Colloquium du CERMICS

Moments, Positive Polynomials, and the Christoffel Function
Jean-Bernard Lasserre
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# Moments, Positive Polynomials, and the Christoffel Function 

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CERMICS COLLOQUIUM, December 2022

# ... in collaboration with <br> D. Henrion, M. Korda, V. Magron, S. Marx, E. Pauwels, M. Putinar, T. Weisser 




## The <br> Christoffel-Darboux Kernel for Data Analysis

Jean Bernard Lasserre, Edouard Pauwels and Mihai Putinar


- Part I:
- The Moment-SOS hierarchy
- SOS-based CERTIFICATES of POSITIVITY
- Illustration of the Moment-SOS hierarchy for optimization
- Part II:
- The Christoffel function
- Anplications and link with Optimization
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## A brief overview of the methodology on 3 examples

Let $\mathbf{P}$ be the initial problem to solve, for instance

- Optimization $\min _{\mathbf{x}}\{f(\mathbf{x}): \mathbf{x} \in \mathbf{K}\}$
- Optimal control

$$
\begin{aligned}
\min _{u} & \int_{0}^{1} h(\mathbf{x}(t), u(t)) d t \\
& \dot{x}(t)=f(\mathbf{x}(t), u(t)), \quad t \in(0,1) \\
& \mathbf{x}(t) \in X ; \quad u(t) \in U, \quad t \in[0,1]
\end{aligned}
$$

- Compute (or approximate)

$$
\tau=\int_{\mathbf{K}} \mathbf{x}^{\alpha} d \lambda(\mathbf{x}), \quad \alpha \in \Gamma
$$

and in particular, $\operatorname{vol}(\mathbf{K})(\Gamma=\{0\})$ where $\mathbf{K}$ is compact basic semi-algebraic set.

## basic strategy

－挶（i）Search for a measure $\mu$ whose support is the solution

$$
d \mu(\mathbf{x})=\delta_{x^{*}} ; d \mu(x, u, t)=\delta_{x(t), u(t)}(d(x, u)) d t ; d \mu(\mathbf{x})=\lambda_{\mathbf{K}}(d \mathbf{x})
$$

－뭉（ii）compute its moments，and
－唤（iii）recover the solution from moments．
Implementation in Three steps：

## I：LIFTING

I：Build up an infinite－dimensional LP with $\mu$ as unknown：

Constraints of the initial problem become挶 LINEAR constraints
on the unknown moments $\left(\mu_{\alpha}\right)_{\alpha \in \mathbb{N}^{n}}$ of $\mu$

## II：Truncation

Consider only FINITELY MANY candidate moments $y_{\alpha}$ ：
脤 Semidefinite constraints on the scalars $y_{\alpha}$＇s state necessary conditions to qualify them as moments of some measure $\mu$

挶 Solve the resulting finite－dimensional convex（conic） optimization problem to obtain a guaranteed lower bound．

## III：Iterate

．喂 Increase the number of moments considered and iterate so as to obtain a monotone non increasing sequence of lower bounds which converges to the optimal value．

## Consider the polynomial optimization problem:

$$
\mathbf{P}: \quad f^{*}=\min \left\{f(\mathbf{x}): \quad g_{j}(\mathbf{x}) \geq 0, j=1, \ldots, m\right\},
$$

for some polynomials $f, g_{j} \in \mathbb{R}[\mathbf{x}]$.


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## Why Polynomial Optimization?

After all ... $\mathbf{P}$ is just a particular case of Non Linear Programming (NLP)!

## True!

... if one is interested with a LOCAL optimum only!!
망 Many minimization algorithms do the job efficiently.
맚 The fact that $f, g_{i}$ are does not help much!

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## Global Maxima



## BUT for GLOBAL Optimization ... the picture is different!

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## Remember that for the GLOBAL minimum :

$$
f^{*}=\sup _{\lambda}\{\lambda: f(\mathbf{x})-\lambda \geq 0 \quad \forall \mathbf{x} \in \mathbf{K}\} .
$$

(Not true for a LOCAL minimum!)


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## and so to compute (or approximate) $f^{*}$...咹 one needs to handle EFFICIENTLY the difficult constraint

$$
f(\mathbf{x})-\lambda \geq 0 \quad \forall \mathbf{x} \in \mathbf{K}, \quad\left(f-\lambda 1 \in \mathscr{P}(\mathbf{K})_{+}\right)
$$

## i.e. one needs

## for the polynomial $\mathbf{x} \mapsto f(\mathbf{x})-\lambda$ !

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 무웅 one needs to handle EFFICIENTLY the difficult constraint$$
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$$

## i.e. one needs <br> TR TRACTABLE CERTIFICATES of POSITIVITY on $\mathbf{K}$ for the polynomial $\mathbf{x} \mapsto f(\mathbf{x})-\lambda!$

## REAL ALGEBRAIC GEOMETRY helps!!!!

## Indeed, POWERFUL CERTIFICATES OF POSITIVITY EXIST!

## Moreover .... and importantly,

## Such certificates are amenable to

## REAL ALGEBRAIC GEOMETRY helps!!!!

## Indeed, POWERFUL CERTIFICATES OF POSITIVITY EXIST!

## Moreover .... and importantly,

Such certificates are amenable to PRACTICAL COMPUTATION!

## SOS-based certificate

A polynomial $p$ is a sum-of-squares (SOS) if and only if

$$
p(\mathbf{x})=\sum_{k=1}^{s} q_{k}(\mathbf{x})^{2}, \quad \forall \mathbf{x} \in \mathbb{R}^{n}
$$

for some polynomials $q_{k}$.

망ㅇ Detecting whether a given polynomial $p$ is SOS can be done efficiently by solving a SEMIDEFINITE PROGRAM

> 图 A SEMIDEFINITE PROGRAM (SDP) is a CONIC, CONVEX OPTIMIZATION PROBLEM that can be solved EFFICIENTLY (up to arbitrary fixed precision)

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팡 CONVEX OPTIMIZATION PROBLEM that can be solved EFFICIENTLY (up to arbitrary fixed precision)

## Illustration for univariate polynomials

Let $v_{d}(t)=\left(1, t, t^{2}, \ldots, t^{d}\right)$ and let $p$ be of even degree $2 d$.

