Tutorial Perturbation-Duality Scheme

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Outline

Duality for linear programs

Perturbation-duality scheme in convex analysis

Duality for pure integer linear programs

Generalized perturbation duality scheme

Developments and examples

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Classic Lagrangian duality (the case of inequality constraints)

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Generalized perturbation duality scheme

Developments and examples

Lagrangian relaxation in LP

Chvátal dual problem for pure integer linear programming

Sensitivity analysis

Couplings for optimal transport

Illustration of the scheme in Linear Programming (LP)

- Constraint matrix $A \in \mathbb{R}^{m \times n}$
- Cost vector $k \in \mathbb{R}^n$
- Anchor $\bar{b} \in \mathbb{R}^m$

Initial minimization problem

$$\begin{array}{ll}
\inf & \langle k \mid x \rangle \\
x \in \mathbb{R}^n \\
Ax = \overline{b} \\
x \ge 0
\end{array}$$

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Step 1. Perturbation of the initial minimization problem

Embedding the problem in a family of minimization problems

$$orall b \in \mathbb{R}^m, \ arphi(b) = \inf_{\substack{x \in \mathbb{R}^n \ Ax = b \ x > 0}} \langle k \mid x
angle$$

- lntroducing a perturbation space: \mathbb{R}^m
- ▶ Perturbation function $\varphi : \mathbb{R}^m \to \overline{\mathbb{R}} = \mathbb{R} \cup \{-\infty\} \cup \{+\infty\}$

unbounded unfeasible

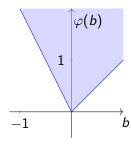
▶ Value of the initial problem: $\varphi(\bar{b})$

Examples of epigraphs of the perturbation functions for LP Example

Let $\varphi:\mathbb{R}\to\mathbb{R}$ be defined as

$$orall b\in\mathbb{R}\;,\;\;arphi(b)=\displaystyle egin{array}{cc} \inf & x_1+2x_2\ & x\in\mathbb{R}^2\ & x_1-x_2=b\ & x>0 \end{array}$$

Then $\varphi(b) = \max\{-2b, b\}$, $\forall b \in \mathbb{R}$



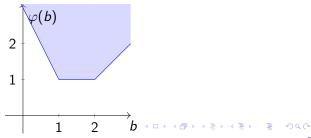
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Examples of epigraphs of the perturbation functions for LP Example

Let $\varphi:\mathbb{R}\to\mathbb{R}$ be defined as

$$arphi(b) = \inf_{\substack{x \in \mathbb{R}^3 \\ x_1 + x_2 + 3x_3 = 1 \\ x_1 + 2x_2 + 4x_3 = b \\ x \ge 0} x_1 + 2x_2 + 4x_3 = b$$

Then $\varphi(b) = \max\{3-2b, 1, b-1\}, \forall b \in \mathbb{R}$

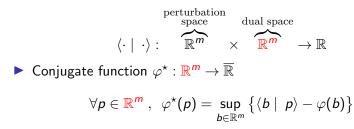


Step 2. Dual space, coupling and conjugate function

Perturbation function

$$orall b \in \mathbb{R}^m, \ arphi(b) = \inf_{egin{array}{cc} x \in \mathbb{R}^n \ Ax = b \ x \geq 0 \end{array}} \langle k \mid x
angle$$

Introducing the bilinear coupling



Conjugate function and Lagrangian

$$\varphi^{\star}(p) = \sup_{b \in \mathbb{R}^{m}} \left\{ \langle b \mid p \rangle - \varphi(b) \right\}$$

=
$$\sup_{b \in \mathbb{R}^{m}} \left\{ \langle b \mid p \rangle - \inf_{\substack{Ax=b \\ x \ge 0}} \langle k \mid x \rangle \right\}$$

=
$$\sup_{b \in \mathbb{R}^{m}} \left\{ \langle b \mid p \rangle + \sup_{\substack{Ax=b \\ x \ge 0}} \langle -k \mid x \rangle \right\}$$

=
$$\sup_{x \ge 0} \left\{ \sup_{\substack{Ax=b \\ b \in \mathbb{R}^{m}}} \langle b \mid p \rangle + \langle -k \mid x \rangle \right\}$$

=
$$\sup_{x \ge 0} \left\{ \langle Ax \mid p \rangle - \langle k \mid x \rangle \right\}$$

=
$$\langle \bar{b} \mid p \rangle - \inf_{x \ge 0} \underbrace{\left\{ \langle \bar{b} - Ax \mid p \rangle + \langle k \mid x \rangle \right\}}_{\text{Lagrangian: } \mathcal{L}(x,p)}$$

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Step 3. Biconjugate and weak duality

▶ Biconjugate function $\varphi^{\star\star'} : \mathbb{R}^m \to \overline{\mathbb{R}}$

$$orall b \in \mathbb{R}^m, \ arphi^{\star\star'}(b) = \sup_{p \in \mathbb{R}^m} \left\{ \langle b \mid p
angle - arphi^\star(p)
ight\}$$

▶ We obtain weak duality for all $b \in \mathbb{R}^m$

$$\sup_{\substack{p \in \mathbb{R}^{m} \\ p^{T}A \leq k \\ \text{dual problem}}} \langle b \mid p \rangle = \varphi^{\star\star'}(b) \leq \varphi(b) = \qquad \begin{array}{c} \inf_{\substack{x \in \mathbb{R}^{n} \\ Ax = b \\ x \geq 0 \end{array}} \langle k \mid x \rangle \\ Ax = b \\ x \geq 0 \end{array}$$
At the anchor \overline{b}

$$\varphi^{\star\star'}(\overline{b}) = \sup_{p \in \mathbb{R}^{m}} \left\{ \langle \overline{b} \mid p \rangle - \varphi^{\star}(p) \right\}$$

Step 4. Conditions for strong duality

Proposition

Let $A \in \mathbb{R}^{m \times n}$ and $k \in \mathbb{R}^n$ If there exists $\bar{b} \in \mathbb{R}^m$ such that $\varphi(\bar{b}) \in \mathbb{R}$

(the corresponding LP is bounded and feasible)

then for all $b \in \mathbb{R}^m$

$$\begin{pmatrix} \sup_{\substack{p \in \mathbb{R}^m \\ p^T A \leq k}} \langle b \mid p \rangle \\ = \end{pmatrix} \underbrace{\varphi^{\star \star'}(b) = \varphi(b)}_{\text{strong duality}} \begin{pmatrix} \inf_{\substack{x \in \mathbb{R}^n \\ Ax = b \\ x \geq 0}} \langle k \mid x \rangle \\ = \underbrace{f_{X \in \mathbb{R}^n}}_{X \geq 0} \begin{pmatrix} e_{X \in \mathbb{R}^n} \\ e_{X \geq 0} \end{pmatrix}$$

Remark

1. $\varphi(b)$ can take the value $+\infty$ 2. If there is no $\overline{b} \in \mathbb{R}^m$ such that $\varphi(\overline{b}) \in \mathbb{R}$ then $\varphi^{**'}(b) = -\infty, \ \forall b \in \mathbb{R}^m$ $\varphi(b) = \begin{cases} -\infty, \ \forall b \in \text{dom } \varphi \\ +\infty, \ otherwise \end{cases}$

Proof of strong duality for LP. Sketch of the proof

We assume there is a bounded and feasible program

 $\exists b \in \mathbb{R}^m, \ \varphi(b) \in \mathbb{R}$

We use the properness lemma to show that every program is bounded

φ is proper

We show that epi φ is a closed convex set (by showing that epi φ is a polyhedron)

 φ is a closed convex function

We apply Rockafellar's result [Rockafellar, 1974, Theorem 5] to get strong duality

 $\varphi^{\star\star} = \varphi$

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Proof of strong duality for LP. Proper functions

Definition

- Let $f : \mathbb{R}^m \to \overline{\mathbb{R}} = \mathbb{R} \cup \{-\infty, +\infty\}$
 - $\blacktriangleright \text{ dom } f = \{b \in \mathbb{R}^m : f(b) < +\infty\}$
 - The function f is said to be proper if dom f ≠ Ø and -∞ < f(b), ∀b ∈ ℝ^m

