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**Exercise 1** (Dual formulation). Let  $g : \mathbb{R}^n \to \mathbb{R}^m$ . Show that

1.  $\mathbb{I}_{g(x)=0} = \sup_{\lambda \in \mathbb{R}^m} \lambda^\top g(x)$ 

2. 
$$\mathbb{I}_{g(x)\leq 0} = \sup_{\lambda\in\mathbb{R}^m_+}\lambda^+g(x)$$

3.  $\mathbb{I}_{g(x)\in C} = \sup_{\lambda\in -C^+} \lambda^\top g(x)$  where C is a closed convex cone, and  $C^+ := \{\lambda \in \mathbb{R}^m \mid \lambda^\top c \ge 0, \forall c \in C\}.$ 

**Exercise 2** (Linear Programming). Consider the following linear problem (LP)

$$(P) \quad \underset{x \ge 0}{\operatorname{Min}} \quad c^{\top} x$$
$$s.t. \quad Ax =$$

b

- 1. Show that the dual of (P) is an LP.
- 2. Show that the dual of the dual of (P) is equivalent to (P).

**Exercise 3** (Quadratically Constrained Quadratic Programming). *Consider the problem* 

$$(QCQP) \quad \underset{x \in \mathbb{R}^n}{\operatorname{Min}} \quad \frac{1}{2} x^\top P_0 x + q_0^\top x + r_0$$
$$\frac{1}{2} x^\top P_i x + q_i^\top x + r_i \le 0 \quad \forall i \in \mathbb{R}^n$$

where  $P_0 \in S_{++}^n$  and  $P_i \in S_{+}^n$ .

- 1. Show by duality that there exists, for  $\mu \in \mathbb{R}^m_+$ ,  $P_\mu, q_\mu$  and  $r_\mu$ , we have  $g(\mu) = -\frac{1}{2}q_\mu P_\mu^{-1} + r_\mu$  such that  $val(P) \ge g(\mu)$ .
- 2. Give an easy condition under which  $val(P) = \sup_{\mu \ge 0} g(\mu).$

**Exercise 4** (Conic Programming). Let  $K \subset \mathbb{R}^n$  be a closed convex pointed cone, and denote  $x \preceq_K y$  iff  $y \in x + K$ . Consider the following program, with  $A \in M_{m,n}$  and  $b \in \mathbb{R}^m$ .

$$(P) \quad \underset{x \in \mathbb{R}^n}{\min} \quad c^{\top} x$$
$$s.t. \quad Ax = b$$
$$x \preceq_K 0$$

- 1. Show that (P) is a convex optimization problem.
- 2. Denote  $\mathcal{L}(x,\lambda,\mu) = c^{\top}x + \lambda^{\top}(Ax b) + \mu^{\top}x$ . Show that  $\operatorname{val}(P) = \operatorname{Min}_{x \in \mathbb{R}^n} \sup_{\lambda \in \mathbb{R}^m, \mu \in K^+} \mathcal{L}(x,\lambda,\mu)$ .
- 3. Give a dual problem to (P).

**Exercise 5** (Duality gap). Consider the following problem

$$\begin{array}{ll} & \underset{x \in \mathbb{R}, y \in \mathbb{R}^+_*}{\operatorname{Min}} & e^{-x} \\ & s.t. & x^2/y \leq 0 \end{array}$$

- 1. Find the optimal solution of this problem.
- 2. Write and solve the (Lagrangian) dual problem. Is there a duality gap ?

**Exercise 6** (Two-way partitionning). Let  $W \in S_n$  be a symmetric matrix, consider the follow- $\in [m]$  problem.

$$\begin{array}{ll} (P) & \underset{x \in \mathbb{R}^n}{\operatorname{Min}} & x^\top W x \\ & s.t. & x_i^2 = 1 & \forall i \in [n] \end{array}$$

1. Consider a set of n element that you want to partition in 2 subsets, with a cost  $c_{i,j}$ if i and j are in the same set, and a cost  $-c_{i,j}$  if they are in a different set. Justify that it can be solved by solving (P).

- 2. Is (P) a convex problem ?
- 3. Show that, for any  $\lambda \in \mathbb{R}^n$  such that  $W + \text{diag}(\lambda) \succeq 0$ , we have  $\text{val}(P) \ge -\sum \lambda_i$ . Deduce a lower bound on val(P).

Exercise 7 (Linear SVM : duality). Consider the following problem (see : https:// www.youtube.com/watch?v=IOetFPgsMUc for background)

$$\min_{\substack{w \in \mathbb{R}^d, b \in \mathbb{R}}} \frac{1}{2} \|w\|^2$$
s.t.  $y_i(w^\top x_i + b) \ge 1 \qquad \forall i \in [n]$ 
 $\eta_i \ge 0 \qquad \forall i \in [n]$ 

- 1. In which case can we guarantee strong duality ?
- 2. Write the dual of this optimization problem and express the optimal primal solution  $(w^{\sharp}, b^{\sharp})$  in terms of the optimal dual solution.