mathtt{Ag}},\cdot)\mid\widetilde{\mathbb{J}}_{\mathtt{Ag}}\vee\widetilde{\mathbb{J}}_{\mathtt{Pr}}\vee\widetilde{\mathbb{D}}_{\mathtt{Pr}}\big]\;,\;\;\forall u_{\mathtt{Ag}}\in\mathbb{U}_{\mathtt{Ag}}$ 

Bayesian Nash equilibria can be obtained by bi-level optimization under assumption SICE

Suppose assumption SICE holds true

The (upper level) optimization problem for the follower Ag

 $\min_{u_{\mathrm{Ag}} \in \mathbb{U}_{\mathrm{Ag}}} \mathcal{E}_{\mathbb{Q}} \big[ j_{\mathrm{Ag}}(u_{\mathrm{Ag}}, \cdot) \mid \widetilde{\mathbb{J}}_{\mathrm{Ag}} \big]$ 

provides (under technical assumptions, by a measurable selection theorem) an  $\widetilde{J}_{Ag}$ -measurable solution

$$\overline{\widetilde{\lambda}}_{\mathtt{Ag}}:\mathbb{U}_{\mathtt{Pr}} imes\Omega o\mathbb{U}_{\mathtt{Ag}}\;,\;\;\sigma(\overline{\widetilde{\lambda}}_{\mathtt{Ag}})\subset\widetilde{\mathbb{J}}_{\mathtt{Ag}}$$

Then, the (lower level) optimization problem for the leader Pr is

$$\min_{\widetilde{\lambda}_{\mathtt{Pr}} \in \widetilde{\Lambda}_{\mathtt{Pr}}^{sd}} \mathcal{E}_{\mathcal{P}_{\mathtt{Pr}}} \left[ j_{\mathtt{Pr}} \circ S_{(\overline{\widetilde{\lambda}}_{\mathtt{Ag}}, \widetilde{\lambda}_{\mathtt{Pr}})} \right]$$

Here is what becomes the solution map on the reduced history space

> By sequentiality, the solution map S<sub>(λ<sub>Ag</sub>,λ<sub>Pr</sub>)</sub> satisfies the co-cycle property

> > $S_{(\lambda_{\mathtt{Ag}},\lambda_{\mathtt{Pr}})} = (\lambda_{\mathtt{Ag}} \circ S_{\lambda_{\mathtt{Pr}}}, S_{\lambda_{\mathtt{Pr}}}) = (\lambda_{\mathtt{Ag}}, \mathrm{Id}_{\mathbb{U}_{\mathtt{Pr}} imes \Omega}) \circ S_{\lambda_{\mathtt{Pr}}}$

• If we introduce a reduced solution map  $S_{\widetilde{\lambda}_{Pr}} = (\widetilde{\lambda}_{Pr}, \mathrm{Id}_{\Omega})$ 

$$\Omega \xrightarrow{S_{\widetilde{\lambda}_{\Pr}}} \mathbb{U}_{\Pr} \times \Omega , \ \omega \mapsto \left( \widetilde{\lambda}_{\Pr}(\omega), \omega \right) ,$$

we can now write  $S_{(\lambda_{Ag},\lambda_{Pr})} = (\widetilde{\lambda}_{Ag}, \mathrm{Id}_{\mathbb{U}_{Pr} \times \Omega}) \circ S_{\widetilde{\lambda}_{Pr}}$ , that is,

$$\mathcal{S}_{(\lambda_{\mathtt{Ag}},\lambda_{\mathtt{Pr}})}:\Omega\xrightarrow{\mathcal{S}_{\widetilde{\lambda}_{\mathtt{Pr}}}}\mathbb{U}_{\mathtt{Pr}}\times\Omega\xrightarrow{(\widetilde{\lambda}_{\mathtt{Ag}},\mathrm{Id}_{\mathbb{U}_{\mathtt{Pr}}}\times\Omega)}\mathbb{U}_{\mathtt{Ag}}\times\mathbb{U}_{\mathtt{Pr}}\times\Omega$$

that is,

$$S_{(\lambda_{\mathtt{Ag}},\lambda_{\mathtt{Pr}})}:\omega\mapsto (\widetilde{\lambda}_{\mathtt{Pr}}(\omega),\omega)\mapsto (\widetilde{\lambda}_{\mathtt{Ag}}(\widetilde{\lambda}_{\mathtt{Pr}}(\omega),\omega),\widetilde{\lambda}_{\mathtt{Pr}}(\omega),\omega)$$

Strategy independence of conditional expectation (SICE)