Lemma If there is $\bar{b} \in \mathbb{R}^m$ such that $-\infty < \varphi(\bar{b})$

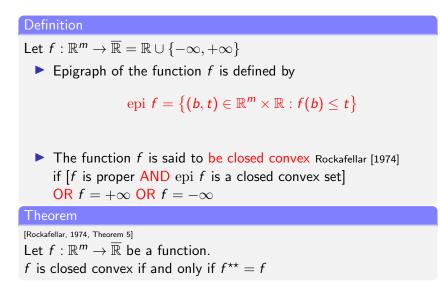
the corresponding LP is bounded

then the value function φ is proper

all LPs are bounded

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Proof of strong duality for LP. Closed convex functions



Proof of strong duality for LP. Argmin lemma

Lemma

Let $b \in \mathbb{R}^m$ such that $\varphi(b) \in \mathbb{R}$ then for $t \in \mathbb{R}$

$$\varphi(b) \leq t \iff \exists x \in \mathbb{R}^n \text{ s.t. } \underbrace{Ax = b, x \geq 0}_{x \text{ is a feasible point}} \text{ and } \langle k \mid x \rangle \leq t$$

Remark This Lemma is also true for $b \in \mathbb{R}^m$ such that $\varphi(b) = -\infty$ TODO proof using Minkowski-Weyl's theorem [?, Theorem 3.52]

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Proof of strong duality for LP. $epi \ \varphi$ is a polyhedron

Proposition

Let $A \in \mathbb{R}^{m \times n}$ and $k \in \mathbb{R}^n$ define the value function $\varphi : \mathbb{R}^m \to \overline{\mathbb{R}}$ by

$$orall b \in \mathbb{R}^m, \ arphi(b) = \inf_{egin{array}{cc} x \in \mathbb{R}^n \ Ax = b \ x \geq 0 \end{array}} \langle k \mid x
angle$$

Then $epi \varphi$ is a polyhedron

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Proof that $epi \varphi$ is a polyhedron $A \in \mathbb{R}^{m \times n}, k \in \mathbb{R}^n$ Let $b \in \mathbb{R}^m$, we assume that $\varphi(b) < +\infty$ $\varphi(b) < t$ $\iff \inf_{\substack{x \in \mathbb{R}^n \\ Ax = b \\ x > 0}} \langle k \mid x \rangle \le t$ Using argmin lemma $\iff \min_{x \in \mathbb{R}^n} \langle k \mid x \rangle \le t$ $A\bar{x}=b$ $x \ge 0$ $\iff \exists x \in \mathbb{R}^n \text{ s.t. } Ax = b , \ x \ge 0 , \ \langle k \mid x \rangle - t \le 0$ $\iff \operatorname{epi} \varphi = \pi_{(b,t)} \Big\{ (b,t,x) \in \mathbb{R}^m \times \mathbb{R} \times \mathbb{R}^n : \begin{cases} Ax = b \\ x \ge 0 \\ \langle k \mid x \rangle - t < 0 \end{cases} \Big\}$

Thus $epi \varphi$ is the projection of a polyhedron. So, $epi \varphi$ is a polyhedron.

Proof of strong duality for LP. Sketch of the proof

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 $\varphi^{\star\star} = \varphi$

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Remark on weak and strong duality

- Weak duality is a global notion on a family of minimization problem and a family of maximization problems
- Strong duality is a local notion between a minimization problem and a maximization problem produced by the perturbation duality scheme

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Example when strong duality is not achieve for LP

$$-\infty = \begin{pmatrix} \sup & p_1 + 2p_2 \\ p \in \mathbb{R}^2 & \\ p_1 + p_2 = -1 & \\ p_1 + p_2 = 0 & \end{pmatrix}$$

 $= \varphi^{\star\star}(1,2) < \varphi(1,2) =$

$$\begin{pmatrix} \inf & -x_1 \\ x \in \mathbb{R}^2 \\ x_1 + x_2 = 1 \\ x_1 + x_2 = 2 \\ x \ge 0 \end{pmatrix} = +\infty$$

Outline of the presentation Duality for linear programs

Duality for linear programs with the PDS

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- Sensitivity analysis
- Couplings for optimal transport

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Summary of the perturbation-duality scheme for LP Rockafellar [1974]

1. We perturb a minimization problem

$$\forall b \in \mathbb{R}^{m}, \ \varphi(b) = \inf_{\substack{x \in \mathbb{R}^{n} \\ Ax = b \\ x \ge 0}} \langle k \mid x \rangle$$

2. We pair the primal space \mathbb{R}^m and a dual space \mathbb{R}^m

$$\langle \cdot | \cdot \rangle : \mathbb{R}^m \times \mathbb{R}^m \to \mathbb{R}$$

3. We biconjugate the perturbation function φ

$$\begin{pmatrix} \sup_{\substack{p \in \mathbb{R}^m \\ p^T A \leq k}} & e^{\star \star'}(b) \leq \varphi(b), \quad \forall b \in \mathbb{R}^m \\ \forall b \in \mathbb{R}^m \end{pmatrix}$$
weak duality is guaranteed

4. Strong duality at the anchor \bar{b} when φ is proper or $\bar{b} \in \text{dom } \varphi$

Back to the coupling

Bilinear coupling
$$\langle \cdot | \cdot \rangle$$
: $\mathbb{R}^m \times \mathbb{R}^m \to \mathbb{R}$
 \mathbb{R}^m can be identified to the functional space $L^m = \{F : \mathbb{R}^m \to \mathbb{R} | F \text{ is linear}\}$, thus the coupling $p \in \mathbb{R}^m \leftrightarrow F \in L^m$
 $\langle b | p \rangle \leftrightarrow F(b)$
Conjugate function $\varphi^* : L^m \to \mathbb{R}$

$$orall F \in L^m \;, \; arphi^\star(F) = \sup_{b \in \mathbb{R}^m} ig\{F(b) - arphi(b)ig\}$$

Where we stand and where we go

- We have illustrated the perturbation duality scheme (PDS) on Linear Programs
- We will present the PDS in the convex case
- Then we will present the PDS in the generalized convex case

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The Fenchel conjugacy

Definition

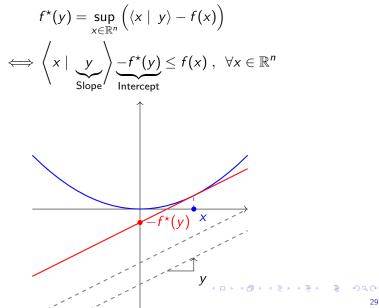
Two vector spaces \mathcal{X} and \mathcal{Y} , paired by a bilinear form $\langle \cdot | \cdot \rangle$ (in the sense of convex analysis), give rise to the classic Fenchel conjugacy

$$f^{\star}(y) = \sup_{x \in \mathcal{X}} \left(\langle x \mid y \rangle - f(x) \right), \ \forall y \in \mathcal{Y}$$

for any function $f : \mathcal{X} \to \overline{\mathbb{R}}$

Fenchel conjugate	Fourier transform				
sup o +					
$+ \rightarrow \times$					
$\sup_{x \in \mathcal{X}} \left(\langle x \mid y \rangle - f(x) \right)$	$\int_{\mathcal{X}} e^{\langle x \mid y \rangle} f(x) dx$				

Representing Fenchel conjugacy For $y \in \mathbb{R}^n$



The biconjugate function is a minorant of the function

Definition

Let $f : \mathcal{X} \to \overline{\mathbb{R}}$ be a function Its biconjugate $f^{\star\star'} : \mathcal{X} \to \overline{\mathbb{R}}$ is defined by

$$f^{\star\star'}(x) = \sup_{y\in\mathcal{Y}} \langle x\mid y
angle - f^{\star}(y)$$