$$
p(t)=\sum_{k=1}^{2 d} p_{k} t^{k} \quad\left(=\left\langle\mathbf{p}, v_{2 d}(t)\right\rangle\right)
$$

is SOS if and only if there exists $Q \succeq 0$ such that

$$
p(t)=\left[\begin{array}{c}
1 \\
t \\
\cdots \\
t^{d}
\end{array}\right]^{T} \underbrace{\left[\begin{array}{cccc}
a & b & c & \cdots \\
b & d & e & \cdots \\
c & e & f & \cdots \\
\cdots & \cdots & \cdots & \cdots
\end{array}\right]}_{Q \succeq 0}\left[\begin{array}{c}
1 \\
t \\
t^{2} \\
\cdots \\
t^{d}
\end{array}\right]
$$

$$
\begin{aligned}
Q & \succeq 0 \Rightarrow Q=\sum_{j=1}^{s} \mathbf{q}_{\mathbf{k}} \mathbf{q}_{k}{ }^{T} \\
{\left[\begin{array}{c}
1 \\
t \\
t^{2}
\end{array}\right]^{T} Q\left[\begin{array}{c}
1 \\
t \\
\cdots \\
t^{d}
\end{array}\right] } & =\sum_{k=1}^{s}\left(\left[\begin{array}{c}
1 \\
t \\
\cdots \\
t^{d}
\end{array}\right]^{T} \mathbf{q}_{\mathrm{k}}\right)\left(\mathbf{q}_{\mathrm{k}}{ }^{\top}\left[\begin{array}{c}
1 \\
t \\
\ldots \\
t^{d}
\end{array}\right]\right) \\
& =\sum_{k=1}^{s} q_{k}(t)^{2}
\end{aligned}
$$

Conversely if

$$
p(t)=\sum_{k=1}^{s} q_{k}(t)^{2}
$$

then write

$$
\begin{aligned}
p(t) & =\sum_{k=1}^{s}\left(\left[\begin{array}{c}
1 \\
t \\
\cdots \\
t^{d}
\end{array}\right]^{T} \mathbf{q}_{\mathbf{k}}\right)\left(\mathbf{q}_{\mathbf{k}}^{T}\left[\begin{array}{c}
1 \\
t \\
\cdots t^{d}
\end{array}\right]\right) \\
& =\sum_{k=1}^{s}\left[\begin{array}{c}
1 \\
t \\
\cdots \\
t^{d}
\end{array}\right]^{T}(\underbrace{\sum_{k=1}^{s} \mathbf{q}_{\mathbf{k}} \mathbf{q}_{\mathbf{k}}^{T}}_{Q \geq 0})\left[\begin{array}{c}
1 \\
t \\
\cdots t^{d}
\end{array}\right]
\end{aligned}
$$

## Example

Let $t \mapsto f(t)=6+4 t+9 t^{2}-4 t^{3}+6 t^{4}$. Is $f$ an SOS? Do we have

$$
f(t)=\left[\begin{array}{c}
1 \\
t \\
t^{2}
\end{array}\right]^{T} \underbrace{\left[\begin{array}{lll}
a & b & c \\
b & d & e \\
c & e & f
\end{array}\right]}_{Q \succeq 0}\left[\begin{array}{c}
1 \\
t \\
t^{2}
\end{array}\right]
$$

We must have:

$$
a=6 ; 2 b=4 ; d+2 c=9 ; 2 e=-4 ; f=6
$$

And so we must find a scalar $c$ such that


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t \\
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$$

We must have:

$$
a=6 ; 2 b=4 ; d+2 c=9 ; 2 e=-4 ; f=6 .
$$

And so we must find a scalar c such that

$$
Q=\left[\begin{array}{ccc}
6 & 2 & c \\
2 & 9-2 c & -2 \\
c & -2 & 6
\end{array}\right] \succeq 0 .
$$

With $c=-4$ we have

$$
Q=\left[\begin{array}{ccc}
6 & 2 & -4 \\
2 & 17 & -2 \\
-4 & -2 & 6
\end{array}\right] \succeq 0
$$

et

$$
\begin{gathered}
Q=2\left[\begin{array}{c}
\sqrt{(2 / 2)} \\
0 \\
\sqrt{(2)} / 2
\end{array}\right]\left[\begin{array}{c}
\sqrt{(2 / 2)} \\
0 \\
\sqrt{(2)} / 2
\end{array}\right]^{\prime}+9\left[\begin{array}{c}
2 / 3 \\
-1 / 3 \\
-2 / 3
\end{array}\right]\left[\begin{array}{c}
2 / 3 \\
1 / 3 \\
-2 / 3
\end{array}\right]^{\prime} \\
+18\left[\begin{array}{c}
1 / \sqrt{(18)} \\
4 / \sqrt{(18)} \\
-1 / \sqrt{(18)}
\end{array}\right]\left[\begin{array}{c}
1 / \sqrt{(18)} \\
4 / \sqrt{(18)} \\
-1 / \sqrt{(18)}
\end{array}\right]^{\prime}
\end{gathered}
$$

$$
f(t)=\left(1+t^{2}\right)^{2}+\left(2-t-2 t^{2}\right)^{2}+\left(1+4 t-t^{2}\right)^{2}
$$

## Let $\mathbf{K}:=\left\{\mathbf{x}: g_{j}(\mathbf{x}) \geq 0, \quad j=1, \ldots, m\right\}$ be compact (with $g_{1}(\mathbf{x})=M-\|\mathbf{x}\|^{2}$, so that $\mathbf{K} \subset \mathbf{B}(0, M)$ ).



$$
\begin{gathered}
\text { Let } \mathbf{K}:=\left\{\mathbf{x}: g_{j}(\mathbf{x}) \geq 0, \quad j=1, \ldots, m\right\} \\
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\end{gathered}
$$



## Theorem (Putinar's Positivstellensatz)

If $f \in \mathbb{R}[\mathbf{x}]$ is strictly positive ( $f>0$ ) on $\mathbf{K}$ then:

$$
\dagger \quad f(\mathbf{x})=\sigma_{0}(\mathbf{x})+\sum_{j=1}^{m} \sigma_{j}(\mathbf{x}) g_{j}(\mathbf{x}), \quad \forall \mathbf{x} \in \mathbb{R}^{n},
$$

for some SOS polynomials $\left(\sigma_{j}\right) \subset \mathbb{R}[\mathbf{x}]$.

## However ... In Putinar's theorem

... nothing is said on the DEGREE of the SOS polynomials $\left(\sigma_{j}\right)$ !

## BUT ... GOOD news

망 Testing whether $\dagger$ holds for some $\operatorname{SOS}\left(\sigma_{j}\right) \subset \mathbb{R}[\mathbf{x}]$

## However ... In Putinar's theorem

... nothing is said on the DEGREE of the SOS polynomials $\left(\sigma_{j}\right)$ !

## BUT ... GOOD news ..!!

嗼 Testing whether $\dagger$ holds for some $\operatorname{SOS}\left(\sigma_{j}\right) \subset \mathbb{R}[\mathbf{x}]$ with a degree bound, is SOLVING an SDP!

