

# A business dinner problem

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## A “scheduling” problem

A problem submitted to the **Junior company** of the Ecole des Ponts a few years ago.

- A pool of **providers** and a pool of **customers**.
- Each provider and each customer have to meet once.
- Two providers have to meet at most once.

A series of dinners has to be planned. Tables are available, with a given capacity.

Two people **meet** if they take their dinner at a same table.

**Objective.** Minimize the number of dinners in order to achieve the constraints above.

# An example

Feasible schedule in 6 dinners for an instance with

- 2 tables, 5 suppliers, 6 customers
- at most 2 suppliers per table, at most 3 customers per table

dinner 1:	table 1:	suppliers: 1,2 customers: 1,2,3	table 2:	suppliers: 3,4 customers: 4,5
dinner 2:	table 1:	suppliers: 3 customers: 2	table 2:	suppliers: 5 customers: 5
dinner 3:	table 1:	suppliers: 2 customers: 4,5	table 2:	suppliers: 4 customers: 2,6
dinner 4:	table 1:	suppliers: 3 customers: 1,3	table 2:	suppliers: 5 customers: 2,6
dinner 5:	table 1:	suppliers: 1,5 customers: 4	table 2:	suppliers: 2,3 customers: 6
dinner 6:	table 1:	suppliers: 4,5 customers: 1,3	table 2:	suppliers: 1 customers: 5,6

# Parameters and Results

## Parameters:

- Number of tables:  $t$
- Number of suppliers:  $s$
- Number of customers:  $c$
- At most  $\sigma$  suppliers per table
- At most  $\gamma$  customers per table

## Contribution:

- Feasible schedule with not too many dinners for any value of the parameters
- Lower bound for any value of the parameters
- Special cases with closed formula
- Open questions

# Related works: combinatorial design

## Kirkman Schoolgirl problem

*15 schoolgirls must be split into five groups of three girls each day of the week so that no two girls are in the same group twice.*

Similarity with our problem: suppliers = schoolgirls, tables = groups, etc.

## Social Golfer problem

## Howell designs

## Oberwolfach problem (Ringel, 1967)

*Is it possible to seat an odd number  $n$  of mathematicians at  $m$  round tables of sizes  $(\ell_1, \dots, \ell_m)$  with  $\sum_i \ell_i = n$  in  $(n-1)/2$  dinners so that each mathematician sits next to everyone else exactly once?*

Conjecture: 'yes' (except for exceptional values of the  $\ell_i$ 's)

## Bipartite Analogue of the Oberwolfach problem

## Oberwolfach Rectangular Table Negotiation problem

None of these are optimization problems.

## Lower bounds

- Number of tables:  $t$
- Number of suppliers:  $s$
- Number of customers:  $c$
- At most  $\sigma$  suppliers per table
- At most  $\gamma$  customers per table

Two suppliers at most once together; each supplier and each customer exactly once together.

Lower bounds on the number of dinners:

$$lb_1 = \left\lceil \frac{s}{\sigma} \right\rceil \quad lb_2 = \left\lceil \frac{c}{\gamma} \right\rceil \quad lb_3 = \left\lceil \frac{s}{t\sigma} \left\lceil \frac{c}{\gamma} \right\rceil \right\rceil$$
$$lb_4 = \left\lceil \frac{\sqrt{s}}{t\gamma} \left( (c - \gamma) \max \left( \sqrt{\frac{\gamma}{c - \gamma}}, 1 \right) + \frac{\gamma}{\max \left( \sqrt{\frac{\gamma}{c - \gamma}}, 1 \right)} \right) \right\rceil$$
$$lb_5 = \max_{j \in \{2, \dots, \sigma\}} \left\lceil \frac{s}{t} \left( \frac{2}{j} \left\lceil \frac{c}{\gamma} \right\rceil - \frac{s - 1}{j(j - 1)} \right) \right\rceil$$

## Lower bounds $lb_1$ and $lb_2$

- Number of tables:  $t$
- Number of suppliers:  $s$
- Number of customers:  $c$
- At most  $\sigma$  suppliers per table
- At most  $\gamma$  customers per table

During a dinner, any customer sits at a table with at most  $\sigma$  suppliers. To meet all suppliers, he needs at least  $\lceil s/\sigma \rceil$  dinners.