Proposition

Let $f : \mathcal{X} \to \overline{\mathbb{R}}$ be a function Then $f^{\star\star'} : \mathbb{R}^m \to \overline{\mathbb{R}}$ satisfies

 $f^{\star\star'} \leq f$

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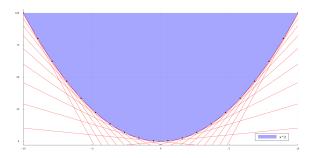
Fenchel-Moreau Theorem

[Bauschke and Combettes, 2017, Theorem 13.37]

Theorem

Let $(\mathcal{H}, \langle \cdot | \cdot \rangle)$ be a real Hilbert space Let $f : \mathcal{H} \to] - \infty, +\infty]$ proper Then f is lower semicontinous and convex if and only if

$$f = f^{\star\star'}$$



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Dual problems: perturbation scheme [Rockafellar, 1974]

sets	optimization	primal	pairing	dual
	set W	space $\mathbb U$	$\mathbb{U} \stackrel{\langle \cdot, \cdot \rangle}{\leftrightarrow} \mathbb{V}$	space $\mathbb V$
variables	decision	perturbation	$\langle u, v \rangle$	sensitivity
	$w \in \mathbb{W}$	$u \in \mathbb{U}$	$\in \mathbb{R}$	$m{v}\in\mathbb{V}$
bivariate		Rockafellian		Lagrangian
functions		$\mathcal{R}:\mathbb{W}\times\mathbb{U}\to\overline{\mathbb{R}}$		$\mathcal{L}:\mathbb{W} imes\mathbb{V} o\overline{\mathbb{R}}$
definition				$\mathcal{L}(w, v) =$
				$\inf_{u\in\mathbb{U}} \left\{ \mathcal{R}(w,u) - \langle u, v \rangle \right\}$
property				$-\mathcal{L}(w,\cdot) = ig(\mathcal{R}(w,\cdot)ig)^{\star}$
property				$-\mathcal{L}(w,\cdot)$ is \star' -convex
				(hence $\mathcal{L}(w, \cdot)$ is concave usc)
univariate		perturbation function		dual function
functions		$\varphi:\mathbb{U}\to\overline{\mathbb{R}}$		$\psi: \mathbb{V} \to \overline{\mathbb{R}}$
definition		$\varphi(u) = \inf_{w \in \mathbb{W}} \mathcal{R}(w, u)$		$\psi(v) = \inf_{w \in \mathbb{W}} \mathcal{L}(w, v)$
property				$-\psi = \varphi^{\star}$

Anchor $0 \in \mathcal{X}$ and dual maximization problem (weak duality) $\varphi^{\star\star'}(0) = \sup_{y \in \mathcal{Y}} \{-\psi(y)\} \le \inf_{w \in \mathbb{W}} \mathcal{R}(w, 0) = \varphi(0)$

Strong duality iff φ is *-convex at 0 iff $\varphi^{**'}(0) = \varphi(0)$

Dual problems given by Fenchel conjugacy

▶ Set \mathbb{W} , function $h : \mathbb{W} \to \overline{\mathbb{R}}$ and original minimization problem

 $\inf_{w\in\mathbb{W}}h(w)$

▶ Embedding/perturbation scheme given by a nonempty set \mathbb{U} , an anchor $\overline{u} \in \mathbb{U}$ and a Rockafellian $\mathcal{R} : \mathbb{W} \times \mathbb{U} \to \overline{\mathbb{R}}$ such that

 $h(w) = \mathcal{R}(w, \overline{u}), \ \forall w \in \mathbb{W}$

Paired spaces U and V, and Lagrangian L : W × V → R given by

$$\mathcal{L}(w, v) = \inf_{u \in \mathbb{U}} \left\{ \mathcal{R}(w, u) - \langle u - \overline{u} \mid v \rangle \right\}$$

Original minimization problem

$$\inf_{w\in\mathbb{W}}\sup_{v\in\mathbb{V}}\mathcal{L}(w,v)=\inf_{w\in\mathbb{W}}h(w)$$

Duality gap

Dual maximization problem

 $\sup_{v\in\mathbb{V}}\inf_{w\in\mathbb{W}}\mathcal{L}(w,v)$

Weak duality always holds true

$$\sup_{v\in\mathbb{V}}\inf_{w\in\mathbb{W}}\mathcal{L}(w,v)\leq\inf_{w\in\mathbb{W}}h(w)$$

When it exists, the duality gap is the nonnegative difference

Strong duality holds true, or there is no duality gap, when

$$\sup_{v\in\mathbb{V}}\inf_{w\in\mathbb{W}}\mathcal{L}(w,v)=\inf_{w\in\mathbb{W}}h(w)$$

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Classic Lagrangian duality

• Let $\theta = (\theta_1, \dots, \theta_p) : \mathbb{W} \to \mathbb{R}^p$ be a mapping, and $\overline{u} \in \mathbb{R}^p$

We consider the optimization problem

$$\min_{ heta(w)\leq \overline{u}}h(w)=\min_{eta_1(w)\leq \overline{u}_1}h(w)\ heta_{
ho}(w)\leq \overline{u}_{
ho}$$

▶ In that case, take the perturbation scheme with $U = \mathbb{R}^{p}$ and

$$\mathcal{R}(w,u) = h(w) \dotplus \iota_{\{\theta(w) \le u\}} = h(w) \dotplus \sum_{j=1}^{p} \iota_{\{\theta_j(w) \le u_j\}}$$

▶ which gives the Lagrangian $\mathcal{L} : \mathbb{W} \times \mathbb{V} \to \overline{\mathbb{R}}$, with $\mathbb{V} = \mathbb{R}^{p}$ and

$$\mathcal{L}(w,v) = h(w) + \langle \theta(w) - \overline{u} \mid v \rangle = h(w) + \sum_{j=1}^{p} v_j (\theta_j(w) - \overline{u})$$

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Illustration of the scheme in PILP

- Constraint matrix $A \in \mathbb{R}^{m \times n}$
- ▶ Anchor $\bar{b} \in \mathbb{R}^m$

Initial minimization problem

$$\begin{array}{ll}
\inf & \langle k \mid x \rangle \\
x \in \mathbb{Z}^n \\
Ax = \overline{b} \\
x \ge 0
\end{array}$$

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Step 1. Perturbation of the initial minimization problem

$$\forall b \in \mathbb{R}^{m}, \ \varphi(b) = \inf_{\substack{x \\ k = b}} \langle k \mid x \rangle$$
$$Ax = b$$
$$x \ge 0$$
$$x \in \mathbb{Z}^{n}$$

- Perturbation space: R^m
- ▶ Perturbation function $\varphi : \mathbb{R}^m \to \overline{\mathbb{R}} = \mathbb{R} \cup \{-\infty\} \cup \{+\infty\}$

unbounded unfeasible

▶ Value of the initial problem: $\varphi(\bar{b})$

Example of epigraph of the perturbation function for a PILP

Step 2. Dual space, coupling and conjugate function

Set of subadditive functions

 $\mathcal{S}^m = \{F : \mathbb{R}^m \to \overline{\mathbb{R}} | F(b_1 + b_2) \le F(b_1) \dotplus F(b_2), \forall b_1, b_2\}$

Subadditive evaluation coupling

$$c_{\mathcal{S}}: \mathbb{R}^m \times \mathcal{S}^m \to \overline{\mathbb{R}}$$

 $c_{\mathcal{S}}(b, F) = F(b)$

► Conjugate function $\varphi^{c_{\mathcal{S}}} : \mathcal{S}^{m} \to \overline{\mathbb{R}}$ $\forall F \in \mathcal{S}^{m}, \ \varphi^{c_{\mathcal{S}}}(F) = \sup_{b \in \mathbb{R}^{m}} \left\{ c_{\mathcal{S}}(b,F) + (-\varphi(b)) \right\}$