## Immediate application

In ANY application where one need to impose that a polynomial $f$ (to be determined) must be positive on $\mathbf{K}$, then :

## DECLARE

$$
\dagger \quad f(\mathbf{x})=\sigma_{0}(\mathbf{x})+\sum_{j=1}^{m} \sigma_{j}(\mathbf{x}) g_{j}(\mathbf{x}), \quad \forall \mathbf{x} \in \mathbb{R}^{n}
$$

with the additional constraint $\operatorname{deg}\left(\sigma_{j} g_{j}\right) \leq 2 t$ for all $j=1, \ldots, m$.
where the degree-parameter $t$ is YOUR CHOICE!
Then identifying both sides of the identity yields

- 幈 Linear constraints on the coefficients of $f$ and $\sigma$- 唌 Semidefinite constraints coming from
on the $\sigma_{j}$ 's


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Then identifying both sides of the identity yields:

- 喛 Linear constraints on the coefficients of $f$ and $\sigma_{j}$,- 娮 Semidefinite constraints coming from SOS conditions on the $\sigma_{j}$ 's
- In fact, polynomials NONNEGATIVE ON A SET K $\subset \mathbb{R}^{n}$ are ubiquitous. They also appear in many important applications (outside optimization),


## among which

Optimization, Probability, Optimal and Robust Control, non-linear PDEs, Game theory, Signal processing, multivariate integration, etc.

## Dual side: The K-moment problem

Given a real sequence $\mathbf{y}=\left(y_{\alpha}\right), \alpha \in \mathbb{N}^{n}$, does there exist a Borel measure $\mu$ on $\mathbf{K}$ such that

$$
\dagger \quad y_{\alpha}=\int_{\mathbf{K}} x_{1}^{\alpha_{1}} \cdots x_{n}^{\alpha_{n}} d \mu, \quad \forall \alpha \in \mathbb{N}^{n} \quad ?
$$

If yes then $y$ is said to have a representing measure supported on $\mathbf{K}$.

$$
\text { Let } \mathbf{K}:=\left\{\mathbf{x}: g_{j}(\mathbf{x}) \geq 0, \quad j=1, \ldots, m\right\}
$$

be compact (with $g_{1}(\mathbf{x})=M-\|\mathbf{x}\|^{2}$, so that $\mathbf{K} \subset \mathbf{B}(0, M)$ ).

Theorem (Dual side of Putinar's Theorem)
A seauence $\mathbf{v}=\left(v_{\sim}\right), \alpha \in \mathbb{N}^{n}$, has a representing measure supported on K IF AND ONLY IF for every $d=0,1$

맚아 The real symmetric matrix $\mathbf{M}_{2}(y)$ is called the MOMENT MATRIX associated with the sequence $y$

榢 The real symmetric matrix $\mathbf{M}_{d}\left(g_{j} y\right)$ is called the LOCALIZING MATRIX associated with the sequence $y$ and the polynomial $g_{j}$.

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$$
(\star) \quad \mathbf{M}_{d}(y) \succeq 0 \quad \text { and } \quad \mathbf{M}_{d}\left(g_{j} \mathbf{y}\right) \succeq 0, \quad j=1, \ldots, m
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## Remarkably,

the Necessary \& Sufficient conditions ( $\star$ ) for existence of a representing measure are stated only in terms of countably many LINEAR MATRIX INEQUALITIES (LMI) on the sequence $y!($ No mention of the unknown representing measure in the conditions.)

맚 Moment matrix $\mathbf{M}_{1}(\mathbf{y})$ in dimension 2 with $d=1$


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the Necessary \& Sufficient conditions ( $\star$ ) for existence of a representing measure are stated only in terms of countably many LINEAR MATRIX INEQUALITIES (LMI) on the sequence $y!($ No mention of the unknown representing measure in the conditions.)

咷 Moment matrix $\mathbf{M}_{1}(\mathbf{y})$ in dimension 2 with $d=1$ :

$$
\mathbf{M}_{1}(\mathbf{y})=\left(\begin{array}{cccc} 
& 1 & x_{1} & x_{2} \\
1 & y_{00} & y_{10} & y_{01} \\
x_{1} & y_{10} & y_{20} & y_{11} \\
x_{2} & y_{01} & y_{11} & y_{02}
\end{array}\right)
$$

## lose localizing matrix $\mathbf{M}_{1}(g \mathbf{y})$ in dimension 2 with $d=1$ and $g(\mathbf{x})=1-x_{1}^{2}-x_{2}^{2}$ :

$$
\left(\begin{array}{cccc} 
& 1 & x_{1} & x_{2} \\
1 & y_{00}-y_{20}-y_{02} & y_{10}-y_{30}-y_{12} & y_{01}-y_{21}-y_{03} \\
x_{1} & y_{10}-y_{30}-y_{12} & y_{20}-y_{40}-y_{22} & y_{11}-y_{31}-y_{13} \\
x_{2} & y_{01}-y_{21}-y_{03} & y_{11}-y_{31}-y_{13} & y_{02}-y_{22}-y_{04}
\end{array}\right)
$$

ALGEBRAIC SIDE
POSitivity ON $K$



- In fact, polynomials NONNEGATIVE ON A SET $\mathbf{K} \subset \mathbb{R}^{n}$ are ubiquitous. They also appear in many important applications (outside optimization),


## ... modeled as

particular instances of the so called Generalized Moment Problem, among which: Probability, Optimal and Robust Control, Game theory, Signal processing, multivariate integration, etc.

## GMP: The primal view

The GMP is the infinite-dimensional LP:

$$
\inf _{\mu_{i} \in M\left(\mathbf{K}_{i}\right)}\left\{\sum_{i=1}^{s} \int_{\mathbf{K}_{i}} f_{i} d \mu_{i}: \quad \sum_{i=1}^{s} \int_{\mathbf{K}_{i}} h_{i j} d \mu_{i} \geqq b_{j}, \quad j \in J\right\}
$$

with $M\left(\mathbf{K}_{i}\right)$ space of Borel measures on $\mathbf{K}_{i} \subset \mathbb{R}^{n_{i}}, i=1, \ldots, s$.

## GMP: The dual view

The DUAL GMP* is the infinite-dimensional LP:

$$
\sup _{\lambda_{j}}\left\{\sum_{j \in J} \lambda_{j} b_{j}: \quad f_{i}-\sum_{j \in J} \lambda_{j} h_{i j} \geq 0 \text { on } \mathbf{K}_{i}, \quad i=1, \ldots, s\right\}
$$

## And one can see that

## the constraints of GMP* state that the functions


must be NONNEGATIVE on certain sets $\mathrm{K}_{i}, i=1$,

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$$

## And one can see that ...

the constraints of GMP* state that the functions

$$
\mathbf{x} \mapsto f_{i}(\mathbf{x})-\sum_{j \in J} \lambda_{j} h_{i j}(\mathbf{x})
$$

must be NONNEGATIVE on certain sets $\mathbf{K}_{i}, i=1, \ldots, s$.

## The moment-SOS hierarchy

- is an iterative numerical scheme to (help) solve the GMP.
- It consists of using a certain type of positivity certificate (e.g., Putinar's certificate) in potentially any application where such a characterization is needed.
- Global optimization is only one example.

of increasing size!.


## The moment-SOS hierarchy

- is an iterative numerical scheme to (help) solve the GMP.
- It consists of using a certain type of positivity certificate (e.g., Putinar's certificate) in potentially any application where such a characterization is needed.
- Global optimization is only one example.