#### Assumption SICE

There exists a probability  $\mathbb{Q}$  on  $\widetilde{\mathbb{H}} = \mathbb{U}_{Pr} \times \Omega$  such that

$$\mathcal{P}_{\mathtt{Ag}} \circ S^{-1}_{\widetilde{\lambda}_{\mathtt{Pr}}} = \mathcal{T}_{\widetilde{\lambda}_{\mathtt{Pr}}} \mathbb{Q} \text{ with } \mathcal{E}_{\mathbb{Q}} \big[ \mathcal{T}_{\widetilde{\lambda}_{\mathtt{Pr}}} \mid \widetilde{\mathbb{I}}_{\mathtt{Ag}} \big] > 0 \;, \; \forall \widetilde{\lambda}_{\mathtt{Pr}} \in \widetilde{\Lambda}_{\mathtt{Pr}}^{ad}$$

and that

$$\mathcal{E}_{\mathbb{Q}}\big[j_{\mathtt{Ag}}(u_{\mathtt{Ag}},\cdot)\mid\widetilde{\mathbb{J}}_{\mathtt{Ag}}\big]=\mathcal{E}_{\mathbb{Q}}\big[j_{\mathtt{Ag}}(u_{\mathtt{Ag}},\cdot)\mid\widetilde{\mathbb{J}}_{\mathtt{Ag}}\vee\widetilde{\mathbb{J}}_{\mathtt{Pr}}\vee\widetilde{\mathbb{D}}_{\mathtt{Pr}}\big]\;,\;\;\forall u_{\mathtt{Ag}}\in\mathbb{U}_{\mathtt{Ag}}$$

Under assumption SICE, we have that

$$\begin{split} \mathcal{E}_{\mathcal{P}_{\boldsymbol{a}}}\Big[j_{\boldsymbol{a}}\circ S_{(\lambda_{\boldsymbol{a}},\lambda_{\boldsymbol{b}})}\Big] &= \mathcal{E}_{\mathcal{P}_{\boldsymbol{a}}}\Big[j_{\boldsymbol{a}}\circ (\widetilde{\lambda}_{\mathtt{Ag}},\mathrm{Id}_{\mathbb{U}_{\mathtt{Pr}}\times\Omega})\circ S_{\widetilde{\lambda}_{\mathtt{Pr}}}\Big] \\ &= \mathcal{E}_{\mathcal{P}_{\mathtt{Ag}}\circ S_{\widetilde{\lambda}_{\mathtt{Pr}}}^{-1}}\Big[j_{\boldsymbol{a}}\circ (\widetilde{\lambda}_{\mathtt{Ag}},\mathrm{Id}_{\mathbb{U}_{\mathtt{Pr}}\times\Omega})\Big] \\ &= \mathcal{E}_{\mathbb{Q}}\Big[\mathcal{T}_{\widetilde{\lambda}_{\mathtt{Pr}}}j_{\boldsymbol{a}}\circ (\widetilde{\lambda}_{\mathtt{Ag}},\mathrm{Id}_{\mathbb{U}_{\mathtt{Pr}}\times\Omega})\Big] \end{split}$$

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## Bayesian Nash equilibrium under assumption SICE

#### Bayesian Nash equilibrium

Under assumption SICE, the couple  $(\widetilde{\lambda}_{Ag}, \widetilde{\lambda}_{Pr}) \in \widetilde{\Lambda}_{Ag}^{ad} \times \widetilde{\Lambda}_{Pr}^{ad}$  of reduced admissible strategies is a Bayesian Nash equilibrium if (in case of payoffs)

$$\mathcal{E}_{\mathbb{Q}}\Big[j_{\mathtt{Ag}}\circ(\overline{\widetilde{\lambda}}_{\mathtt{Ag}},\mathrm{Id}_{\mathbb{U}_{\mathrm{Pr}} imes\Omega})\Big]\geq \quad \mathcal{E}_{\mathbb{Q}}\Big[j_{\mathtt{Ag}}\circ(\widetilde{\lambda}_{\mathtt{Ag}},\mathrm{Id}_{\mathbb{U}_{\mathrm{Pr}} imes\Omega})\Big]$$
  
 $orall\widetilde{\lambda}_{\mathtt{Ag}}\in\widetilde{\Lambda}^{ad}_{\mathtt{Ag}}$ 

$$\begin{split} \mathcal{E}_{\mathcal{P}_{\mathtt{Pr}}} \Big[ j_{\mathtt{Pr}} \circ S_{(\overline{\widetilde{\lambda}}_{\mathtt{Ag}}, \overline{\widetilde{\lambda}}_{\mathtt{Pr}})} \Big] \geq & \mathcal{E}_{\mathcal{P}_{\mathtt{Pr}}} \Big[ j_{\mathtt{Pr}} \circ S_{(\overline{\widetilde{\lambda}}_{\mathtt{Ag}}, \overline{\lambda}_{\mathtt{Pr}})} \Big] \\ & \forall \widetilde{\lambda}_{\mathtt{Pr}} \in \widetilde{\Lambda}_{\mathtt{Pr}}^{ad} \end{split}$$