Step 3. Biconjugate and weak duality

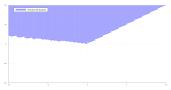
▶ Biconjugate function $\varphi^{c_S c_S'} : \mathbb{R}^m \to \overline{\mathbb{R}}$

$$\forall b \in \mathbb{R}^m, \ \varphi^{c_{\mathcal{S}}c_{\mathcal{S}}'}(b) = \sup_{F \in \mathcal{S}^m} \left\{ c_{\mathcal{S}}(b,F) - \varphi^{c_{\mathcal{S}}}(F) \right\}$$

We obtain weak duality

$$\underbrace{\sup_{F \in S^m} \left\{ F(b) + \inf_{x \in \mathbb{Z}^n_+} \{ \langle k \mid x \rangle - F(Ax) \} \right\}}_{\text{dual problem}} = \varphi^{c_S c_S'}(b) \le \varphi(b) = \begin{array}{c} \inf_x & \langle k \mid x \rangle \\ Ax = b \\ x \ge 0 \\ x \in \mathbb{Z}^n \end{array}$$

Step 4. Strong duality for the subadditive dual problem



▶ If
$$Ax_1 = b_1$$
 and $Ax_2 = b_2$, then $A(x_1 + x_2) = b_1 + b_2$
▶ So

$$\varphi(b_1+b_2) \leq \varphi(b_1) \dotplus \varphi(b_2) , \ \forall b_1, b_2 \in \mathbb{R}^m$$

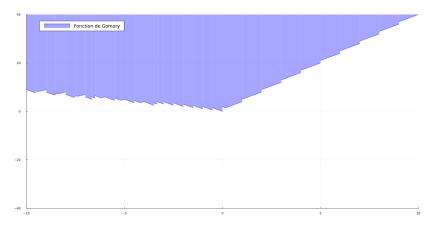
 \blacktriangleright It results that the function φ is subadditive

$$\varphi(b) + \inf_{x \in \mathbb{Z}_+^n} \underbrace{\{\langle k \mid x \rangle - \varphi(Ax)\}}_{\geq 0} \leq \varphi^{c_S c_S'}(b)$$

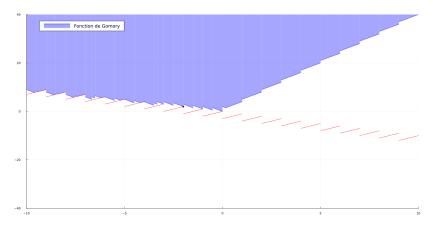
▶ Thus we have strong duality $\forall b \in \mathbb{R}^m$

$$\underbrace{\sup_{F \in S^m} \left\{ F(b) + \inf_{x \in \mathbb{Z}^n_+} \left\{ \langle k \mid x \rangle - F(Ax) \right\} \right\}}_{\text{dual problem}} = \varphi^{c_S c_S'}(b) = \varphi(b) = \begin{cases} \inf_x & \langle k \mid x \rangle \\ Ax = b \\ x \ge 0 \end{cases}$$

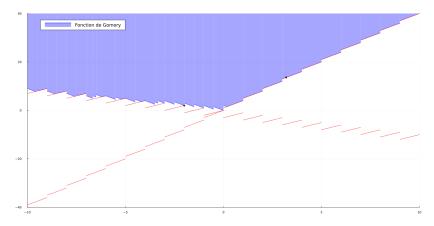
$$\varphi(b) = \max\{3b + \lceil b \rceil, 2b + -3\lceil b \rceil, -3b + \lceil 2b \rceil + \lceil \frac{3}{10}b \rceil\}$$



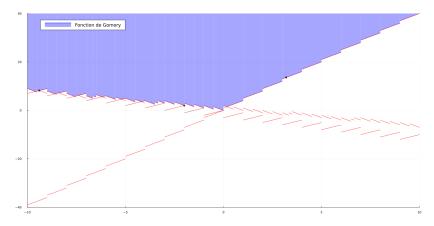
$$\varphi(b) = \max\{3b + \lceil b \rceil, 2b + -3\lceil b \rceil, -3b + \lceil 2b \rceil + \lceil \frac{3}{10}b \rceil\}$$



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$$\varphi(b) = \max\{3b + \lceil b \rceil, 2b + -3\lceil b \rceil, -3b + \lceil 2b \rceil + \lceil \frac{3}{10}b \rceil\}$$



Outline of the presentation Duality for linear programs

Duality for pure integer linear programs

Duality for pure integer linear (PILP) programs with the PDS Summary

- Couplings for optimal transport

Summary of the perturbation-duality scheme for PILP

1. We perturb a minimization problem

$$\forall \mathbf{b} \in \mathbb{R}^m, \ \varphi(\mathbf{b}) = \inf_{\substack{x \\ k = \mathbf{b} \\ x \in \mathbb{Z}^n_+}} \langle k \mid x \rangle$$

2. We pair the primal space \mathbb{R}^m and a dual space \mathcal{S}^m

$$c_{\mathcal{S}}(\cdot, \cdot) : \mathbb{R}^m imes \frac{\mathcal{S}^m}{\mathcal{S}} \to \overline{\mathbb{R}}$$

 $c_{\mathcal{S}}(b, F) = F(b)$

Reminder: set of subadditive functions

$$\mathcal{S}^m = \{F : \mathbb{R}^m \to \overline{\mathbb{R}} | F(b_1 + b_2) \leq F(b_1) \dotplus F(b_2), \ \forall b_1, b_2\}$$

3. We biconjugate the perturbation function φ

$$\underbrace{\varphi^{c_{\mathcal{S}}c_{\mathcal{S}}'}(b) \leq \varphi(b) \ , \ \forall b \in \mathbb{R}^{m}}_{\mathcal{S}}$$

weak duality is guaranteed

4. Strong duality as φ is subadditive

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Outline of the presentation

Duality for linear programs

Perturbation-duality scheme in convex analysis

Duality for pure integer linear programs

Generalized perturbation duality scheme

Developments and examples

Outline of the presentation Duality for linear programs

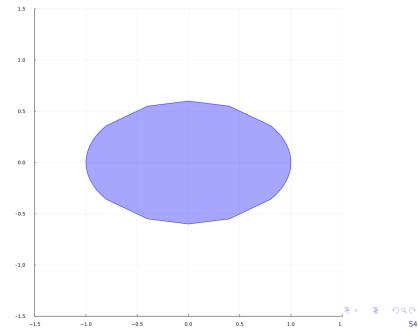
Duality for pure integer linear (PILP) programs with the PDS

Generalized perturbation duality scheme

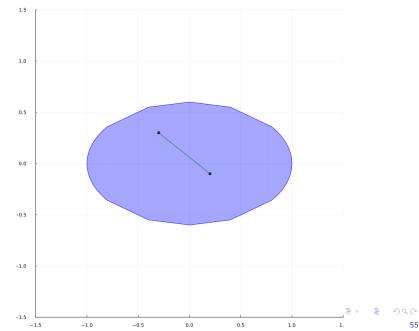
Generalized convexity

Couplings for optimal transport

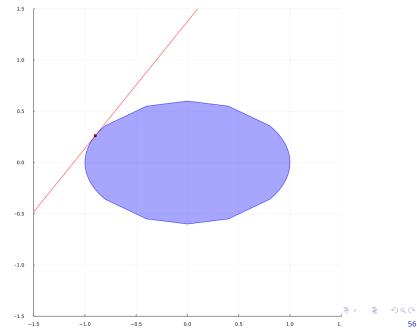
A closed convex set



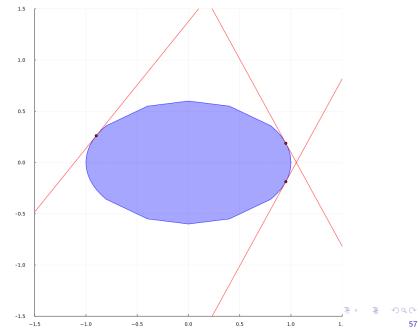
Usual definition of convexity by the interior