> In many situations this amounts to solving a HIERARCHY of SEMIDEFINITE PROGRAMS
... of increasing size!.


- Has already been proved successful in applications with modest problem size, notably in optimization, control, robust control, optimal control, estimation, computer vision, etc.

중 If sparsity then problems of larger size can be addressed

Global optimization


Volume of semialgebraic set


Reachable set


Super resolution

Optimal control


Region of attraction


Maximum invariant sets


PDE analysis \& control


## Example: Global optimization

# Global OPTIM $\rightarrow f^{*}=\inf _{\mathbf{x}}\{f(\mathbf{x}): \mathbf{x} \in \mathbf{K}\}$ is the SIMPLEST example of the GMP 



咦 A GMP with only one unknown measure $\mu$ and only one moment-constraint $\int_{\mathbf{K}} 1 d \mu=1$

## Example: Global optimization

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## because ...

$$
f^{*}=\inf _{\mu \in \mathscr{M}(\mathbf{K})_{+}}\left\{\int_{\mathbf{K}} f d \mu: \int_{\mathbf{K}} 1 d \mu=1\right\}
$$

(19) A GMP with only one unknown measure $\mu$ and only one moment-constraint $\int_{\mathbf{K}} 1 d \mu=1$

## Example: Global optimization

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$$

ATP AMP with only one unknown measure $\mu$ and only one moment-constraint $\int_{\mathbf{K}} 1 d \mu=1$

## Remember also that for the GLOBAL minimum :

$$
f^{*}=\sup _{\lambda}\{\lambda: f(\mathbf{x})-\lambda \geq 0 \quad \forall \mathbf{x} \in \mathbf{K}\}
$$

## Then for each $t$ solve:

$$
\begin{gathered}
\rho_{t}=\sup _{\lambda, \sigma_{j}}\left\{\lambda: f(\mathbf{x})-\lambda=\sigma_{0}(\mathbf{x})+\sum_{j=1}^{m} \sigma_{j}(\mathbf{x}) g_{j}(\mathbf{x}), \forall \mathbf{x} \in \mathbb{R}^{n}\right. \\
\left.\operatorname{deg}\left(\sigma_{j} g_{j}\right) \leq t, \quad j=0, \ldots, m\right\}
\end{gathered}
$$

鲒 $\rho_{t} \leq \rho_{t+1} \leq f^{*}$ for all $t$ and $\rho_{t} \uparrow f^{*}$ as $t \rightarrow \infty$.

## Alternatively, for each $t$ solve:

$$
\left.\left.\begin{array}{rl}
\rho_{t}^{*}=\inf _{y}\left\{L_{y}(f): \quad \text { (think of } \int f d \mu\right. \text { ) } \\
y_{0}=1 \\
\mathbf{M}_{t}(y) & \succeq 0 \\
\mathbf{M}_{t-t_{j}}\left(g_{j} y\right) & \succeq 0 \quad \forall j=1, \ldots, m
\end{array}\right\} \Leftarrow y_{\alpha}=\int_{\mathbf{K}} x^{\alpha} d \mu\right\}
$$

## Theorem (Lass 2000)

ㅁㅏㅜㅇ $\square$
Moreover, generically $\rho_{t}^{*}=f^{*}$ and one may extract global minimizers from the optimal (truncated moment) solution y

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呢 $\rho_{t} \leq \rho_{t}^{*} \leq f^{*}$ for all $t$ and $\rho_{t}^{*} \uparrow f^{*}$ as $t \rightarrow \infty$.
Moreover, generically $\rho_{t}^{*}=f^{*}$ and one may extract global minimizers from the optimal (truncated moment) solution $y^{*}$.

Ex: Consider the optimization problem: $\min \{f(x): x \in[0,1]\}$ :

$$
x \mapsto f(x):=\sum_{j=1}^{4} a_{j} x^{j} ; \quad[0,1]=\{x: x(1-x) \geq 0\},
$$

## SDP relaxation

$$
\begin{aligned}
\operatorname{SDP} \quad f^{*}= & \min _{y}\left\{\sum_{j=1}^{4} a_{j} y_{j}:\left[\begin{array}{lll}
y_{0} & y_{1} & y_{2} \\
y_{1} & y_{2} & y_{3} \\
y_{2} & y_{3} & y_{4}
\end{array}\right] \succeq 0\right. \\
& \left.\operatorname{SDP}\left[\begin{array}{ll}
y_{1}-y_{2} & y_{2}-y_{3} \\
y_{2}-y_{3} & y_{3}-y_{4}
\end{array}\right] \succeq 0 ; y_{0}=1 .\right\}
\end{aligned}
$$



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\end{array}\right] \succeq 0 ; y_{0}=1 .\right\}
\end{aligned}
$$

잡 $y^{*}=\left(1, x^{*},\left(x^{*}\right)^{2},\left(x^{*}\right)^{3},\left(x^{*}\right)^{4}\right)$, and

$$
f(x)-f^{*}=\underbrace{\sigma_{0}(x)}+\underbrace{\sigma_{1}(x)} \quad x(1-x) .
$$

SOS of degree 4 SOS of degree 2
－Finite convergence is＂generic＂
－Extraction of minimizers from an optimal solution of the dual（linear algebra）
－If the problem is SOS－convex then convergence takes place at the first step of the hierarchy
－The＂same algorithm＂for many combinatorial optimization problems（just use $x_{i}^{2}=x_{i}$ to model boolean variables） which still provides better lower bounds than ad－hoc tailored algorithms．
［I马大ㅇ Has become key tool to prove／disprove Khot＇s Unique Games Conjecture in computational complexity．

唳 The NPA－hierarchy is a non－commutative version of the Moment－SOS hierarchy to address some quantitative problems in Quantum Information．

## A recovery issue

In static optimization, an optimal solution is a point $\mathbf{x}^{*} \in \mathbb{R}^{n}$.
Generically, some semidefinite relaxation at step $t$ of the Moment-SOS hierarchy is exact and:

To recover $\mathbf{x}^{*}$ from its optimal solution $\mathbf{y}^{*}=\left(y_{\alpha}^{*}\right)_{\alpha \in \mathbb{N}_{2 t}^{n}}$ can be done via a linear algebra subroutine.
(arer If $\mathbf{x}^{*}$ is unique then it is even trivial as $\mathbf{x}^{*}$ is just the subvector of degree-1 moments of $\mathbf{y}^{*}$

## However,

in many other problems like optimal control, PDE's, computational geometry, an optimal solution is a function $f: \Omega \rightarrow \mathbb{R}$ (e.g. a trajectory $\{\mathbf{x}(t): t \in[0,1]\})$
... and the Moment-SOS hierarchy provides a sequence of scalars $\left(\mu_{\alpha, j}\right)_{\alpha, j}$ which approximates moments

$$
\mu_{\alpha, j}^{*}=\int_{\Omega} \mathbf{x}^{\alpha} y^{j} d \mu^{*}(\mathbf{x}, y), \quad \alpha \in \mathbb{N}^{n}, j \in \mathbb{N}
$$

of the measure $d \mu^{*}(\mathbf{x}, y)=\delta_{\{f(\mathbf{x})\}}(d y) \phi(d \mathbf{x})$ whose support IS the graph $\{(\mathbf{x}, f(\mathbf{x})): \mathbf{x} \in \Omega\}$ of the optimal solution $f$.