During a dinner, any supplier sits at a table with at most  $\gamma$  customers. To meet all customers, he needs at least  $\lceil c/\gamma \rceil$  dinners.

## Lower bounds $lb_3$ and $lb_5$

$x_{ij}$  = number of times supplier  $i$  sits at a table with  $j$  suppliers (him included)

$$\begin{aligned} \min \quad & \frac{1}{t} \sum_{i=1}^s \sum_{j=1}^{\sigma} \frac{1}{j} x_{ij} \\ \text{s.t.} \quad & \sum_{j=1}^{\sigma} x_{ij} \geq \left\lceil \frac{c}{\gamma} \right\rceil \quad i \in [s] \\ & \sum_{j=1}^{\sigma} (j-1)x_{ij} \leq s-1 \quad i \in [s] \\ & x_{ij} \in \mathbb{R}_+ \quad i \in [s], j \in [\sigma]. \end{aligned}$$

Optimal value equal to

$$\frac{s}{t} \max \left( \frac{1}{\sigma} \left\lceil \frac{c}{\gamma} \right\rceil, \max_{j \in \{2, \dots, \sigma\}} \left( \frac{2}{j} \left\lceil \frac{c}{\gamma} \right\rceil - \frac{s-1}{j(j-1)} \right) \right).$$

**Proof.** Strong duality. Maximization of a piecewise linear concave function  $\mathbb{R} \rightarrow \mathbb{R}$ .



# Exact solutions

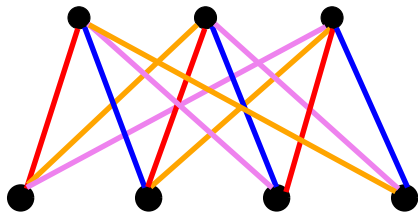
If  $\sigma = 1$ , then the optimal number of dinners is

$$\max \left( s, \left\lceil \frac{c}{\gamma} \right\rceil, \left\lceil \frac{s}{t} \left\lceil \frac{c}{\gamma} \right\rceil \right\rceil \right).$$

**Proof.** If  $\gamma = 1$ , proper edge-coloring of  $K_{s,c}$  with each color present at most  $t$  times and  $lb_1, lb_2, lb_3$ . For  $\gamma > 1$ , make groups.

$$t = s = 3, c = 4, \sigma = \gamma = 1$$

providers



customers

4 colors  $\Rightarrow$  4 dinners

# Howell designs

A square array is a **Howell design** if

1. every cell is either empty or contains an unordered pair of elements (*symbols*) chosen from a set of size  $2n$ ,
2. every symbol occurs exactly once in each row and each column,
3. every unordered pair of symbols occurs at most once in the array.

$t = 3$  tables,  $s = 4$  suppliers,  
 $\lceil c/\gamma \rceil = 3$  groups of customers (the columns),  
at most  $\sigma = 2$  suppliers per table, in 3 dinners (the rows)

1,2	3,4	
3	1	2,4
4	2	1,3

$t = 5$  tables,  $s = 6$  suppliers,  $\lceil c/\gamma \rceil = 5$   
groups of customers (the columns), at  
most  $\sigma = 2$  suppliers per table, in 5  
dinners (the rows)

	6	2,3	4,5	1
6	3,4	1	2	5
3,5	1,2		6	4
2,4		5	1	3,6
1	5	4,6	3	2

# Exact solutions

If  $\sigma = 2$ ,  $s > c/\gamma$ ,  $t \geq \min(c/\gamma, s/2)$  and  $(\lceil c/\gamma \rceil, s) \neq (2, 4)$ , the optimal number of dinners is

$$\max \left( \left\lceil \frac{c}{\gamma} \right\rceil, \left\lceil \frac{s}{2} \right\rceil \right).$$

**Proof.** Almost directly from Howell design and  $lb_1, lb_2$ .

If there are 4 tables, 7 suppliers, 26 customers, at most 2 suppliers per table, at most 4 customers per table, the optimal number of dinners is 7.