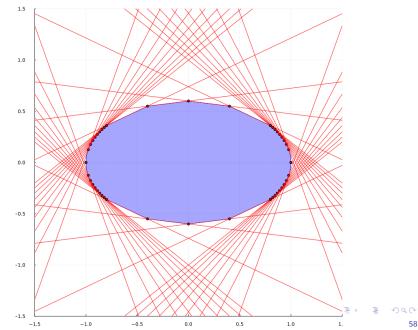
Equivalent definition for closed-convexity by the exterior



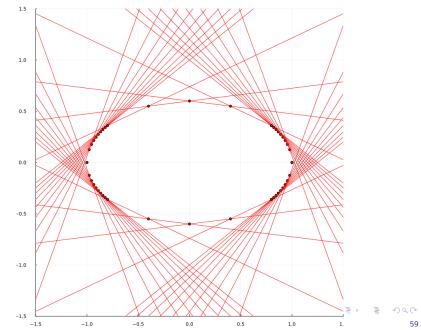
Equivalent definition for closed-convexity by the exterior



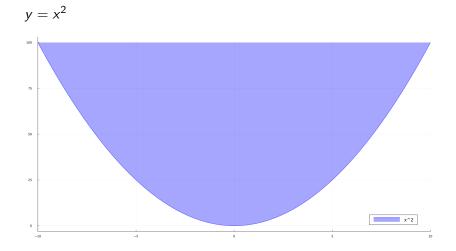
Equivalent definition for closed-convexity by the exterior



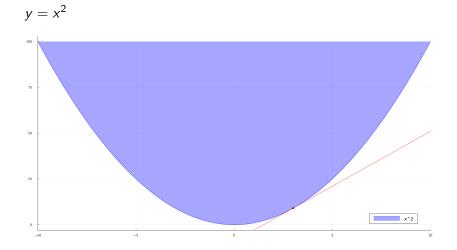
Approximation by finite number of cuts



Epigraph of a closed-convex function

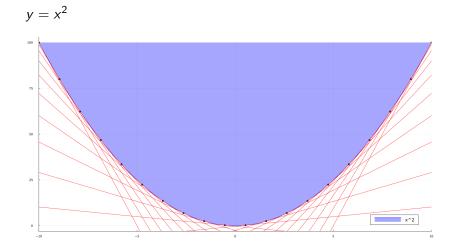


Epigraph of a closed-convex function



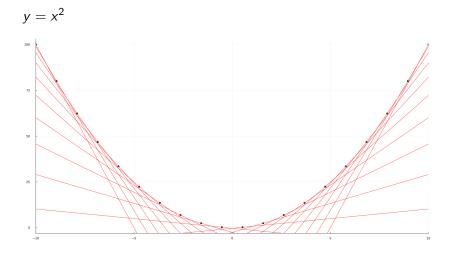
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The epigraph is above its tangents



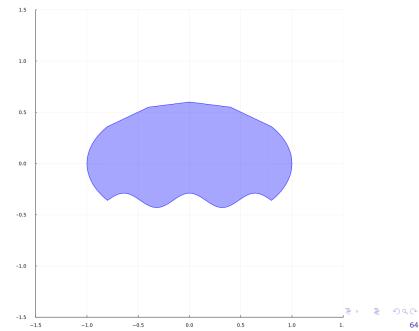
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Approximation by a finite number of cuts

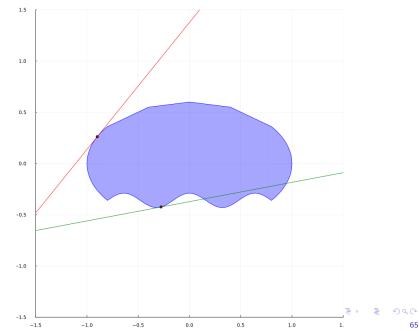


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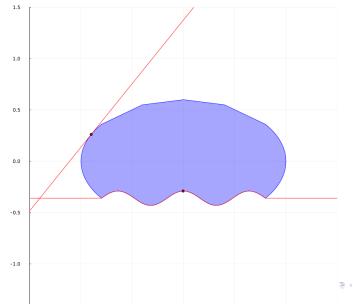
Example of a nonconvex set



Some tangents won't stay outside!



Generalized convexity: we change the shape of the tangents!



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Generalized convexity: we change the shape of the tangents!

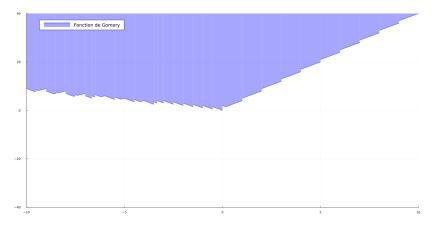
$$T(x) = \langle x \mid \alpha \rangle + \beta , \ \forall x \in \mathbb{R}^n$$

 $\begin{array}{l} \text{Scalar product } \langle \cdot \mid \cdot \rangle : \mathbb{R}^n \times \mathbb{R}^n \\ \text{Slope: } \alpha \in \mathbb{R}^n \\ \text{Intercept: } \beta \in \mathbb{R} \end{array}$

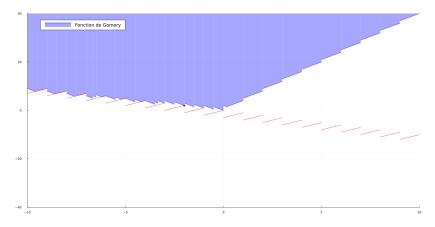
 $S(u) = c(u, v) + \beta$ Coupling $c : U \times V \rightarrow \overline{\mathbb{R}} = \mathbb{R} \cup \{-\infty, +\infty\}$ "Slope": $v \in V$ Intercept: $\beta \in \mathbb{R}$

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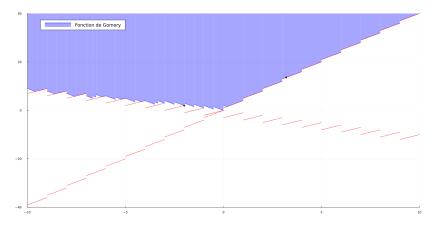
$$G(b) = \max\{3b + \lceil b \rceil, 2b + -3\lceil b \rceil, -3b + \lceil 2b \rceil + \lceil \frac{3}{10}b \rceil\}$$



$$G(b) = \max\{3b + \lceil b \rceil, 2b + -3\lceil b \rceil, -3b + \lceil 2b \rceil + \lceil \frac{3}{10}b \rceil\}$$

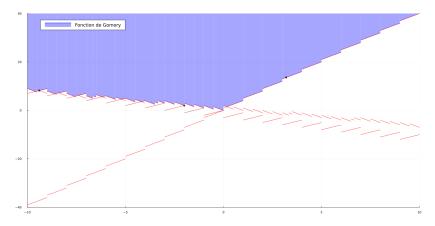


$$G(b) = \max\{3b + \lceil b \rceil, 2b + -3\lceil b \rceil, -3b + \lceil 2b \rceil + \lceil \frac{3}{10}b \rceil\}$$



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$$G(b) = \max\{3b + \lceil b \rceil, 2b + -3\lceil b \rceil, -3b + \lceil 2b \rceil + \lceil \frac{3}{10}b \rceil\}$$



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Outline of the presentation Duality for linear programs

Duality for pure integer linear (PILP) programs with the PDS

Generalized perturbation duality scheme

Generalized perturbation duality scheme

Couplings for optimal transport

Moreau lower and upper additions

The Moreau lower addition extends the usual addition with
 (+∞) + (-∞) = (-∞) + (+∞) = -∞
 The Moreau upper addition extends the usual addition with
 (+∞) + (-∞) = (-∞) + (+∞) = +∞

Background on couplings and Fenchel-Moreau conjugacies

- Let be given two sets \mathcal{X} ("primal") and \mathcal{Y} ("dual")
- Consider a coupling function $c : \mathcal{X} \times \mathcal{Y} \to \overline{\mathbb{R}} = [-\infty, +\infty]$
- We also use the notation $\mathcal{X} \stackrel{c}{\leftrightarrow} \mathcal{Y}$ for a coupling