For instance in optimal control (OC), one uses the weak formulation of OC
IT웅 infinite-dimensional LP on occupation (Young) measure $\mu$

- Controlled dynamics of OC

맙 linear constraints on moments of $\mu$ via integration of polynomial test functions

- integral cost functional

四 linear criterion $\langle h, \mu\rangle$ on $\mu$.

- state/control constraints $=$ support constraints on $\mu$四 semidefinite conditions on moments of $\mu$ (by Putinar theorem)


Then it remains to extract $f$ from knowledge of the $\left(\mu_{\alpha, j}\right)$...
ㄴㅏㅜㄱ This can be done by several techniques (including $L^{2}$-polynomial approximation via a standard application of the Christoffel-Darboux kernel) not detailed here.

呢 We claim that a non-standard application of the CD kernel provides a simple and easy to use tool (with no optimization involved) which can help solve problems not only in data analysis, but also in approximation and interpolation of (possibly discontinuous) functions. In particular one is able to recover a discontinuous function with no Gibbs phenomenon.


Interpolation


Recovery


## Part two:

## The Christoffel function

## The <br> Christoffel-Darboux Kernel for Data Analysis

Jean Bernard Lasserre, Edouard Pauwels and Mihai Putinar



Jean B. Lasserre*
Moments, Positive Polynomials, and the Christoffel Function

## Motivation

Consider the following cloud of 2D-points (data set) below


The red curve is the level set

$$
S_{\gamma}:=\left\{\mathbf{x}: Q_{d}(\mathbf{x}) \leq \gamma\right\}, \quad \gamma \in \mathbb{R}_{+}
$$

of a certain polynomial $Q_{d} \in \mathbb{R}\left[x_{1}, x_{2}\right]$ of degree $2 d$.

Notice that $S_{\gamma}$ captures quite well the shape of the cloud.

## Not a coincidence!

四 Surprisingly, low degree $d$ for $Q_{d}$ is often enough to get a pretty good idea of the shape of $\Omega$ (at least in dimension $p=2,3)$


## Cook up your own convincing example

Perform the following simple operations on a preferred cloud of $2 D$-points: So let $d=2, p=2$ and $s(d)=\binom{p+d}{p}$.

- Let $\mathbf{v}_{d}(\mathbf{x})^{T}=\left(1, x_{1}, x_{2}, x_{1}^{2}, x_{1} x_{2}, \ldots, x_{1} x_{2}^{d-1}, x_{2}^{d}\right)$. be the vector of all monomials $x_{1}^{i} x_{2}^{j}$ of total degree $i+j \leq d$
- Form the real symmetric matrix of size $s(d)$

$$
\mathbf{M}_{d}:=\frac{1}{N} \sum_{i=1}^{N} \mathbf{v}_{d}(\mathbf{x}(i)) \mathbf{v}_{d}(\mathbf{x}(i))^{T}
$$

where the sum is over all points $(\mathbf{x}(i))_{i=1 \ldots, N} \subset \mathbb{R}^{2}$ of the data set.

洈 Note that $\mathbf{M}_{d}$ is the MOMENT-matrix $\mathbf{M}_{d}\left(\mu^{N}\right)$ of the empirical measure

$$
\mu^{N}:=\frac{1}{N} \sum_{i=1}^{N} \delta_{\mathbf{x}(i)}
$$

associated with a sample of size $N$, drawn according to an unknown measure $\mu$.
[10 supported at the point $\mathbf{x}(i)$ of $\mathbb{R}^{2}$.

Recall that the moment matrix $\mathbf{M}_{d}(\mu)$ is real symmetric with rows and columns indexed by $\left(\mathbf{x}^{\alpha}\right)_{\alpha \in \mathbb{N}_{d}^{p}}$, and with entries

$$
\mathbf{M}_{d}(\mu)(\alpha, \beta):=\int_{\Omega} \mathbf{x}^{\alpha+\beta} d \mu=\mu_{\alpha+\beta}, \quad \forall \alpha, \beta \in \mathbb{N}_{d}^{p} .
$$

망 Illustrative example in dimension 2 with $d=1$ :

$$
\mathbf{M}_{1}(\mu):=\left(\begin{array}{cccc} 
& 1 & X_{1} & X_{2} \\
1 & \mu_{00} & \mu_{10} & \mu_{01} \\
X_{1} & \mu_{10} & \mu_{20} & \mu_{11} \\
X_{2} & \mu_{01} & \mu_{11} & \mu_{02}
\end{array}\right)
$$

is the moment matrix of $\mu$ of "degree $d=1$ ".

- Next, form the SOS polynomial:

$$
\begin{aligned}
& \mathbf{x} \mapsto Q_{d}(\mathbf{x}):=\mathbf{v}_{d}(\mathbf{x})^{T} \mathbf{M}_{d}^{-1}\left(\mu^{N}\right) \mathbf{v}_{d}(\mathbf{x}) . \\
& =\left(1, x_{1}, x_{2}, x_{1}^{2}, \ldots, x_{2}^{d}\right) \mathbf{M}_{d}^{-1}\left(\mu^{N}\right)\left(\begin{array}{c}
1 \\
x_{1} \\
x_{2} \\
x_{1}^{2} \\
\ldots \\
x_{2}^{d}
\end{array}\right)
\end{aligned}
$$

- Plot some level sets

$$
S_{\gamma}:=\left\{\mathbf{x} \in \mathbb{R}^{2}: Q_{d}(\mathbf{x})=\gamma\right\}
$$

for some values of $\gamma$, the thick one representing the particular value $\gamma=\binom{2+d}{2}$.

The Christoffel function $\Lambda_{d}: \mathbb{R}^{p} \rightarrow \mathbb{R}_{+}$is the reciprocal

$$
\mathbf{x} \mapsto Q_{d}(\mathbf{x})^{-1}, \quad \forall \mathbf{x} \in \mathbb{R}^{p}
$$

of the SOS polynomial $Q_{d}$.

## 맚ํ It has a rich history in Approximation theory and Orthogonal Polynomials.

값 Among main contributors: Nevai, Totik, Króo, Lubinsky, Simon,

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Let $\Omega \subset \mathbb{R}^{p}$ be the compact support of $\mu$ with nonempty interior, and $\left(P_{\alpha}\right)_{\alpha \in \mathbb{N}^{p}}$ be a family of orthonormal polynomials w.r.t. $\mu$.