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There exists an optimal strategy of the follower Ag that does not depend on the leader Pr strategy

$$\begin{split} \min_{\widetilde{\lambda}_{Ag}\in\widetilde{\Lambda}_{Ag}^{ad}} \mathcal{E}_{\mathbb{Q}}\Big[j_{Ag}\circ(\widetilde{\lambda}_{Ag},\mathrm{Id}_{\mathbb{U}_{Pr}\times\Omega})\Big] &= \min_{\widetilde{\lambda}_{Ag}} \sum_{\widetilde{\lambda}_{Ag}^{-1}(\mathfrak{U}_{Ag})\subset\widetilde{\mathfrak{I}}_{Ag}} \mathcal{E}_{\mathbb{Q}}\Big[j_{Ag}\circ(\widetilde{\lambda}_{Ag},\mathrm{Id}_{\mathbb{U}_{Pr}\times\Omega})\Big] \\ &= \mathcal{E}_{\mathbb{Q}}\Big[\min_{u_{Ag}\in\mathbb{U}_{Ag}} \mathcal{E}_{\mathbb{Q}}\Big[j_{Ag}(u_{Ag},\cdot)\mid\widetilde{\mathfrak{I}}_{Ag}\Big]\Big] \end{split}$$

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### Research questions

How should we talk about games using WIM?

- Can we extend the Bayesian Nash Equilibrium concept to general risk measures?
- How does the notion of subgame perfect Nash equilibrium translate within this framework?
- WIM: game theoretical results
  - What would a Nash theorem be in the WIM setting?
  - When do we have a generalized "backward induction" mechanism?

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- Under proper sufficient conditions on the information structure (extension of perfect recall), can we restrict the search among behavioral strategies instead of mixed strategies?
- Applications of WIM
  - Can we re-organize the games bestiary using WIM?
  - Can we use the WIM framework for mechanism design?

# We obtain a Nash theorem in the WIM setting

#### Theorem

Any finite, solvable, Witsenhausen game has a mixed NE

Proof

- The set of policies is finite, as policies map the finite history set towards finite decision sets
- To each policy profile, we associate a payoff vector
- ▶ We thus obtain a matrix game and we can apply Nash theorem

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## Generalized existence result

Step one, discretization

▶ We introduce  $g_a^{(n)}$  the injection from  $\mathbb{U}_a^{(n)}$  into  $\mathbb{U}_a$ 

$$g_a^{(n)}:\mathbb{U}_a^{(n)}\hookrightarrow\mathbb{U}_a$$

▶ We introduce  $h^{(n)}$  that maps  $\mathbb H$  into  $\mathbb H^{(n)}$  with  $h^{(n)}_{\mathbb H^{(n)}} = Id_{\mathbb H^{(n)}}$ 

• 
$$(\lambda_a^{ad})^{(n)} = \{\lambda_a \in \lambda_a^{(n)}, \sigma(g_a^{(n)} \circ \lambda_a \circ h^{(n)}) \subseteq \mathcal{I}_a\}$$

Current difficulties:

- Definition of the discretization, in particular h<sup>(n)</sup>, to obtain a limit
- Continuity of the solution map

$$\Lambda_A imes \Omega o \mathbb{H}$$
,  $(\lambda, \omega) \mapsto S_\lambda(\omega)$ 

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## Behavioral vs mixed strategies

Mixed strategies are

 $\prod_{p\in P} \Delta\big(\prod_{a\in A_p} \Lambda_a^{ad}\big)$ 

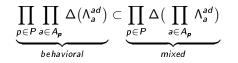
and reflect the synchronization of his agents by the player

Behavioral strategies are

 $\prod_{p \in P} \prod_{a \in A_p} \Delta(\Lambda_a^{ad})$ 

and they do not require any correlating procedure

 Under proper sufficient conditions on the information structure, we expect to prove that some games can be solved over the smaller set of behavioral strategies instead of the large set of mixed strategies



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## Applications

- The WIM is of particular interest for non sequential games
- In particular we envision applications for networks, auctions and decentralized energy systems

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## Mechanism design presented in the intrinsic framework

- The designer (= principal) can extend the natural history set, by offering new decisions to every agent (messages)
- He is free to extend the information fields of the agents as he wishes

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He can partly shape the objective functions of the players

#### Thank you :-)

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