Definition

The *c*-Fenchel-Moreau conjugate of a function $f : \mathcal{X} \to \overline{\mathbb{R}}$, with respect to the coupling *c*, is the function $f^c : \mathcal{Y} \to \overline{\mathbb{R}}$ defined by

$$f^{c}(y) = \sup_{x \in \mathcal{X}} \left(c(x, y) + (-f(x)) \right), \ \forall y \in \mathcal{Y}$$

Fenchel-Moreau conjugate (max, +)	Kernel transform $(+, \times)$	
$\sup_{x\in\mathcal{X}}\left(c(x,y)+\left(-f(x)\right)\right)$	$\int_{\mathcal{X}} c(x, y) f(x) dx$	

Background on couplings and Fenchel-Moreau conjugacies

With the coupling c, we associate the reverse coupling c'

$$c': \mathcal{Y} \times \mathcal{X} \to \overline{\mathbb{R}} \;, \;\; c'(y,x) = c(x,y) \;, \;\; \forall (y,x) \in \mathcal{Y} imes \mathcal{X}$$

The c'-Fenchel-Moreau conjugate of a function g : 𝔅 → ℝ, with respect to the coupling c', is the function g^{c'} : 𝔅 → ℝ

$$g^{c'}(x) = \sup_{y \in \mathcal{Y}} \left(c(x, y) + (-g(y)) \right), \ \forall x \in \mathcal{X}$$

The *c*-Fenchel-Moreau biconjugate *f^{cc'}* : X → ℝ of a function *f* : X → ℝ is given by

$$f^{cc'}(x) = (f^c)^{c'}(x) = \sup_{y \in \mathcal{Y}} \left(c(x, y) + (-f^c(y)) \right), \ \forall x \in \mathcal{X}$$

Dual problems: perturbation scheme [Rockafellar, 1974]

Set
$$\mathbb{W}$$
, function $h: \mathbb{W} \to \overline{\mathbb{R}}$

and original minimization problem

 $\inf_{w\in\mathbb{W}}h(w)$

Embedding/perturbation scheme given by a nonempty set X (perturbations), an element x̄ ∈ X (anchor) and a function (Rockafellian) R : W × X → R̄ such that

$$h(w) = \mathcal{R}(w, \overline{x})$$

Perturbation function

$$\phi(x) = \inf_{w \in \mathbb{W}} \mathcal{R}(w, x)$$

Original minimization problem

$$\phi(\overline{x}) = \inf_{w \in \mathbb{W}} \mathcal{R}(w, \overline{x}) = \inf_{w \in \mathbb{W}} h(w)$$

Dual problems: conjugacy, weak and strong duality

▶ Coupling $\mathcal{X} \stackrel{c}{\leftrightarrow} \mathcal{Y}$, and Lagrangian $\mathcal{L} : \mathbb{W} \times \mathcal{Y} \to \overline{\mathbb{R}}$ given by

$$\mathcal{L}(w,y) = \inf_{x \in \mathcal{X}} \left\{ \mathcal{R}(w,x) \dotplus (-c(x,y)) \right\}$$

Dual function

$$\psi(y) = -\phi^{c}(y) = \inf_{w \in \mathbb{W}} \mathcal{L}(w, y)$$

Dual maximization problem (weak duality)

$$\phi^{cc'}(\overline{x}) = \sup_{y \in \mathcal{Y}} \left\{ c(\overline{x}, y) + \psi(y) \right\} \le \inf_{w \in \mathbb{W}} h(w) = \phi(\overline{x})$$

Strong duality holds true when ϕ is *c*-convex at \overline{x} , that is,

$$\phi^{cc'}(\overline{x}) = \sup_{y \in \mathcal{Y}} \left\{ c(\overline{x}, y) + \psi(y) \right\} = \inf_{w \in \mathbb{W}} h(w) = \phi(\overline{x})$$

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Dual problems with general couplings

Dual problems: perturbation scheme [Rockafellar, 1974]

sets	optimization	primal	coupling	dual	
	set W	set $\mathcal X$	$\mathcal{X} \stackrel{c}{\leftrightarrow} \mathcal{Y}$	set ${\mathcal Y}$	
variables	decision	perturbation	c(x,y)	sensitivity	
	$w \in \mathbb{W}$	$x \in \mathcal{X}$	$\in \overline{\mathbb{R}}$	$y\in \mathcal{Y}$	
bivariate	Rockafellian			Lagrangian	
functions	$\mathcal{R}:\mathbb{W} imes\mathcal{X} o\overline{\mathbb{R}}$			$\mathcal{L}: \mathbb{W} \times \mathcal{Y} \to \overline{\mathbb{R}}$	
definition				$\mathcal{L}(w, y) =$	
				$\inf_{x\in\mathcal{X}}\left\{\mathcal{R}(w,x) \dotplus (-c(x,y))\right\}$	
property				$-\mathcal{L}(w,\cdot) = (\mathcal{R}(w,\cdot))^{c}$	
property				$-\mathcal{L}(w,\cdot)$ is c' -convex	
univariate		perturbation function		dual function	
functions		$\phi: \mathcal{X} \to \overline{\mathbb{R}}$		$\psi: \mathcal{Y} \to \overline{\mathbb{R}}$	
definition		$\phi(x) = \inf_{w \in \mathbb{W}} \mathcal{R}(w, x)$		$\psi(y) = \inf_{w \in \mathbb{W}} \mathcal{L}(w, y)$	
property				$-\psi = \phi^c$	

Anchor $\overline{x} \in \mathcal{X}$ and dual maximization problem (weak duality) $\phi^{cc'}(\overline{x}) = \sup_{y \in \mathcal{Y}} \{c(\overline{x}, y) + \psi(y)\} \leq \inf_{w \in \mathbb{W}} h(w) = \phi(\overline{x})$ Strong duality iff ϕ is *c*-convex at \overline{x} iff $\phi^{cc'}(\overline{x}) = \phi(\overline{x})$

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Duality between Lagrangians and Rockafellians (work in progress)

 $(-\mathcal{L}, \mathcal{R})$ is minimal in the inequality $(-\mathcal{L}(w, y)) \stackrel{.}{+} \mathcal{R}(w, x) \ge c(x, y)$ $-\mathcal{L}(w,\cdot) = (\mathcal{R}(w,\cdot))^c$ and $\mathcal{R}(w,\cdot) = (-\mathcal{L}(w,\cdot))^{c'}$ $-\mathcal{L}(w,\cdot) = (\mathcal{R}(w,\cdot))^c$ and $(\mathcal{R}(w,\cdot))^{cc'} = \mathcal{R}(w,\cdot)$ $\mathcal{R}(w,\cdot) = (-\mathcal{L}(w,\cdot))^{c'}$ and $(-\mathcal{L}(w,\cdot))^{c'c} = -\mathcal{L}(w,\cdot)$

The c-subdifferential is defined as the Rockafellar-Moreau subdifferential

Definition

For any function $f : \mathcal{X} \to \overline{\mathbb{R}}$ and $x \in \mathcal{X}$, the c-subdifferential is

$$\partial_c f(x) = \{y \in \mathcal{Y} \mid c(x', y) + (-f(x')) \\ \leq c(x, y) + (-f(x)), \forall x' \in \mathcal{X}\}$$

The following properties are satisfied

$$y \in \partial_c f(x) \iff f^c(y) = c(x, y) + (-f(x))$$

"à la" Fenchel-Young
$$\partial_c f(x) \neq \emptyset \Rightarrow f^{cc'}(x) = f(x)$$

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Introducing generalized convexity

Balder [1977]

Fenchel conjugatec-conjugate $f^*(v) = \sup_{u \in \mathbb{R}^m} \langle u \mid v \rangle - f(u)$ $g^c(v) = \sup_{u \in U} c(u, v) + (-g(u))$ Fenchel biconjugatec-biconjugate $f^{\star\star'}(u) = \sup_{v \in \mathbb{R}^m} \langle u \mid v \rangle - f^\star(v)$ $g^{cc'}(u) = \sup_{v \in V} c(u, v) + (-g^c(v))$ \star - convex functionsc-convex functions $\iff f = f^{\star\star'}$ $\iff g = g^{cc'}$

Perturbation-duality scheme with generalized convexity

1. We perturb a minimization problem

 $\varphi:\mathbb{R}^m\to\overline{\mathbb{R}}$

2. We pair the primal space \mathbb{R}^m and a dual space V

 $c: \mathbb{R}^m \times V \to \overline{\mathbb{R}}$

3. We biconjugate the perturbation function φ

$$\underbrace{\varphi^{\mathsf{cc'}}(b) \leq \varphi(b) \ , \ \forall b \in \mathbb{R}^m}_{\mathsf{cc'}}$$

Weak duality is guaranteed!