The vector space $\mathbb{R}[\mathbf{x}]_{d}$ viewed as a subspace of $L^{2}(\mu)$ is a Reproducing Kernel Hilbert Space (RKHS). Its reproducing kernel

$$
(\mathbf{x}, \mathbf{y}) \mapsto K_{d}^{\mu}(\mathbf{x}, \mathbf{y}):=\sum_{|\alpha| \leq d} P_{\alpha}(\mathbf{x}) P_{\alpha}(\mathbf{y}), \quad \forall \mathbf{x}, \mathbf{y} \in \mathbb{R}^{p},
$$

is called the Christoffel-Darboux kernel.

## The reproducing property

$$
\mathbf{x} \mapsto q(\mathbf{x})=\int_{\Omega} K_{d}^{\mu}(\mathbf{x}, \mathbf{y}) q(\mathbf{y}) d \mu(\mathbf{y}), \quad \forall q \in \mathbb{R}[\mathbf{x}]_{d}
$$

榢 useful to determinate the best degree-d polynomial approximation

$$
\inf _{q \in \mathbb{R}[\mathbf{x}]_{d}}\|f-q\|_{L^{2}(\mu)}
$$

of $f$ in $L^{2}(\mu)$. Indeed:

$$
\begin{aligned}
\mathbf{x} \mapsto \widehat{f_{d}}(\mathbf{x}) & :=\sum_{\alpha \in \mathbb{N}_{d}^{p}}(\overbrace{\left.\int_{\Omega} f(y) P_{\alpha}(y) d \mu\right)}^{\widehat{f_{d, \alpha}}} P_{\alpha}(\mathbf{x}) \in \mathbb{R}[\mathbf{x}]_{d} \\
& =\arg \min _{q \in \mathbb{R}[\mathbf{x}]_{d}}\|f-q\|_{L^{2}(\mu)}
\end{aligned}
$$

## Theorem

The Christoffel function $\wedge_{d}^{\mu}: \mathbb{R}^{p} \rightarrow \mathbb{R}_{+}$is defined by:

$$
\xi \mapsto \Lambda_{d}^{\mu}(\xi)^{-1}=\sum_{|\alpha| \leq d} \boldsymbol{P}_{\alpha}(\xi)^{2}=K_{d}^{\mu}(\xi, \xi), \quad \forall \xi \in \mathbb{R}^{p},
$$

and it also satisfies the variational property:

$$
\Lambda_{d}^{\mu}(\xi)=\min _{P \in \mathbb{R}[\mathbf{x}]_{d}}\left\{\int_{\Omega} P^{2} d \mu: P(\xi)=1\right\}, \quad \forall \xi \in \mathbb{R}^{p}
$$

맙 Alternatively

$$
\Lambda_{d}^{\mu}(\xi)^{-1}=\mathbf{v}_{d}(\xi)^{T} \mathbf{M}_{d}(\mu)^{-1} \mathbf{v}_{d}(\xi), \quad \forall \xi \in \mathbb{R}^{p}
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$$

[1] Importantly, and crucial for applications, the Christoffel function identifies the support $\Omega$ of the underlying measure $\mu$.

## Theorem

Let the support $\Omega$ of $\mu$ be compact with nonempty interior. Then:

- For all $\mathbf{x} \in \operatorname{int}(\Omega): K_{d}^{\mu}(\mathbf{x}, \mathbf{x})=O\left(d^{p}\right)$.
- For all $\mathbf{x} \in \operatorname{int}\left(\mathbb{R}^{p} \backslash \Omega\right): K_{d}^{\mu}(\mathbf{x}, \mathbf{x})=\Omega(\exp (\alpha d))$ for some $\alpha>0$.

줎 $\operatorname{In}$ particular, as $d \rightarrow \infty$,

$$
d^{p} \wedge_{d}^{\mu}(\mathbf{x}) \rightarrow 0 \text { very fast whenever } \mathbf{x} \notin \Omega .
$$

## Growth rates for $K_{d}^{\mu}(\mathbf{x}, \mathbf{x})=\Lambda_{d}^{\mu}(\mathbf{x})^{-1}$.



## Some other properties

- Under some (restrictive) assumption on $\Omega$ and $\mu$

$$
\lim _{d \rightarrow \infty} s(d) \wedge_{d}^{\mu}(\xi)=f_{\mu}(\xi) \omega(\xi)^{-1}
$$

where $\omega$ is the density of an equilibrium measure intrinsically associated with $\Omega$.
For instance with $p=1$ and $\Omega=[-1,1], \omega(\xi)=\sqrt{1-\xi^{2}}$.

- If $\mu$ and $\nu$ have same support $\Omega$ and respective densities $f_{\mu}$ and $f_{\nu}$ w.r.t. Lebesgue measure on $\Omega$, positive on $\Omega$, then:


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- If $\mu$ and $\nu$ have same support $\Omega$ and respective densities $f_{\mu}$ and $t_{\nu}$ w.r.t. Lebesgue measure on $\Omega$, positive on $\Omega$, then:

$$
\lim _{d \rightarrow \infty} \frac{\Lambda_{d}^{\mu}(\xi)}{\Lambda_{d}^{\nu}(\xi)}=\frac{f_{\mu}(\xi)}{f_{\nu}(\xi)}, \quad \forall \xi \in \Omega
$$

呢 useful for density approximation

##  $\xi$ such that $\Lambda_{d}^{\mu^{N}}(\xi)<\binom{p+d}{p}^{-1}$.

맚ㄱ Such a strategy (even with relatively low degree $d$ ) is as efficient as more elaborated techniques, (the degree d), and
(ax) Lass. \& Pauwels (2016) Sorting out typicality via the inverse moment matrix SOS polynomial, Lass. \& Pauwels (2019) The empirical Christoffel function with applications in data analysis,

Lass. (2022) On the Christoffel function and classification in data analysis.
(27) For instance one may decide to classify as outliers all points $\xi$ such that $\Lambda_{d}^{\mu^{N}}(\xi)<\binom{p+d}{p}^{-1}$.

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吡 Lass. \& Pauwels (2016) Sorting out typicality via the inverse moment matrix SOS polynomial, NIPS 2016.
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1439-1468
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## Manifold learning

A measure $\mu$ on compact set $\Omega$ is completely determined by its moments and therefore it should not be a surprise that its moment matrix $\mathbf{M}_{d}(\mu)$ contains a lot of information.

嗗 We have already seen that its inverse $\mathbf{M}_{d}(\mu)^{-1}$ defines the Christoffel function.

When $\mu$ is degenerate and its support $\Omega$ is contained in a zero-dimensional real algebraic variety $V$ then the kernel of $\mathbf{M}_{d}(\mu)$ identifies the generators of a corresponding ideal of $\mathbb{R}[\mathbf{x}]$ (the vanishing ideal of $V$ ).

For instance let $\Omega \subset \mathbb{S}^{p-1}$ (the Euclidean unit sphere of $\mathbb{R}^{p}$ )


Then the kernel of $\mathbf{M}_{d}(\mu)$ contains vectors of coefficients of polynomials in the ideal generated by the quadratic polynomial $\mathbf{x} \mapsto g(\mathbf{x}):=1-\|\mathbf{x}\|^{2}$.