# Exact solutions

If  $t = \lceil s/2 \rceil$ ,  $\sigma = 2$ , and  $\lceil c/\gamma \rceil \geq \frac{3}{2}s$ , then the optimal number of dinners is

$$2 \left\lceil \frac{c}{\gamma} \right\rceil - s + 1.$$

**Proof.** Combination of the previous two exact solutions and  $lb_5$ .

If there are 5 tables, 9 suppliers, 41 customers, at most 2 suppliers per table, at most 3 customers per table, the optimal number of dinners is 20.

## Feasible schedules

$r(t, \mathbf{s}, \mathbf{c}, \sigma, \gamma) :=$  optimal value

$$r(t_1, \mathbf{s}, \mathbf{c}, \sigma, \gamma) \leq \left\lceil \frac{t_2}{t_1} \right\rceil r(t_2, \mathbf{s}, \mathbf{c}, \sigma, \gamma)$$

$$r(t, \mathbf{s}, \mathbf{c}, \sigma_1, \gamma) \leq \left\lceil \frac{\sigma_2}{\sigma_1} \right\rceil r(t, \mathbf{s}, \mathbf{c}, \sigma_2, \gamma)$$

$$r(t, \mathbf{s}, \mathbf{c}, \sigma, \gamma_1) \leq r(t, \mathbf{s}, \lceil \mathbf{c}/\gamma_2 \rceil, \sigma, \lceil \gamma_1/\gamma_2 \rceil) \quad \text{if } \gamma_1 \geq \gamma_2$$

General upper bounds on  $r(t, \mathbf{s}, \mathbf{c}, \sigma, \gamma)$

$$ub_1 = \left\lceil \frac{2}{\sigma} \right\rceil \left\lceil \frac{1}{t} \min \left( \left\lceil \frac{\mathbf{c}}{\gamma} \right\rceil, \mathbf{s} \right) \right\rceil \max \left( \left\lceil \frac{\mathbf{c}}{\gamma} \right\rceil, \left\lceil \frac{\mathbf{s}}{2} \right\rceil \right)$$

$$ub_2 = \left\lceil \frac{1}{t} \left\lceil \frac{\mathbf{s}}{\sigma} \right\rceil \right\rceil \left( 1 - \sigma + \sigma \max \left( \left\lceil \frac{\mathbf{c}}{\gamma} \right\rceil, 2 \left\lceil \frac{\mathbf{s}}{\sigma} \right\rceil \right) \right) \quad \text{if } \lceil \mathbf{s}/\sigma \rceil \leq \lceil \mathbf{c}/\gamma \rceil$$

Inefficient when  $\sigma \geq 3$ .

# Open question: making groups of customers

Example from the beginning:

$t = 2$  tables,  $s = 5$  suppliers,  $c = 6$  customers, at most  $\sigma = 2$  suppliers per table, at most  $\gamma = 3$  customers per table

dinner 1:	table 1:	suppliers: 1,2 customers: 1,2,3	table 2:	suppliers: 3,4 customers: 4,5
dinner 2:	table 1:	suppliers: 3 customers: 2	table 2:	suppliers: 5 customers: 5
dinner 3:	table 1:	suppliers: 2 customers: 4,5	table 2:	suppliers: 4 customers: 2,6
dinner 4:	table 1:	suppliers: 3 customers: 1,3	table 2:	suppliers: 5 customers: 2,6
dinner 5:	table 1:	suppliers: 1,5 customers: 4	table 2:	suppliers: 2,3 customers: 6
dinner 6:	table 1:	suppliers: 4,5 customers: 1,3	table 2:	suppliers: 1 customers: 5,6

## Open question: making groups of customers

$t = 2$  tables,  $s = 5$  suppliers,  $c = 6$  customers, at most  $\sigma = 2$  suppliers per table, at most  $\gamma = 3$  customers per table

Optimal solution:

dinner 1:	table 1:	suppliers: 1,2 customers: 1,2,3	table 2:	suppliers: 3,4 customers: 4,5,6
dinner 2:	table 1:	suppliers: 3 customers: 1,2,3	table 2:	suppliers: 2,5 customers: 4,5,6
dinner 3:	table 1:	suppliers: 4,5 customers: 1,2,3	table 2:	suppliers: 1 customers: 4,5,6

*Is there always an optimal solution in which the customers are split into groups, the members of each group staying together for all dinners?*

## Open question: $t = \gamma = 1$

*What is the optimal solution when there is only one table and exactly one customer per table?*



Thank you.