4. Strong duality when φ is *c*-convex

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Functions in the perturbation duality scheme

bivariate functions	univariate functions definition		property
Rockafellian			
$\mathcal{R}: \mathbb{W} \times \mathcal{X} \to \overline{\mathbb{R}}$			
Lagrangian		$\mathcal{L}(w, y) =$	
$\mathcal{L}: \mathbb{W} \times \mathcal{Y} \to \overline{\mathbb{R}}$		$\inf_{x \in \mathcal{X}} \left\{ \mathcal{R}(w, x) \dotplus \left(-c(x, y) \right) \right\}$	$-\mathcal{L}(w, \cdot) = (\mathcal{R}(w, \cdot))^{c}$
	perturbation		
	$\varphi: \mathcal{X} \to \overline{\mathbb{R}}$	$\varphi(x) = \inf_{w \in \mathbb{W}} \mathcal{R}(w, x)$	
	dual		
	$\Psi: \mathcal{Y} \to \overline{\mathbb{R}}$	$\Psi(y) = \inf_{w \in \mathbb{W}} \mathcal{L}(w, y)$	$-\Psi = \varphi^c$
	dual objective		
		(-1) (-1) (-1)	$\Phi_{\bar{x}}(y) = c(\bar{x}, y) + \Psi(y)$
	$\Phi_{\bar{X}}:\mathcal{Y}\to\overline{\mathbb{R}}$	$\Phi_{\bar{x}}(y) = c(\bar{x}, y) + (-\varphi^c(y))$	$\varphi^{cc'}(\bar{x}) = \sup_{y \in \mathcal{Y}} \Phi_{\bar{x}}(y)$

Outline of the presentation

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Developments and examples

Outline of the presentation Duality for linear programs

Duality for pure integer linear (PILP) programs with the PDS

Developments and examples

Lagrangian relaxation in LP

Couplings for optimal transport (日) (部) (注) (注) (三)

Summary of the perturbation-duality scheme

1. We partially perturb

$$\forall b \in \mathbb{R}^{m_1}, \ \varphi(b_1) = \inf_{\substack{x \\ A_1x = b_1 \\ A_2x = b_2 \\ x \ge 0 \\ x \in \mathbb{R}^n}} \langle k \mid x \rangle$$

2. We pair the primal space \mathbb{R}^{m_1} and a dual space \mathbb{R}^{m_1}

 $\langle \cdot | \cdot \rangle : \mathbb{R}^{m_1} \times \mathbb{R}^{m_1} \to \mathbb{R}$

3. We biconjugate the perturbation function φ

$$\varphi^{\star\star'}(b_1) = \sup_{\lambda \in \mathbb{R}^{m_1}} \quad \inf_{x} \quad \langle k \mid x \rangle + \langle \lambda \mid b_1 - A_1 x \rangle \quad \leq \varphi(b_1) \\ A_2 x = b_2 \\ x \geq 0 \\ x \in \mathbb{R}^n$$

4. Strong duality when φ is lsc convex

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Outline of the presentation Duality for linear programs

Duality for pure integer linear (PILP) programs with the PDS

Developments and examples

Chvátal dual problem for pure integer linear programming

Couplings for optimal transport

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Definition of Chvátal functions

Definition

The class of Chvátal functions C^m is the smallest class of functions $D \subset \{f | f : \mathbb{Q}^m \to \mathbb{Q}\}$ such that

> $b \in \mathbb{Q}^{m} \mapsto \lambda b \in D, \quad \forall b \in \mathbb{Q}^{m} \quad \text{(linear functions)}$ $\alpha F_{1} + \beta F_{2} \in D, \quad \forall F_{1}, F_{2} \in D, \quad \alpha, \beta \in \mathbb{Q}_{+} \quad \text{(conic combination)}$ $\lceil F \rceil \in D, \quad \forall F \in D \quad \text{(round-up)}$

Examples in 1D

$$b \mapsto \frac{3}{4}b$$

$$b \mapsto \lceil b \rceil$$

$$b \mapsto \frac{3}{4}b + \frac{7}{10}\lceil b \rceil$$

$$b \mapsto 15b + \frac{39}{22}\lceil \frac{3}{4}b + \frac{7}{10}\lceil b \rceil \rceil + \lceil 16b \rceil$$

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Chvátal perturbation-duality scheme

We define a perturbation function

$$\forall b \in \mathbb{Q}^m, \ \varphi(b) = \inf_{\substack{x \\ x \in \mathbb{Z}^n_+}} \langle k \mid x \rangle$$
$$Ax = b$$
$$x \in \mathbb{Z}^n_+$$

We define a coupling between primal and dual space

$$c_{\mathcal{C}}: \mathbb{Q}^m \times \mathcal{C}^m \to \mathbb{R}$$

$$c_{\mathcal{C}}(b, F) = F(b) , \ \forall b \in \mathbb{Q}^m , \ \forall F \in \mathcal{C}^m$$

We biconjugate the perturbation functions

$$\underbrace{\varphi^{c_{\mathcal{C}}c_{\mathcal{C}}'}(b) \leq \varphi(b)}, \ \forall b \in \mathbb{Q}^{m}$$

weak duality

► We get strong duality $\varphi^{c_{\mathcal{C}}c_{\mathcal{C}}'}(\bar{b}) = \varphi(\bar{b})$

Obtained dual problems

Formulation 1: $\varphi^{c_{\mathcal{C}}c_{\mathcal{C}}'}(\bar{b}) = \sup_{F \in \mathcal{C}^m} \left\{ F(\bar{b}) + \inf_{b \in \mathbb{Q}^m} \{ \varphi(b) - F(b) \} \right\}$ Formulation 2: $\varphi^{c_{\mathcal{C}}c_{\mathcal{C}}'}(\bar{b}) = \sup_{F \in \mathcal{C}^m} \left\{ F(\bar{b}) + \inf_{x \in \mathbb{Z}^n_+} \{ \langle k \mid x \rangle - F(Ax) \} \right\}$

Reminder Jeroslow's dual problem

$$sup F(\bar{b})$$

$$F(A_j) \le k_j$$

$$F(0) \le 0$$

$$F \in C^m$$

Generalized subdifferential and complementary slackness

Proposition

• φ : bounded perturbation function of a MILP

•
$$A = \left(A_j
ight)_{j=1,...,n} \in \mathbb{Q}^{m imes n}$$
 constraint matrix

▶ $\bar{b} \in \mathbb{Q}^n$ anchor

If $\hat{x} \in \{x \in \mathbb{Z}^n_+ | Ax = \overline{b}\}$ and $\widehat{F} \in \mathcal{C}^m$ are "primal"-dual optimal solutions then we have the equivalence

$$egin{aligned} \widehat{F} \in \partial^{c_{\mathcal{C}}} arphi(ar{b}) \ & \iff -k \in \partialig(-\widehat{F} \circ A \dotplus \delta_{\mathbb{Z}^n_+}ig)(\hat{x}) \end{aligned}$$

Furthermore, if $\widehat{F}(A_j) \leq k_j, \forall j = 1, ..., n$, then the following assertion is also equivalent

 $\widehat{F}(0) \leq 0$, $\widehat{F}(\overline{b}) = \varphi(\overline{b})$ and $(k_j - \widehat{F}(A_j))\hat{x}_j = 0$, $\forall j = 1, \dots, n$.