## In fact and remarkably,

$$
\operatorname{rank} \mathbf{M}_{d}(\mu)=p(d)
$$

for some univariate polynomial $p$ (the Hilbert polynomial associated with the algebraic variety) which is of degree $t$ if $t$ is the dimension of the variety.

For instance $t=p-1$ if the support is contained in the sphere $\mathbb{S}^{p-1}$ of $\mathbb{R}^{p}$.
 from empirical moments and the Christoffel function, Found. Comput. Math. 21, pp. 243-273.

四 Again this illustrates how quite sophisticated concepts of algebraic geometry are hidden and encapsulated in the moment matrix $\mathbf{M}_{d}(\mu)$.

뭆 They can be exploited to extract various useful information on the data set.
ney In addition, extraction of this information can be done via quite simple linear algebra techniques.

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## 衡 However

## for non modest dimension of data, matrix inversion of $\mathbf{M}_{d}^{-1}$ does not scale well ...

## 당 O O the other hand

for evaluation $\Lambda_{d}^{\mu}(\xi)$ at a point $\xi \in \mathbb{R}^{p}$, the variational formulation

is the simple quadratic programming problem.

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 for evaluation $\Lambda_{d}^{\mu}(\xi)$ at a point $\xi \in \mathbb{R}^{p}$, the variational formulation$$
\Lambda_{d}^{\mu}(\xi)=\min _{P \in \mathbb{R}[\mathbf{x}]_{d}}\left\{\int_{\Omega} P^{2} d \mu: P(\xi)=1\right\}, \quad \forall \xi \in \mathbb{R}^{p} .
$$

is the simple quadratic programming problem.

$$
\min _{p \in \mathbb{R}^{s(d)}}\left\{p^{T} \mathbf{M}_{d} p: \quad \mathbf{v}_{d}(\xi)^{T} p=1\right\}
$$

which can be solved quite efficiently.

## The Christoffel function for approximation

A typical approach is to approximate $f:[0,1] \rightarrow \mathbb{R}$ in some function space, e.g. its projection on $\mathbb{R}[\mathbf{x}]_{n} \subset L^{2}([0,1])$ :

$$
x \mapsto \hat{f}_{n}(x):=\sum_{j=0}^{n}\left(\int_{0}^{1} f(y) L_{j}(y) d y\right) L_{j}(x)
$$

with an orthonormal basis $\left(L_{j}\right)_{j \in \mathbb{N}}$ of $L^{2}([0,1])$.

## BUT ...



Ex: Chebyshev interpolant
幈 Typical Gibbs phenomenon occurs.

Alternative Positive Kernels with better convergence properties have been proposed, still in the same framework:

Féjer, Jackson kernels, etc.

- Reproducing property of the CD kernel is LOST
- Preserve positivity (e.g when approximating a density)
- Better convergence properties than the CD kernel, in particular uniform convergence (for continuous functions) on arbitrary compact subsets


## An alternative approach, still via the CD-kernel

A counter-intuitive detour: Instead of considering $f:[0,1] \rightarrow \mathbb{R}$, and the associated measure

$$
d \mu(x):=f(\mathbf{x}) d x
$$

on the real line, whose support is $[0,1] \in \mathbb{R}$,

鲆 Rather consider the graph $\Omega \subset \mathbb{R}^{2}$ of $f$, i.e., the set
and the measure
$d \phi(x, y):=\delta_{f(x)}(d y) 1_{[0,1]}(x) d x$

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맚ㅇ Rather consider the graph $\Omega \subset \mathbb{R}^{2}$ of $f$, i.e., the set

$$
\Omega:=\{(x, f(x)): \quad x \in[0,1]\} .
$$

and the measure

$$
d \phi(x, y):=\delta_{f(x)}(d y) 1_{[0,1]}(x) d x
$$

on $\mathbb{R}^{2}$ with degenerate support $\Omega \subset \mathbb{R}^{2}$.

The Moment-SOS Hierarchy



Why should we do that as it implies going to $\mathbb{R}^{2}$ instead of staying in $\mathbb{R}$ ?

## 喁 ... because

- The support of $\phi$ is exactly the graph of $f$, and
- The CF $(x, y) \mapsto \Lambda_{n}^{\phi}(x, y)$ identifies the support of $\phi$ !

So suppose that we know the moments

$$
\phi_{i, j}=\int x^{i} y^{j} d \phi(x, y)=\int_{[0,1]} x^{i} f(x)^{j} d x, \quad i+j \leq 2 d
$$

and let $\varepsilon>0$ and $\lambda$ be the Lebesgue measure on $[0,1]$.

- 뭉 Compute the degree- $d$ moment matrix of $\phi$ :

$$
\mathbf{M}_{d}(\phi):=\int \mathbf{v}_{d}(x, y) \mathbf{v}_{d}(x, y)^{T} d \phi(x, y)
$$

- Compute the Christoffel function

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## Interpolation: same story

So suppose that you are given point evaluations $\left\{f\left(x_{i}\right)\right\}_{i \leq N}$ of an unknown function $f$ on $[0,1]$, and again let

$$
\mathbf{v}_{d}(x, y):=\left(1, x, y, x^{2}, x y, y^{2}, \ldots, x y^{d-1}, y^{d}\right) .
$$

- Compute the degree- $d$ empirical moment matrix:

$$
\mathbf{M}_{d}(\phi):=\sum_{i=1}^{N} \mathbf{v}_{d}\left(\left(x_{i}, f\left(x_{i}\right)\right) \mathbf{v}_{d}\left(x_{i}, f\left(x_{i}\right)\right)^{T}\right.
$$

of the empirical measure $d \phi(x, y):=\frac{1}{N} \sum_{i=1}^{N} \delta_{x(i), f(x(i))}$ on $\mathbb{R}^{2}$, by one pass over the data

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## Choosing

$$
\varepsilon:=2^{3-\sqrt{d}}
$$

ensures convergence properties for bounded measurable functions, e.g. pointwise on open sets with no point of discontinuity.

## Convergence properties as $d \uparrow$

- llo $L^{1}$-convergence
- (T) pointwise convergence on open sets with no point of discontinuity, and so almost uniform convergence.
- 挶 $L^{1}$-convergence at a rate $O\left(d^{-1 / 2}\right)$ for Lipschitz continuous $f$.

In non trivial exemples, the approximation is quite good with small values of $d$, and with no Gibbs phenomenon.



## Ex: Recovery

Below: Recovery of a (discontinuous) solution of the Burgers Equation from knowledge of approximate moments of the occupation measure supported on the solution.


Again note the central role played by the Moment Matrix!
S. Marx, E. Pauwels, T. Weisser, D. Henrion, J.B. Lass. Semi-algebraic approximation using Christoffel-Darboux kernel, Constructive Approximation, 2021

## Christoffel function and Positive polynomials

Let $\Omega \subset \mathbb{R}^{n}$ be the basic semi-algebraic set (with nonempty interior)

$$
\Omega:=\left\{\mathbf{x} \in \mathbb{R}^{n}: g_{j}(\mathbf{x}) \geq 0, \quad j=1, \ldots, m\right\}
$$

with $g_{j} \in \mathbb{R}[\mathbf{x}]_{d_{j}}$ and let $s_{j}=\left\lceil\operatorname{deg}\left(g_{j}\right) / 2\right\rceil$. Let $g_{0}=1$ with $s_{0}=0$.

With $t$ fixed, its associated quadratic module

$$
Q_{t}(\Omega):=\left\{\sum_{j=0}^{m} \sigma_{j} g_{j}: \quad \sigma_{j} \in \Sigma[\mathbf{x}]_{t-s_{j}}\right\} \subset \mathbb{R}[\mathbf{x}]
$$

is a convex cone with nonempty interior,

## and with dual convex cone of pseudo-moments

$$
Q_{t}(\Omega)^{*}:=\left\{y \in \mathbb{R}^{s(t)}: \mathbf{M}_{t-s_{j}}\left(g_{j} y\right) \succeq 0, \quad j=0, \ldots, m\right\},
$$

where $s(t)=\binom{n+t}{n}$.

## Notice that if $\mathbf{M}_{t}(y)^{-1} \succ 0$ for all $t$,

then one may define a family of polynomials $\left(P_{\alpha}\right)_{\alpha \in \mathbb{N}^{n}} \subset \mathbb{R}[\mathbf{x}]$ orthonormal w.r.t. $y$, meaning that

$$
L_{y}\left(P_{\alpha} \cdot P_{\beta}\right)=\delta_{\alpha=\beta}, \quad \alpha, \beta \in \mathbb{N}^{n}
$$

and exactly as for measures, the Christoffel function $\Lambda_{t}^{y}$

$$
\mathbf{x} \mapsto \Lambda_{t}^{y}(\mathbf{x})^{-1}:=\sum_{|\alpha| \leq t} P_{\alpha}(\mathbf{x})^{2}
$$

## Theorem

For every $p \in \operatorname{int}\left(Q_{t}(\Omega)\right)$ there exists a sequence of pseudo-moments $y \in \operatorname{int}\left(Q_{t}(\Omega)^{*}\right)$ such that

$$
\begin{aligned}
p(\mathbf{x}) & =\sum_{j=0}^{m}\left(\mathbf{v}_{t-s_{j}}(\mathbf{x})^{T} \mathbf{M}_{t}\left(g_{j} y\right)^{-1} \mathbf{v}_{t-s_{j}}(\mathbf{x})\right) g_{j}(\mathbf{x}) \\
& =\sum_{j=0}^{m} \Lambda_{t-s_{j}}^{g_{j} \cdot y}(\mathbf{x})^{-1} g_{j}(\mathbf{x})
\end{aligned}
$$

where $(g \cdot y)$ is the sequence of pseudo-moments

$$
(g \cdot y)_{\alpha}:=\sum_{\gamma} g_{\gamma} y_{\alpha+\gamma}, \quad \alpha \in \mathbb{N}^{n} \quad\left(\text { if } g(\mathbf{x})=\sum_{\gamma} g_{\gamma} \mathbf{x}^{\gamma}\right)
$$

In addition $L_{y}(p)=\sum_{j=0}^{m}\binom{n+t-s_{j}}{n}$.


## The proof combines

- 搌 a result by Nesterov on a one-to-one correspondence between $\operatorname{int}\left(Q_{t}(\Omega)\right)$ and $\operatorname{int}\left(Q_{t}(\Omega)^{*}\right)$, and
- 㺃 the fact that

$$
\mathbf{v}_{t-s_{j}}(\mathbf{x})^{T} \mathbf{M}_{t}\left(g_{j} y\right)^{-1} \mathbf{v}_{t-s_{j}}(\mathbf{x})=\Lambda_{t-s_{j}}^{g_{j} \cdot y}(\mathbf{x})^{-1} .
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## In other words:

If $p \in \operatorname{int}\left(Q_{t}(\Omega)\right)$ then in Putinar's certificate

$$
p=\sum_{j=0}^{m} \sigma_{j} g_{j}, \quad \sigma_{j} \in \mathbb{R}[\mathbf{x}]_{t-s_{j}}
$$

of positivity of $p$ on $\Omega$,


$$
\sigma_{j}(\mathbf{x}):=\Lambda_{t-s_{j}}^{g_{j} \cdot y}(\mathbf{x})^{-1}, \quad j=0, \ldots, m
$$

for some sequence of pseudo-moments $y \in \operatorname{int}\left(Q_{t}(\Omega)^{*}\right)$.

## 傕 In particular,

every SOS polynomial $p$ of degree $2 d$, in the interior of the SOS-cone, is the reciprocal of the CF of some linear functional $y \in \mathbb{R}[\mathbf{x}]_{2 d}^{*}$. That is:

$$
p(\mathbf{x})=\mathbf{v}_{d}(\mathbf{x})^{T} \mathbf{M}_{d}(\mathbf{y})^{-1} \mathbf{v}_{x}(\mathbf{x})=\Lambda_{d}^{y}(\mathbf{x})^{-1}, \quad \forall \mathbf{x} \in \mathbb{R}^{n} .
$$

## CF - Pell's equation - equilibrium measure

四 What is the link between $p \in \operatorname{int}\left(Q_{t}(\Omega)\right)$ and the mysterious linear functional $y$ ?

## Theorem

For some sets $\Omega, 1 \in \operatorname{int}\left(Q_{t}(\Omega)\right)$ and

$$
\begin{equation*}
1=\frac{1}{\sum_{j=0}^{m} s\left(t-t_{j}\right)} \sum_{j=0}^{m} \Lambda_{t-s_{j}}^{g_{j} \cdot \phi}(\mathbf{x})^{-1} g_{j}(\mathbf{x}) \tag{1}
\end{equation*}
$$

where $\phi$ is the equilibrium measure of $\Omega$.
(1) can be called a generalized polynomial Pell's equation satisfied by the CFs $\Lambda_{t-s_{j}}^{g_{j} \cdot \phi}(\mathbf{x})^{-1}$.

## Disintegration

Recall that if $\mu$ is a measure on a Borel set $\Omega:=X \times Y$, then it disintegrates as

$$
d \mu(x, y)=\underbrace{\hat{\mu}(d y \mid x)}_{\text {conditional }} \underbrace{\phi(d x)}_{\text {marginal }}
$$

with marginal $\phi$ on $X$ and conditional $\hat{\mu}(d y \mid x)$ on $Y$ given $x \in X$.

## Theorem (Lass (2022))

The Christoffel function $\Lambda_{d}^{\mu}(x, y)$ disintegrates into

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for some measure $\nu_{x, d}$ on $\mathbb{R}$.

## Crucial in the proof is the use of the previous duality result of Nesterov.

## THANK YOU!