Outline of the presentation Duality for linear programs

Duality for pure integer linear (PILP) programs with the PDS

Developments and examples

Sensitivity analysis Couplings for optimal transport (日) (部) (注) (注) (三)

Subdifferential of the perturbation function (sensitivity analysis)

The perturbation function is

$$arphi(u) = \inf_{w \in \mathbb{W}} \mathcal{R}(w, u) \ , \ \forall u \in \mathbb{U}$$

Theorem [Rockafellar, 1974, Theorem 16, p. 40] For $\overline{v} \in \mathbb{V}$, the following conditions are equivalent

- 1. $\overline{\nu} \in \arg \max_{v \in \mathbb{V}} \inf_{w \in \mathbb{W}} \mathcal{L}(w, v)$ and $\max_{v \in \mathbb{V}} \inf_{w \in \mathbb{W}} \mathcal{L}(w, v) = \inf_{w \in \mathbb{W}} h(w)$
- 2. $\overline{v} \in \partial \varphi(\overline{u})$

3.
$$\inf_{w \in \mathbb{W}} h(w) = \inf_{w \in \mathbb{W}} \mathcal{L}(w, \overline{v})$$

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Subdifferential of the perturbation function (sensitivity analysis)

The convex case

Theorem [Rockafellar, 1974, Theorem 18, p. 41] Suppose that

- ▶ the function $\mathcal{R} : \mathbb{W} \times \mathbb{U} \to \overline{\mathbb{R}}$ is convex
- ► there exists w ∈ W such that the function u → R(w, u) is bounded above in a neighborhood of u

Then there exists $\overline{v} \in \mathbb{V}$ such that

- 1. $\overline{v} \in \arg \max_{v \in \mathbb{V}} \inf_{w \in \mathbb{W}} \mathcal{L}(w, v)$ and $\max_{v \in \mathbb{V}} \inf_{w \in \mathbb{W}} \mathcal{L}(w, v) = \inf_{w \in \mathbb{W}} h(w)$
- 2. $\overline{v} \in \partial \varphi(\overline{u})$
- 3. $\inf_{w \in \mathbb{W}} h(w) = \inf_{w \in \mathbb{W}} \mathcal{L}(w, \overline{v})$

Slater qualification constraint

The convex case

Theorem [Rockafellar, 1974, p. 45]

Suppose that

• the functions h and $\theta_1, \ldots, \theta_p$ are is convex

• there exists $w \in \mathbb{W}$ such that

$$\theta_1(w) < \overline{u}_1, \ldots, \theta_p(w) < \overline{u}_p$$

Then there exists $\overline{v} \in \mathbb{V}$ such that

- 1. $\overline{\nu} \in \arg \max_{v \in \mathbb{V}} \inf_{w \in \mathbb{W}} \mathcal{L}(w, v)$ and $\max_{v \in \mathbb{V}} \inf_{w \in \mathbb{W}} \mathcal{L}(w, v) = \inf_{w \in \mathbb{W}} h(w)$
- 2. $\overline{v} \in \partial \varphi(\overline{u})$
- 3. $\inf_{w \in \mathbb{W}} h(w) = \inf_{w \in \mathbb{W}} \mathcal{L}(w, \overline{v})$

Outline of the presentation Duality for linear programs

Duality for pure integer linear (PILP) programs with the PDS

Developments and examples

Couplings for optimal transport

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Optimal transport

The optimal transport problem is

$$\inf_{\pi\in\Pi(P_0,P_1)}\int_{\mathbb{A}\times\mathbb{B}}k(a,b)d\pi(a,b)$$

where

- \blacktriangleright the sets $\mathbb A$ and $\mathbb B$ are two Polish spaces
- ▶ we denote by P(A × B), P(A) and P(B) the corresponding probability spaces (rectangle and marginals)
- the set Π(P₀, P₁) is made of probabilities π ∈ P(A × B) on the rectangle, whose marginals are P₀ ∈ P(A) and P₁ ∈ P(B)
- ▶ the measurable cost function k : A × B →] − ∞, +∞], where k(a, b) represents the cost to move from a ∈ A towards b ∈ B

We introduce a suitable coupling between probabilities and functions

- We denote by C⁰_b(A) and C⁰_b(B) the spaces of continuous bounded functions
- We introduce the bilinear coupling

$$\mathcal{P}(\mathbb{A}) imes \mathcal{P}(\mathbb{B}) \stackrel{eta}{\longleftrightarrow} C^0_b(\mathbb{A}) imes C^0_b(\mathbb{B})$$

 $etaig((P_0, P_1); (f, g)ig) = \int_{\mathbb{B}} g(b) dP_1(b) - \int_{\mathbb{A}} f(a) dP_0(a)$

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The *k*-conjugacy appears naturally as a sub-product of computations with the β -conjugacy

The optimal cost is, as function of marginals,

$$K(P_0,P_1) = \inf_{\pi \in \Pi(P_0,P_1)} \int_{\mathbb{A} \times \mathbb{B}} k(a,b) d\pi(a,b)$$

and its β -conjugate is

$$egin{aligned} \mathcal{K}^eta(f,g) &= \sup_{a\in\mathbb{A},b\in\mathbb{B}} \left[g(b)-f(a)-k(a,b)
ight] \ &= \sup_{b\in\mathbb{B}} \left(g(b)+\underbrace{f^{-k}(b)}_{k ext{-conjugate}}
ight) \end{aligned}$$

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Conjugacy properties in optimal transport

$$\begin{split} \mathcal{K}(P_0,P_1) &= \inf_{\pi \in \Pi(P_0,P_1)} \int_{\mathbb{A} \times \mathbb{B}} k(a,b) d\pi(a,b) \\ \mathcal{D}(f,g) &= \sup_{a \in \mathbb{A}, b \in \mathbb{B}} \left[g(b) - f(a) - k(a,b) \right] = \sup_{b \in \mathbb{B}} \left(g(b) + f^{-k}(b) \right) \end{split}$$

We have the following conjugacy equalities and inequalities

$$\begin{split} \mathcal{K}^{\beta}(f,g) &= D(f,g) = D^{\beta'\beta}(f,g) \\ \mathcal{K}(P_0,P_1) &\geq \mathcal{K}^{\beta\beta'}(P_0,P_1) = D^{\beta'}(P_0,P_1) \\ &= \sup_{f,g} \left(\int_{\mathbb{B}} g(b) dP_1(b) - \int_{\mathbb{A}} f(a) dP_0(a) - D(f,g) \right) \\ &\geq \sup_{g-f \leq k} \left(\int_{\mathbb{B}} g(b) dP_1(b) - \int_{\mathbb{A}} f(a) dP_0(a) \right) \end{split}$$

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Perturbation for optimal transport

$$\inf_{\pi \in \Pi(P_0, P_1)} \int_{\mathbb{A} \times \mathbb{B}} k(a, b) d\pi(a, b)$$

$$= \inf_{\pi \in \mathcal{P}(\mathbb{A} \times \mathbb{B})} \left(\int_{\mathbb{A} \times \mathbb{B}} k(a, b) d\pi(a, b) \dotplus \delta_{\Pi(P_0, P_1)}(\pi) \right)$$

$$\geq \sup_{f \in C_b^0(\mathbb{A}), g \in C_b^0(\mathbb{B})} \left(-\left(\pi \mapsto \int_{\mathbb{A} \times \mathbb{B}} k(a, b) d\pi(a, b) \right)^{\alpha}(f, g) + \left(-\delta_{\Pi(P_0)}^{-\alpha} \right)^{\alpha}(f, g) + \left(-\delta_{\Pi(P_0)}^{-\alpha} \right)^{\alpha}(f, g) + \left(\int_{\mathbb{A}} f(a) dP_0 \right)^{\alpha}(f, g) + \left(\int_{\mathbb{A}} f(g) dP_0 \right)^{\alpha}(f$$

Thank you for your attention !

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